

PERFORMANCE MONITORING OF PDP-11 COMPUTERS

by

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A thesis submitted to the Faculty of Graduate Studies
and Research in partial fulfillment of the requirements
for the degree of Master of Science.

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MONTREAL, P.Q., CANADA
August, 1976

ABSTRACT

Various system monitors and performance analysis methods have been surveyed. Software and hardware techniques are compared with reference to certain applications. An investigation of a general model for performance measurement tools has resulted in the development of HIMOS, a hybrid interactive monitor system. It is designed for a PDP-11 minicomputer installation. The software and hardware parts of the monitor are controlled from another PDP-11 computer. The system can work in sampling and event driven mode. Its ability to interfere actively with the performance of the measured system can be used for dynamic changes of the measurement session. The control system provides online and offline analysis of event reports. Monitor output is directed to a graphic display unit and can be concurrently recorded on tape or disk. The usage of the monitor in several application fields is illustrated through examples.

RESUME

Plusieurs moniteurs de système d'ordinateurs ainsi que des méthodes d'analyse de performance ont été étudiés. Des techniques de programmation et d'électronique ont été comparées. Un modèle général d'instrument de mesure de performance a été examiné et il en a résulté le développement de HIMOS, un système de moniteur hybride et interactif. Il est conçu pour une installation de mini-ordinateur PDP-11. Le programme et l'électronique du moniteur sont contrôlés par un autre PDP-11. Le système peut prendre des échantillons ou réagir sur des événements. Il peut intervenir sur la performance de l'ordinateur observé, changeant ainsi la séquence de mesure. Le système de contrôle permet l'analyse des rapports d'événements en ligne ouverte ou fermée. Les résultats sont dirigés sur un écran cathodique et peuvent être enregistrés simultanément sur bande ou sur disque. L'emploi du moniteur est illustré par différents exemples.

ACKNOWLEDGEMENT

I would like to thank Professor N. Marovac for getting me interested in this project. Professor G.F.G. Patzer as my thesis supervisor provided guidance concerning the scope of the text. His encouragement and constructive criticism has helped me to bring this work to completion. I wish to thank Arnold Epstein for designing the hardware part of this project and also for the many fruitful discussions we have had together.

Finally I thank Ginette for her assistance in typing this thesis and for her patience and encouragement during the past year.

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CHAPTER 1

Introduction to Performance Evaluation

1.1 Need for Evaluation

Prior to the description of the implementation and application of a practical tool for performance measurements in a computer system a discussion of the more general aspect of performance evaluation for which those tools are required will be presented.

The need for performance evaluation has not always been as apparent as it is today. First generation computer systems had a very simple structure and the central processor unit (CPU) was the dominant part of the installation. Peripheral devices such as magnetic tape drives, paper tape readers and punches were mainly used to maintain the traffic of program, data, and results between the system and the outside world. The devices were under direct control of the CPU. Computation and I/O were not overlapping because the CPU was dedicated at any time to one or the other activity. For scientific applications the power of a system was mainly dependent on the CPU speed. For commercial use the processing speed was generally limited by the performance of peripheral equipment.

The manager of a computer centre was able to charge his customers by the clock on the wall. His main concern was to insure the availability of his system. Availability can be defined as a measure of the likelihood that a system is operating properly at a given moment, or as the percentage

of time during which it is working properly.

With the advent of second and third generation systems, the complexity of information processing systems has grown considerably. For a human observer it became impossible to determine which job was executing at a certain time in a multiprogrammed environment. A scientific user could no longer expect the execution time of his program to decrease by factor two when the processing speed of the CPU was doubled. The interactions between hardware, system software, and user programs are far too complex as to allow this kind of simplified deduction.

It was not only for accounting purposes that a whole industry for performance measurement tools came into existence. Very soon it was felt that it was not enough just to write system and user software with the only goal to make it work. As the cost for maintaining computer systems went up one became more concerned about their efficient use. The performance of an existing installation had to be increased and criteria for future hard and software development had to be gathered. As a large variety of systems is offered on the market, there must be a method to compare the different products with regard to the intended application.

With the growing complexity of modern operating systems it is difficult to fully understand their behaviour. A system analyst has to know which parameters affect which aspect of the performance if he wants to tune a system. He may be confronted with side effects which he did not expect. His decisions have to be based on information that he can only

obtain through performance measurements.

1.2 Some Definitions

In order to define performance of information processing systems some terms have to be introduced first.

Throughput is the rate at which a given workload can be handled by the system. The workload can consist of payroll processing where throughput would measure the number of payroll statements per hour. If we consider a network node, this measure indicates the number of messages switched per hour. In a University environment we would be interested in the number of ALGOL or FORTRAN statements compiled per minute. A detailed discussion of workload characterization is given by Ferrari [1].

Depending on the application, another measure can be more important: Turnaround time is defined as the delay between the presentation of input to a system and the receipt of output from it [2, p.13]. Turnaround time is mostly used in connection with batch processing systems where it can specify the time span from putting cards into the hopper of a card reader and getting a complete listing on the line printer. Other papers include human factors such as operator handling time, think time, debugging time etc. into their definition (see for example [3, p.262]). But in this case we should speak about the time which is required for one test and debugging cycle.

Closely related to turnaround time is the term "response time". However this applies more to real time applications. For industrial process control response time could measure

the time elapsed between the arrival of a signal indicating the position of a work piece and the emission of some correcting signal from the computer. In a time sharing environment, response time gives the period between a user's request from the terminal and the moment at which the system is able to service this request.

It should be pointed out, that system overhead (the processing time which is spent in the operating system as opposed to the time spent in a user program) is not a fundamental measure of performance. The goal of performance improvement is not to decrease overhead but to increase throughput and to decrease response time. Overhead can be used as a descriptive parameter but it does not necessarily affect performance.

In the previous examples, the term "performance" has been used in different ways. Intuitively we can have a certain feeling for what it stands. However performance as an independent entity does not exist. It can be discussed only in the context of a specific application or a set of applications.

For example a high level language compiler may be judged by its compilation speed. This is an adequate measure in an educational environment whereas other users are much more concerned about the quality of the object code produced or the storage requirements for the compiler.

In a commercial computer center, performance is related to cost effectiveness. However this measure has to be interpreted carefully with regard to the individual application [3]. A fairly basic measure of performance,

giving the mean cost of delay to each job is suggested by Greenberger [4]. Generally we want some measure of effectiveness that is related to the ability to process jobs where jobs can be user programs or operating system tasks.

1.3 Performance Evaluation Techniques

One approach to performance evaluation is to describe the target system by means of an equivalent model. A model is an abstraction of the real system containing only the significant variables and relations. It is usually much simpler than the system it models. On one hand, the input to the model can consist of data about the actual workload of an existing computer installation. On the other hand information about the expected task volume of a projected computer system can be used. As this technique has to employ some simplifications in order to keep the volume of the model at a feasible size, the result can only give an approximation of the real performance. This method is widely used in the field of system design and research.

More precise results about the performance of a working system can be obtained from the analysis of performance measurements. Especially for the application in computing centres where observations about the actual system usage are necessary, this technique has known a growing popularity during the last ten years.

Calingaert [2] has given a very good review of the different performance evaluation techniques. The decision to use one or the other of these methods depends on the individual application.

1.3.1 Analytic Models

In analytic models, the relation between system variables and performance parameters is described by a set of mathematical equations. In accordance with real system behaviour it is assumed that events such as the arrival of a new job at the input appear in a stochastic manner. This means that variables can change their values in a random fashion but the range of values and the probability of the occurrence of a certain value is known.

We may for example know a given workload for the system but we cannot predict the sequence in which the jobs of this workload will arrive at the input. Let us consider a simple model of a batch processing system given by Graham [5,p.406]. In figure 1.1 new jobs which enter the system at the input are stored in a first-in, first-out (FIFO) queue. Time is divided into units called quanta, each of which is exactly Q seconds long. At the end of each quantum a new job can arrive at the tail of the queue. The job at the head of the queue is allocated to the processor where it is executed until completion. Once the execution is terminated it leaves the system at the output and the next job is shifted to the head of the queue. If the queue is empty, the processor remains idle.



Figure 1.1

To construct an analytical model, the job arrival time and the execution times for each job are expressed by probability distributions rather than in absolute terms. For example it can be assumed that a new job arrives in the system at the end of each quantum Q with probability pQ . The resulting job arrival distribution is a special case of the discrete Bernoulli distribution. The job's execution time can be chosen independently from a geometric distribution, s , assuming that each job requires an exact multiple of Q for processing. The probability that a job requires n quanta for execution is then expressed by $s(n)$.

We can modify the previous model to study the round-robin scheduling policy which may be used in time-sharing systems (fig.1.2). In this case the processor is only allocated to one job for exactly one quantum of time. If the execution was not completed at the end of one quantum the job is fed back to the end of the queue. There it waits together with other incomplete or newly arrived jobs for its next time slice in the processor. Assuming that a job's execution time is exactly nQ , it will have travelled n times through the system before it has completed.

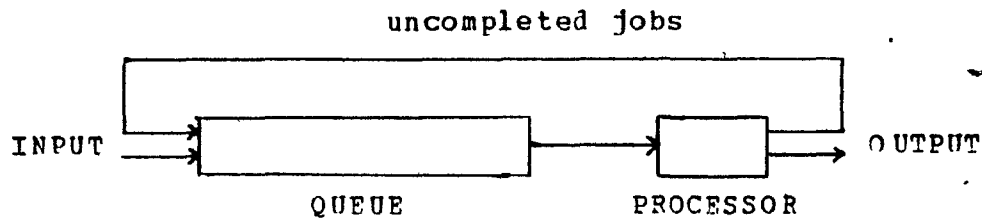


Figure 1.2

Another version where this model is extended by a second queue to buffer the incoming jobs is given by Blatney et. al. [6,p.412]. According to a selection algorithm, jobs are chosen from this auxiliary store to be placed at the end of the queue of the previous example which now serves as intermediate storage area for uncompleted jobs. Further levels of detail can easily be added to the model, but as is shown in the above mentioned paper, a deep analysis of the model can already become quite complex.

For the first two examples, Kleinrock [7] has derived equations for the expected total time that a job spends in the system under the above stated assumptions.

A more complex model which analyzes the system throughput as a function of input/output and processing overlap is described by Hellerman and Smith [8]. Several simplifying assumptions had to be adopted to keep the mathematical framework within practical limits. This shows already a severe drawback from the technique of analytical models. For complex systems, the relation between different variables may even not be expressable by mathematical functions.

Analytic models are most useful for the evaluation of subsystems of limited size (e.g. drum paging efficiency

[9]). As they need a certain degree of detail for their descriptive parameters they are in general used for evaluation of existing systems. At earlier stages of system development where this information is not available, the simulation technique is more appropriate.

1.3.2 Logical Models

While for analytical models, the system structure has to be translated into mathematical terms the logical model reflects this structure more directly. There are several different ways to express a system structure.

The directed graph model is very similar to a flowchart of some software. System states are represented by nodes and transitions from one state into another are shown as directed arcs. More details on this technique can be found in the literature [10,11].

Another interesting approach to logical modelling is in the use of Markovian models. Foley [13] has described such a model, using a discrete time semi-Markov chain where the time interval between successive state transitions is fixed. The chain is described by a transition probability matrix whose elements are the probabilities of entering a certain state from the current state of the system. The transition probabilities can either represent expected or estimated values, or they can be obtained from measurements on the executing target system.

Simulation which is the most frequently used technique among logical models will be discussed in more detail in the following section.

The first step in simulation is to define a dynamic model of the investigated computer system. The model consists of classes of entities, the attributes of these classes, a set of activities, and a set of events. A job that is to be executed by the system represents, for example, one class of entities. It can be described by attributes such as memory requirement, I/O activities, execution time etc. Activities are, for example, arithmetic computations or waiting for I/O completion. A job that enters the system causes an event; I/O interrupts can also be classified as an event. Events have no duration but in general they induce some activity.

The basic concept of simulation is time. Although simulation time is usually not equal to the real system time, the simulator controls events and activities with regard to its own time scale. Thus the simulator follows the sequence of events which indicate changes of the system state over time. Models for simulation of digital computer systems increment time in discrete intervals whose length correspond to the simulation time between consecutive events.

To simulate a system, a set of different experiments is conducted on the model. In other words, a sequence of different workloads is presented to the input of the model and results about the system performance on this input are obtained.

Another way to drive the model is to measure certain parameters such as job arrival time, distribution of page faults etc. in a real system and to use this data as input to the model. During simulation, other performance

parameters can be varied to find their optimal values. In this case we speak about trace driven modelling.

If required, the simulation results of a series of experiments can be formulated as mathematical equations describing the relationship between system variables and performance parameters. This abstraction can then be used to carry out further experiments more economically. This indicates that there is no absolute separation between analytic models and simulation models. For example, some subcomponents of a system can also be described analytically within a simulator. Of course, it is much easier to modify some aspects of the system in a simulation model just by replacing attribute definitions than by modifying the complex equations of the other model. In simulation no attempt is made to isolate relations between variables; observations are merely carried out on how they change their values with time.

In practice, the task of creating a simulation model is simplified by special purpose simulation languages such as GPSS (General Purpose Simulation System) or CSS (Computer System Simulator) [14]. Those languages provide facilities for varying parameters, for different experiments.

Especially at the planning stage of computer system development the simulation model represents the most convenient technique for performance evaluation. It is not only the most flexible of all models but it also can be constructed to any level of detail. Concurrency of activities can easily be modelled whereas analytic models are not able to express this feature which is becoming more

and more important in modern computing systems.

However all modelling techniques have common disadvantages. The obtained results can only be as good as the model. Analytical models are often oversimplified to reduce their complexity whereas simulation models are very expensive. Simplifications and assumptions in less known areas of the system can lead to undetected distortion of results. The validity of a model can be examined by comparing the results with those obtained on a real system but of course this is only possible if this system already exists. As the performance of a system is a function of the input, the experiments have to be designed at least as carefully as the model.

1.3.3 Performance Evaluation Programs

The techniques described in this paragraph apply to installed operational systems. They are mainly used to compare performance of different systems or to verify expected improvements of system behaviour. Executing the same assembly run on two different machines, we may observe different execution times just by looking at our watch. But even by timing the runs with a clock with a resolution of one microsecond would not tell us why machine A is faster than machine B. If we did the same test with another sample job we could obtain the opposite result, computer B being the faster one. In order to achieve more objective results, artificial workloads consisting of special programs or sequences of programs have been developed to evaluate the performance of computing systems.

The simplest technique in the class of performance evaluation programs is the instruction mix. The frequency of a machine instruction is derived from a dynamic trace under real workload. In the instruction mix, each instruction execution time is then weighted by a coefficient based on the frequency measurements. A table of instruction weights for a scientific and for a commercial mix are given in [15]. Knowing the instruction execution times of the target CPU a value representing the CPU performance can easily be calculated. The result depends on the dynamic trace which determines the relative importance of each instruction with regard to the overall performance.

The kernel technique uses another attempt to define a representative workload. It is composed of one or more assembler routines, each of which solves a specific computational problem, e.g. matrix inversion or evaluation of a polynomial. The set of problems that form a kernel is considered to be representative of the predominant application of the system.

The kernel technique has the advantage in that it is more CPU independent than the instruction mix as it tests performance on a task level rather than on the instruction level. However an objective comparison of different CPU's depends on the objective programming. The advantages and disadvantages of the last two techniques are discussed further in [2].

So far, the I/O characteristics of a system have not been evaluated. Benchmark problems include this feature in their workload definition. A popular benchmark problem for

commercial installations is the file updating program [16]. A processor reads a master file from disk or tape, updates it according to a detail file and produces a report file containing a record of each transaction performed.

Many attempts have been made to standardize benchmarks, but in general, they still reflect a special application. This technique may be valuable to gain insight into the expected performance of a system mainly used for a particular application environment but it cannot give a global measure of system performance.

Synthetic jobs as proposed by Buchholz [17] are more flexible than benchmarks. In those programs, written in high level languages, various resources may be parametrically changed to fit the characteristics of a real workload. This is done by several tasks which are executed cyclically. As every task requests for another resource (memory, CPU-time, I/O channels, etc.), the loop parameters can be adjusted to the system size and the workload to be simulated.

The techniques described so far evaluate system performance under the condition that the test program is the only one executed in the system. In this case the execution time of the job is used for evaluation or comparison. For timesharing systems we would rather want to measure the response time under an artificial workload. This means several users have to be simulated submitting different tasks concurrently to the system. A description of an internal driver simulating interactive users for the MULTICS time sharing system is outlined in [18]. Each user is

simulated by a process which can be considered an interactive benchmark consisting of a fixed sequence of commands and fixed think times.

Criticizing performance evaluation programs, one can say they rather define performance than evaluate it. Yet, if well understood, they can give some information about performance with regard to the yielded application. If they are used along with performance measurement techniques (which will be discussed in the next chapter), the obtained results can help to improve system performance under a well defined workload.

CHAPTER 2

Performance Measurement Techniques

Performance measurement or monitoring can be defined as the process of obtaining useful information on the performance of software and software controlled computer hardware. The earliest measurement tools were developed around 1956. They were able to supply information such as job execution time or core utilization which was obtained by core resident software monitors. In 1958 IBM developed the first laboratory version of an external hardware device for performance measurement. Tables A.1 and A.2 give a historical overview of some of the monitoring tools and their capabilities that have been implemented so far.

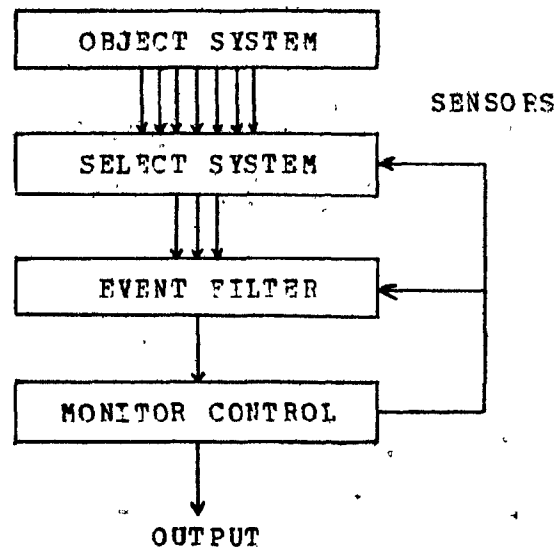


Figure 2.1: General model of a system monitor

The basic structure of a monitoring facility can be

described by the model in figure 2.1, regardless of hardware or software being used to extract the information. The computer system which is to be measured is called the object or host system. Sensors collect signals which indicate changes in the state of this system. These changes are also called events. The sensors can be implemented by software or hardware. Their implementation will be discussed in more detail in the following chapters. A hardware sensor for example can indicate if the wait light on the computer console is on or off. On the other hand, a software sensor can give an interrupt to the monitor whenever execution of a program passes over a predefined point in the code. The select system chooses desired observation points at times when their states are of interest. The measurements then proceed through the next stage where the flow of events is filtered so as to reject signals that are irrelevant for the particular monitoring session. Both selecting systems are supervised by the monitor control which receives the preprocessed signals and produces some sort of output.

Early monitoring devices in general recorded the collected data on magnetic tape. Programs were written to process these tapes after a monitoring session, reducing the large amount of data and creating statistical reports. The problem of interpreting monitor results is considered in [31]. Later monitors reduce and process the incoming data in real time besides producing a detailed record on mass storage. Graphic displays provide a convenient way of visualizing immediately system activities as they take place.

2.1 Objectives for Monitoring

Monitoring is used to obtain exact information about the various activities in a computing system while it is executing some tasks. One of these tasks is the operating system itself which normally is the most complex program of a given system. In general computer manufacturers offer a range of operating systems for different applications. However the detailed specifications which depend on the individual computing environment have to be entered in the form of parameters to the system. Many of these values, e.g. the maximum number of connected terminals or the number of tape and disk drives can be determined right at the beginning. However for tuning the operating system to achieve optimal execution a system monitor is almost indispensable.

2.1.1 Measurement of System Activities

Before optimizing system or user software we have to get a picture of the real system activities during execution. Measurements can give a utilization report for the different modules of an operating system. Based on this information, more frequently used modules can be made core resident while others can be stored on disk or drum. The statistical analysis of a disk usage report could show a high probability that module B is accessed immediately after module A on the same device. In this case seek times can be reduced by locating both files in the sequence A,B close together on the same device.

Many times device bottlenecks have slowed down a whole computer system. It can be very painful to pinpoint the source of troubles if no monitoring facilities are available. Using a monitor, an unexpected overload of an input/output channel can easily be detected just by sampling the channel activities during peak periods. This can also lead to predicting future facility requirements of the computing centre.

System measurement may also help to improve the coding of software. Cantrell and Ellison [19] talk about "performance bugs" by which they mean errors in evaluation or judgement of performance optimization. If the performance of some software is important it should be performance debugged by measurement and analysis. Inefficient parts of the program can be detected and recoded.

Another way to use monitors which have already been proved to be useful is in connection with simulation to produce what is called "trace driven modelling" [15]. A very detailed profile of the job stream, down to an almost microscopic level is observed and supplied as input to the simulation model.

Most operating systems are documented by numerous manuals and even if their descriptions are accurate (which is not always the case), system understanding is better facilitated by monitoring devices that show what really happens in the computer. A performance monitor with graphical output can help teach different aspects of operating systems. Such a device could also be used as feedback to the operators enabling them to make decisions.

The main objective of monitoring in this thesis is to offer a tool for operating system research and development. The monitor can be used to verify the efficiency of algorithms for paging or task scheduling. For the MULTICS operating system [20] this application has proved to be very useful [18]. Performance measurement can also be used for general software research such as program correctness proofs, automatic flowcharting, tracing, and debugging (see for example the SNUPER COMPUTER project [21]). Some of these possibilities are discussed further in chapter 6.

2.1.2 Prevention, Detection and Recovery from Failures

As mentioned in the last chapter, availability is an important factor in computer system performance. The conventional way to increase availability is to schedule certain maintenance periods during which the system is not accessible to users.

Permanent monitoring of the activity of system devices, including the CPU can help reduce maintenance time and yet insure high reliability of the system. Most operating systems are already able to capture fault conditions and produce a corresponding message. Monitoring facilities can go a step further, producing a precise report on all error conditions to show the maintenance personnel where to concentrate their efforts. Further improvements are obtained when the monitor analyzes fault conditions and executes some of the diagnostic routines on the devices that reported failures. When this is done during idle time of the system, the time for maintenance can be reduced even.

further.

(In every computing system there are some resources which are vital to its operations. If the disk drive which contains operating system modules, part of main memory or even the CPU itself is working improperly then the system is "down". At installations that control important processes such as air traffic control or centrals for electric power switching, backup equipment for the most important resources is indispensable. In these systems a monitor could quickly detect an error condition and initiate the necessary actions to replace the defective device by its backup. In computer networks a node causing problems can be located and messages can automatically be routed through another path avoiding this node.

With the growing complexity of operating systems and multiprocessor hardware the reliability of computer installations becomes an increasing problem. Many times users are frustrated by system crashes which are due to implementation errors. For the system manager it is often a very difficult task to reconstruct and analyze the reasons for the malfunction. Dynamic verification and consistency checks by a measurement device could help analyze and even prevent such events. Plausibility checks can, for example, detect unreasonable overusage of resources which may be a hint for failure in the near future. This situation is typical for deadlocks which occur when the resource requirements of different tasks cannot be resolved. One approach for dynamic verification of operating system decisions is described by Fabry in [22].

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2.2 Software Monitors

The earliest monitoring techniques developed used programs in core to monitor what the system was doing. These monitors are an integral part of the host system, i.e. the system that is to be analyzed. Thus, the monitor is both logically and physically a part of this system. The primary attribute of an internal software measurement technique is that it has access to and can selectively record system data stored in memory or CPU registers (depending on the computer architecture, device registers and buffers can be examined too).

When activated, the monitor interferes with the system operation to record and analyze values of interest. This leads to some considerations for the development of software monitors:

- The monitor must not alter any system variables except the ones that belong to the monitor itself.
- Resources used by the monitor, e.g. CPU time, memory space, I/O channels have to be well defined.
- The presence of the monitor must be logically invisible to the system.

Normally a software monitor is written in assembly language. The reasons for this are twofold: first, the execution time can be kept at a minimum and second, the access routines to system registers and memory locations have to be written in a low level language.

A survey of existing software monitors is given in appendix A, table A.1.

2.2.1 Sampling Technique

The basic idea of the sampling technique is to interrupt the normal flow of execution at periodic intervals where control of the system is passed to the measurement routine. The intervals can be determined by a clock or a combination of events in the system that meet some criterion. The best choice for the sampling rate would be a random time schedule which is known to be statistically independent of any natural execution pattern in the program to be analyzed. This avoids synchronization effects between the monitor and the sampled events which otherwise could amplify or zero out the recording of periodic events.

To obtain a compute time profile of a program, that program is loaded and run in its normal fashion without any modifications and without any need for prior knowledge of any of its characteristics. At the time of a sample, execution is interrupted and values of interest are recorded by the monitor. Some of the typical information gathered at this point is given in the following list:

- CPU active
- Location of next instruction to be executed
- Instruction code
- Location of data referenced by the instruction
- Processor status word (PSW)
- Status of channels, control units, and devices
- I/O originators.

This information, recorded on tape along with the sample time is sometimes called a "system's snapshot". At the

(completion of the run this tape can, for example, be sorted by interrupt location address and the resulting frequency distribution printed. The accuracy of the sampling technique is based on the central limit theorem [23]. The central limit theorem states that the accuracy achieved improves as the number of random samples increases. If 8000 samples are taken, the probability that we will be within 2% of the actual value is 99.99%. After approximately 26000 samples, there is no substantial improvement in accuracy.

2.2.2 Intercept Technique

Another approach to software monitoring is to place intercept points (or hooks) at strategic locations in system software and application programs. When the monitor is enabled and execution passes over such a point, control is passed to the measurement program which records interesting values just as in the sampling mode before passing control back to normal execution. Contrary to the method described in the last chapter the intercept technique is very operating system and machine dependent. This is mainly reflected in the way how the hooks can be placed. In systems with many different vectored interrupts where each of them is responsible for a small group of peripheral devices, the interrupts can be intercepted to obtain information about I/O activity. If additionally the operating system and user programs exchange information by means of some sort of supervisor calls which cause interrupts, there is a convenient way of monitoring this type of activity with the intercept method. In some systems

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an entire class of interrupts may be matched into a unique address. Generally within the system architecture there is further information which identifies the specific nature of the interrupt.

To create intercept points for events other than interrupts, the source code has to be modified by insertion of special intercept code at interesting points of the program. This is for example used to monitor if certain locations such as branch entry points or subroutines are executed. Details of the implementation of sampling and intercept monitors are described in chapter 4.

With the intercept technique the requirement for software monitors to be logically invisible to the system when the monitor is not active is more difficult to meet. While there are no problems to deactivate hooks which capture interrupts just by resetting the original addresses in the interrupt vectors, this can cause more problems for hooks that were placed into the code. As changing the source code to remove those points is impractical, a way has to be found to change the hook command into a NOP (no operation) command. In bigger systems like the IPM 370 series, this can be done by the EXECUTE instruction which has the ability to define a hook instruction either to be an interrupt creating one or to be just a NOP. Microprogrammable machines can support the same feature in a similar way just by redefining a micro instruction. For computers from Digital Equipment (PDP11 family) the problem has to be solved in a different way which will be shown in chapter 4.

Another technique which is for example used at the McGill

computing centre should be mentioned. The software monitor (SLACMON, [51]) is invoked whenever the CPU is idle in order to use up all "spare" time. During those periods, a counter is incremented at a fixed rate. The counter value gives then a measure of the CPU load.

2.2.3 Comparison of Sampling and Intercept Methods

In a commercial computing environment where only the overall system behaviour has to be controlled, the sampling monitor would be a sufficient tool. It is relatively easy to handle and the degree of detail can be adjusted simply by changing the sample frequency. On a CPU with an average instruction time of one micro second and a sample rate of one milli second, statistically relevant information can be obtained after only 10 seconds with a maximum of 1000 instructions between two successive samples. In practice this means that a system profile can be updated in 10 second intervals. However the choice of the sampling frequency is not a trivial one and depends upon the yielded information. Another problem with sampling monitors is their low degree of flexibility. The set of values collected per sample during one session is mostly invariable. The selection will have to be general enough to give all the required information but it has to be detailed enough to minimize redundancy in this information. The number of values recorded per sample cannot be too large to keep the ratio of recording time to sampling period reasonably low and thus limit the overhead.

For applications where specific problems in a system have

to be analyzed and where each occurrence of certain events has to be recorded, the intercept monitor is the more appropriate tool. At initialisation a set of events can be defined along with an individual description of the information to be recorded for each of the events. This feature increases the flexibility of that monitor type and decreases the amount of redundant data recorded.

Assuming we want to monitor disk arm movement to find an optimal arrangement of files on that disk we can specify that only values like physical disk address, seek time and originator identification be recorded. If, in the same run, paging activity is to be monitored, the events indicating page faults etc. can produce a completely different set of values. This distinction cannot be made in sampling monitors where we would have to record the union of both sets for every event which of course increases the impact that the monitor imposes on the system.

In general, the main advantage of the intercept technique is the ability to concentrate on a clearly defined range of system activities and to generate a reproducible report.

The monitor overhead in the sampling technique is readily defined by the sampling rate and the number of values recorded per sample. This is not so easy in the case of the intercept technique. Actually the overhead is a function of the event occurrence. Additionally, an exhaustive overall system profile based on a statistical evaluation cannot be obtained.

As long as events can refer to interrupts, the installation of an intercept monitor is no problem: the corresponding

interrupt vectors just have to be modified to point to the monitor routine. But the installation, activation, and deactivation of hooks in the software code can cause some headaches. As mentioned in chapter 2.2.2 this depends on the system layout. A quite unproblematic solution for computers in the PDP 11 family is shown in chapter 4.

All software monitors require host system resources such as CPU time, core memory, buffer space, I/O channel time, and mass storage. This of course is a disadvantage as the monitor's operation may distort the normal system activities and thus may lead to a situation which is known as the Heissenberg uncertainty. Therefore the impact of the monitor imposed on the measured system has to be investigated carefully before the results of a monitoring session can be analyzed.

2.3 Hardware Monitors

2.3.1 Architecture of Hardware Monitors

In the previous section we have seen that software monitors use resources of the host system. Among other reasons this was one of the drawbacks that led to the development of hardware measurement techniques. Table A.2 shows some of the hardware monitors developed during the last 18 years. A good description of the history of hardware measurement techniques is given in [24].

Hardware monitors collect electrical signals that indicate the state of the host system. Probes are attached to test

points or pins on the back panel of the computer system. These high impedance probes are completely passive to the host system. In other words, they do not affect the performance of the computer in any way. Hardware monitors usually operate in sampling mode rather than in the event driven fashion. In the second case an event is caused by the occurrence of an electrical signal at a probe point.

Once signals are presented to the monitor, they are generally made available at some form of logic plugboard. This is used to accomplish first level data reduction on the signals to determine states of coexistence and mutual exclusiveness. If, for example, the overlap of CPU and I/O channel activity has to be measured, we would "OR" the active signals of all channels and combine this output via a logical "AND" with the CPU active signal. More examples for the use of logical data reduction in hardware monitors are listed in [25].

The resultant signals are directed to several counters or recorded together with a time stamp on a storage media. Some of the more recent hardware monitors perform some real time evaluation of the incoming signals and update the result periodically on a display unit.

The method by which the signals are evaluated and recorded constitutes the major difference between early monitors and those available today. It also defines the type of measurement which the monitor can perform and thus is a means of classifying hardware monitors into major subgroups.

2.3.2 Summary Devices

The earliest hardware monitors developed by IBM were summary type monitors (see appendix A, table A.2). They recorded signals either by timing their duration or counting the number of times they occurred. The counter contents was periodically read out and written on tape. Thus, the summary type monitor could give information about the percentage of time a signal was active over the monitoring period. When counting the signals it gave the frequency of occurrence for the measured events.

It is important to note that, any dynamic properties of the system activities were lost in the accumulated information. However for computing systems in the early 1960's this monitor type was good enough to give a clue to what was happening in the computer. At many installations this technique is still used to produce overall system utilization reports.

2.3.3 Dynamic Monitors

With the increasing complexity of computer systems the need for more detailed activity reports became imperative. Information was required for input to various simulation programs whose accuracy depended on the actual distribution of I/O and computing activities. The summary type monitor would indicate the CPU being much more utilized than expected but the cause of the inefficiency remained unknown. One of the early attempts to overcome this problem by using a dynamic monitor is given in [26]. A dynamic trace of the

event "Jump instruction executed" was used to follow the execution flow and to detect loops.

Instead of merely accumulating the events, dynamic monitors provide a trace in time of events as they occur. The major problem in implementing these devices is that the computer is generally able to change its state faster than the monitor can record it. There are mainly two solutions to the problem: either the monitor's resolution is kept at a frequency at which the storage facilities are capable of registering the events or as in [27] the monitor is given the ability to stop the host system at certain times to follow up with the recording activity. However, the interruption of the host computer is against the principle of noninterference as stated above for the basic hardware monitors. Furthermore it can hardly be tolerated in commercial installations. Different buffering techniques have been implemented to increase the resolution of dynamic monitors without losing too much information.

2.3.4 Modes of Operation

Modern hardware monitors combine dynamic as well as summarizing features to provide information about different aspects of the host system activities. Corresponding to the measurements made by the early monitors there are basically three modes of operation:

1. Event mapping (E-MAP) mode: In this mode the counters used by the early summary monitors are replaced by words in the monitor's random access memory. The address of a memory word is given by the code of the event. The

contents of each memory cell represents either the accumulated count of the events or the total time a particular event is active. In the first case, the content of the according memory cell is incremented by one every time the event occurs. In the second case, the monitored signals are sampled at regular intervals as specified by an internal clock.

2. Time mapping (T-MAP) mode: While an input signal is active, a counter is incremented at a rate determined by an internal clock. When the specified signal becomes inactive a memory location whose address is given by the counter value is incremented by one. The monitor's memory then represents a frequency distribution for the duration of the observed signal. In practise this could be used to obtain the seek time distribution of a disk unit.

3. Store mode (S-MODE): The store mode corresponds to the trace measurements described in chapter 2.3.3. The event codes are simply written in consecutive memory locations as they occur or at the frequency determined by the internal sampling clock. The memory is divided into buffer zones which are periodically dumped on tape to free the storage space. The choice of one of the monitor modes depends on the information requirements. Several examples how these features of hardware monitors can be used are presented by Epstein [28].

2.4 Hybrid Monitors

While hardware monitors can virtually measure any significant physical machine event, they can capture logical events only in some cases. Protect key activity on IBM 360 systems for example is strictly a logical function of the operating system, but the register that contains the code for the protect key of the controlling task is in a fixed machine location and can therefore be monitored.

A number of extremely important system characteristics are never reflected directly in the hardware. Thus there is no efficient way for a hardware monitor to determine how much memory is occupied by user programs, how many jobs are currently in the job queue, the reasons why many jobs may wait for resource allocation before being initiated or how long a particular executive routine may reside in core.

The most evident advantage of hardware monitors is that they do not affect the host system in its performance. They are capable of sensing a wide variety of hardware and even some software events but they are limited in their ability to detect the stimulus for a set of responses. This cause and effect relationship can be better established by software monitors which take advantage of the fact that they reside within the host system. They have access to virtually any parameter which is controlled by the operating system software. Especially in multiprogramming systems, the latter technique can be used to investigate the system's task table and correlate the location of an event with the identification of the originating program. Physical machine

events, however, are beyond the reach of software monitors. In table A.3 three monitoring techniques are compared with regard to the most important criteria.

Fortunately, the capabilities of hardware and software monitors complement each other and thus suggest a solution to this problem. The resulting type of systems monitor is called a hybrid monitor. The underlying premise of the hybrid monitor is that a hardware monitoring device is not invisible to the operating system, but instead, it is treated as an "intelligent" peripheral device.

This concept has been implemented at the MIT Lincoln Laboratory [29] where sensors and event counters are allocated to user programs under the control of software. The user can define which activities are to be sensed and the hardware device interrupts the user program as soon as it recognizes one of these events. The device has been used to derive an online histogram of the subroutine utilization, to investigate virtual memory performance of a single program and other similar tasks.

The distinguishing characteristic of hybrid monitors is that software can be used to determine activities which are to be measured by the hardware device. This implies that the commands to a hardware monitor are no longer given by buttons on the front pannel but they are dynamically supplied and altered by a software program. In table 3 the different monitoring techniques are compared.

For example, suppose the memory activity of a certain program is to be monitored. Soft hooks can be established in the supervisor's dispatching module to detect when the

(program is about to be executed. At this point the software monitor can take control and transmit to the hardware the base address and size of the program. Since memory activity is to be monitored it also directs the hardware to select the memory address register as the source for the measurements.

One of the first hybrid monitors for systems with virtual memory is the DSP-1 [30] which is currently used at the McGill computer centre. It reduces the monitor overhead in the above mentioned example in the following way: A map of memory areas that have to be monitored is established by software and sent to the device which then continuously samples the content of the program counter (PC) and compares this value to the ones in the map. As soon as the PC enters one of the specified memory areas, the monitor starts to record predefined system variables.

As we have seen, the monitoring devices have become more sophisticated with increasing complexity of computer systems. While the first monitors purely recorded everything they were able to measure, the more advanced techniques concentrate on the problem of screening the sensed signals to eliminate redundant or unwanted information prior to the recording. Plugboard logic and event filters were improved but at the end it became obvious that many problems could be solved if a second processor were used to perform the preliminary data reduction. This has led to the last generation of performance measurement tools which will be the subject of discussion in the

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following chapters.

CHAPTER 3

Monitoring with External Processors3.1 A General Instrumentation System

The DSP-1 mentioned in the last chapter is an elaborate tool for performance measurements. The sensory system which had to be programmed for different logical combinations on bulky plugboards in earlier systems is now controlled by commands on a console. However, the most important feature of the DSP-1 is the possibility to transmit those commands in remote mode from the host system, i.e. software residing in the monitored computer can define the preliminary data reduction and other parameters that control the monitor's operation.

As it has been shown, hybrid monitors combine the advantages of hard and software monitors but they are not flexible enough to change the monitoring mode dynamically as a function of the incoming events or to perform more complex real time data reduction. Due to the storage requirements of the map and store modes, memories were added to monitor devices. Their internal control structure required an increasing amount of hardware to supervise the data flow. The DSP-1 monitor incorporates already a special purpose processor to take care of these functions.

In this chapter one further step will be taken and a more flexible instrumentation system will be discussed which includes its own processor and software residing in the

monitor's memory. This model is, at the same time, the basis on which a performance monitor was developed and which will be the subject of the remaining chapters.

3.1.1 The General Model

Logically the monitor consists of three parts as shown in figure 3.1. The software monitor, operating in sampling or event driven mode, resides within the host system where it extracts information which is available in the host memory. The hardware monitor collects electrical signals with its sensory system and performs first level data reduction. Both units send the preprocessed information to the control or supervisor system. On the other hand, the control system has the ability to send data back to either of the monitors, interrupt their operation via control lines and also interrupt the host system itself. This last point is usually not found in commercial monitor devices. However, it has proven to be very useful for applications in research and development. Finally, after having processed the data, the control unit communicates the result over the I/O bus to some peripherals.

In the following three paragraphs we will take a closer look at each of the components. A functional drawing of the configuration is given in appendix A, figure A.1.

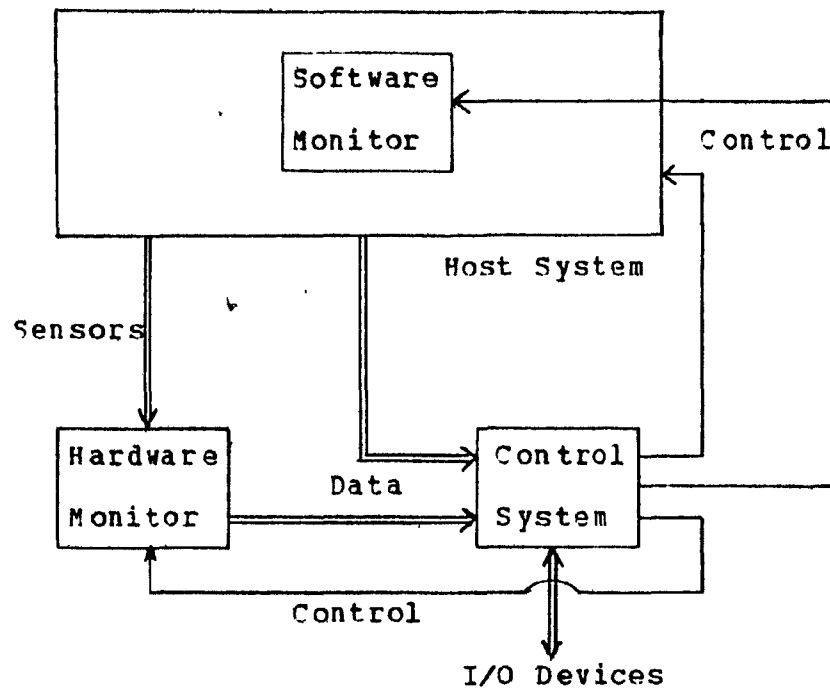


Figure 3.1: General Monitor Model

3.1.2 Software Monitor

This part of the performance measurement system resides in the host memory. For normal applications, the software monitor runs in the intercept mode because the data gathered in sampling mode can be extracted more efficiently by the hardware sensors.

As early as in 1967, as part of the "Snuper Computer" project [21], Estrin developed an idea for hybrid monitors which has been included in this system. Instead of inserting softhooks at strategically important points of programs which are to be monitored, the software is analyzed at compile time and the relative program counter value for those points is determined. The technique used by Estrin is

very similar to those used for program correctness proofs [35] (see also Chapter 5). At load time the absolute PC values of the event locations are determined and passed on to the sensory system which reports an event as soon as one of those locations appears in the PC during execution.

Of course, the method is not easy to implement as it requires changes in the operating system software, namely the compilers and loaders. It has the advantage that the system degradation is placed into utility tasks and the program to be observed can run without any interference. It is also a typical example for a situation where the distinction between operating system and software monitor begins to disappear.

When the monitor, running in intercept mode, has recognized an event it will collect relevant data and send it together with the event identification to the control CPU. This transmission can be done over a standard I/O channel of the host system. The same channel is used in the other direction to send information about the operating mode to the monitor. This has to include interrupt vector locations, a breakdown of events to be captured per interrupt if the vectors are used for multiple functions, and details about the required data per event.

3.1.3 The Sensory System

High impedance probes capture signals that indicate the state of the host CPU and its devices. In practice, a set of probe points is amplified in so called concentrators before the signals arrive at the sensory system. More about

hardware probe points can be found in the literature [32,33]. The signals then pass through logic circuits where the first level data reduction takes place. The combinatorial logic can take many forms. Generally speaking some logical relationship between the signals is determined. Instead of using a plugboard to specify which signals have to be compared and which logical function is to be used, the combinations are defined by the control processor. Additionally this remote technique makes quick changes during a monitoring session possible. The resulting signals are then presented at the input of buffer registers, counters or comparators. In the first case the data is ready to be transmitted to the control CPU or directly stored into memory by means of a DMA channel. Otherwise signals are either counted directly or used to open a gate to increment the counter with a fixed frequency during the signal duration. These two techniques correspond to the STORE or E-MAP mode and the T-MAP mode respectively.

Although the control CPU maintains its own time base it can be useful to send the host time to the monitor. Provided both systems increment time at the same rate, only an initial synchronization is necessary.

Apart from checking the input signals, the use of comparators leads to a new data reduction technique called sequencing. The first monitored value is compared to the contents of a series of registers. The index of the first matching register is remembered. In case there is no match at all, the value is ignored, otherwise the next value is compared with the register following the remembered one and

if the test is successful again, the index of the currently matching register is stored. This can be carried out a certain number of times depending on the size of the sequencer or infinitely if the comparator series is logically closed to form a loop. In a linear arrangement the result of this procedure indicates if a predefined sequence of values occurred in the host system or not. In the case of a closed sequence, a counter can keep track of how many times a loop was executed.

This idea can be carried further to analyze more complex sequences with alternative branches by arranging the comparator registers in the form of a finite state automaton. For any state there is a predefined set of events which can move the automaton in one of the successor states. If the received signal does not belong to the set or if the automaton enters a trap state the sequence of signals is rejected as not belonging to the language (in terms of the theory of formal languages) which was defined by the automaton.

Finally the circuit can have a "learning" ability if the first sequence of n values defines the comparators and is checked against the next n values. In the case of a match, a loop is detected and the comparator values can be transmitted to the control processor to identify the loop or else the second n values are used to redefine the comparator values.

However these attractive features require complex hardware and for economical reasons their implementation will have to wait for further innovations in hardware technology.

3.1.4 The Control Processor

The choice of the control CPU is no easy task. The CPU should have the following properties:

1. Execution speed similar to the host system
2. Fast I/O channels
3. A minimum of 32K words of memory

In general the control system is of a magnitude smaller than the one to be measured. This means that, for example, an IBM 360 installation should be monitored by a mini computer or if a mini computer is to be supervised, it would require a micro computer in the control system. This, however, creates difficulties regarding the first point stated above. In most cases the speed of a CPU decreases with the size. The reason for this lies in the architecture of the central processor control. Large scale systems have many functional units working in parallel. Medium and especially small scale CPU's have a sequential microprogrammed controller which makes them easier to implement but slower in execution. This problem is inherent in all processor controlled monitors. If the control processor is not especially designed for this application, the solution can only be found in the sensory system itself and in alternating or rotating data buffers.

The main bottleneck of monitors lies in the transmission of data between the sensory system and the control processor. This is why most hardware monitors perform some preliminary data reduction with combinatorial logic. This measure is still insufficient for systems which produce events at an

excessive rate. For this reason it is desirable to have a set of high speed registers in the sensory system which can be used as buffers when data bursts arrive at the input. If this is still insufficient there is only one way out of it, namely to interrupt the host system until the monitor has caught up. The interrupt should be directed to the software monitor which can inhibit any processor activity until it gets the signal to release the CPU over the control/host communication channel. At installations where this cannot be tolerated, at least a data overrun indication is recorded along with the data.

A second speed problem can occur at the output channel when data is recorded on mass storage. However, the situation becomes only serious when the monitor is running in store mode where every event is recorded without any data reducing transformation. The monitor speed would then be limited by the speed of the recording device. In all other modes this problem is not very likely to occur as first and second level data reduction as well as buffering techniques should decrease the data rate so that it can be handled easily.

The memory requirement stated in the third point is due to several considerations. First, memory has to hold the decision and evaluation software. Second, it has to provide enough space for buffer zones and third, it has to have room for the E-Map, T-MAP or STORE modes. Some of the information that reaches the monitor does not have to be processed immediately. In this case the operation speed can be accelerated when the processor and its I/O channels are bypassed and the data is stored via direct memory access.

(DMA).

(Apart from managing the data flow from the monitor units the control processor has several other tasks. It has to perform second level data reduction and on-line evaluation and present the results mostly in statistical form to the user. CRT displays or line printers can be used as output devices. At the same time a more detailed record of system events has to be buffered in the control memory to be dumped periodically on a mass storage device, e.g. a magnetic tape.

Finally the processor has to accept commands from a keyboard. This enables the user to set the monitoring mode, the display mode, and other parameters to specify the operations of the whole performance measurement system. A desirable feature would be to have not only a manufacturer supplied set of software routines but a "monitoring operating system" which incorporates the basic tasks of data acquisition. In the CPM-X monitor [34], the most frequently used routines for data reduction are programmed in firmware, which again helps to increase the processing speed of the CPU. The user should have the possibility to write his own evaluation software on top of the utility routines. A considerable amount of flexibility would be added to the system and programs to track down special problems or to modify the evaluation techniques could be written according to individual requirements.

3.2 Application of the System

The interactive nature of this system offers a range of new applications. The most difficult part of a monitoring session is to set up the probe points, define the logic combinations in the sensory system and to specify the events to be intercepted by the software monitor. Finally the secondary data reduction procedure in the control processor has to be chosen. Experience has shown that the first setup never exactly gives the desired information. Several test and correction cycles have to be carried out before the result is satisfactory. With the monitoring system developed in this chapter, the adjustments can be done interactively from the monitor's keyboard.

Let us consider a typical example. During system profile measurements a significant operating system overhead is observed in one of the utility programs. Without relocating probe points or inserting new hooks in the software, the monitor can supply the identification of the troublesome routine as soon as it is loaded from disk to core by associating the physical disk address with the name of the accessed module, assuming that the system device directory is available for the monitor. Knowing the size of the identified routine a sampling run can be set up to observe PC activity for the corresponding core segment. The start of the sampling session can be defined as the event of loading the routine the next time into core. The resulting execution profile will then show which parts of the routine are responsible for the low performance. Commands for this

(tracing procedure can be given in a simple language (possibly similar to BASIC) and the complete process can be supervised from the monitor's console.

The centralized control over the measurement system also enables non specialists to handle the monitor. Operators can ask for continuous feedback for their decisions at the host system console and can immediately see if the system load is optimal. The real time display showing the different system activities is a valuable tool for teaching purposes. The pedagogical value of a dynamic display is obviously better than static explanations on the blackboard.

Finally the interactive monitor is an efficient tool for system development and research. In the experimental computing environment the host system can be interrupted or slowed down. While events have to be otherwise captured "on the fly" and the results are displayed in summary form or evaluated offline, in this application a detailed online analysis is possible. The monitor can be regarded as a sophisticated debugging routine. The hardware probes allow an instruction trace much faster than any conventional trace routine. If the monitor has the ability to alter the state of the host system, e.g. changing the PC, stack pointers etc., all the requirements for a powerful debugging tool are met.

CHAPTER 4

HIMOS - Hybrid Interactive Monitor System4.1 Purpose of HIMOS

One of the initial objectives of the design of a monitoring system for the mini computer installation in the School of Computer Science at McGill University was to provide a method to increase the availability of the computers. It was planned to develop a technique for automatic loading and execution of diagnostic programs as soon as a system failure was detected. During the analysis of the problem the idea was abandoned for several reasons. First it turned out that monitoring was a prerequisite to accomplish the original goal. Although there was a multitude of literature on computer system instrumentation, the subject seemed promising for further research. Second it was found that problems other than availability, yet related to monitoring, were more pressing. At the same time the project offered the possibility to test the real time capability of the mini computer used to control the monitoring process.

The resulting objective was to design an interactive hybrid monitor, controlled by a dedicated processor as outlined in figure A.1. Because the monitored system was underutilized most of the time, the improvement of performance characteristics was not one of the prime considerations although with the increasing workload, this may become an

important point in the future. But, as mentioned before, monitors can be used in many other applications. During the last year mainly three directions emerged for research on the mini computer installation:

1. Multiprocessor systems: This subject is specially attractive as there are already two mainframes available in the department. In the long run it is planned to extend the system by several micro computers to build a network for research purposes which could be supervised by the monitor.
2. Operating systems: Apart from the system software required for point one, there are projects on paging algorithms and task scheduling routines in progress. Recently two new time sharing systems (UNIX [36] and concurrent PASCAL [37]) were acquired and the monitor should be of help in understanding and extending those systems which are not very well documented.
3. Programming languages: This field offers a wide range of monitor applications. It is planned that one of the projects uses the monitor to produce flowcharts dynamically according to the real execution. This topic is closely related to the problem of program correctness proofs for which the monitor offers some practical possibilities (see ch. 5).

These three points can only give a rough overview of the possible applications. One side effect of the monitoring system for example showed up during the development of the monitor itself. The currently used operating system, RT-11 (supplied by Digital Equipment [38]), employs several

techniques which were not described in the accompanying manuals. Already the very first version of the monitor detected several unknown scheduling techniques and the knowledge of their existence induced new implementation ideas.

Finally the monitor will assist in teaching several hardware and software courses. For a course in introduction to computer architecture, hardware probes can be connected to the host system and dynamic changes of the machine state can be demonstrated under reduced execution speed on the graphic display. For students in software oriented courses, the real time creation of the execution profile for a sample program will aid in understanding the relation between static and dynamic program structures. The computer will no longer appear as the "mystic black box" and terms like program counter or interrupt will become more realistic once they are visualized.

4.2 Computing Environment

4.2.1 System Configuration

The computing facilities in the School of Computer Science are subject to continuous improvements in soft and hardware. Therefore references to the system configuration are as of March 1976.

The complete computing equipment was manufactured by Digital Equipment Corporation (D.E.C.). There are mainly two independent systems. One is based on a PDP 11/20 processor with 28 ~ K words of core memory, one 2.5 Megabyte

disk drive, two 288 Kilobyte DEC-tape drives and a Tektronix 4012 graphic display which serves as a console device. This installation is used as the control system for HIMOS.

The host system is built around a PDP 11/45 CPU with 40 K words of core memory, a memory management unit, one floating point processor, one 2.5 Megabyte disk drive and two system consoles consisting of a 300 baud hardcopy terminal (DEC-Writer) and a 1200 baud video display (Beehive).

The two systems are physically located close together. They are connected via two communication channels. One consists of two interconnected serial asynchronous interfaces operating at approximately 10 Kilobaud; the second link operates by means of two parallel asynchronous 16 bit interfaces at an average transmission rate of 40,000 words per second. A third parallel interface establishes the communication between the PDP 11/20 and the hardware monitor. All interfaces were originally designed by DEC but were modified to fit the special needs for high reliability and increased data transfer rates.

The complete system configuration is shown in the appendix, figure A.2.

4.2.2 PDP 11 Family

In order to fully understand some of the monitoring techniques applied in HIMOS, it is necessary to have a basic knowledge of the PDP 11 system architecture. This paragraph will only highlight the features which are relevant to monitor implementation, especially with regard to the PDP 11/45 structure as this computer was chosen for the host.

system. Further details about the PDP 11 structure, the corresponding processor handbooks should be consulted ([39] and [40]). A detailed description of peripherals and interfaces can be found in [41].

The PDP 11/45 is a 16 bit parallel processor which can run in kernel, supervisor or user mode. It can perform DMA data transfer as well as CPU controlled interrupt or flag driven I/O. The interrupt system has a 4 level priority scheme and the CPU can be on one of 8 different priority levels. The 16 general purpose registers (GPR's) are functionally divided into two sets of 6 registers each, 3 stackpointers (SP), each for one of the three different operating modes and the program counter (PC). The address space ranges from 0 to 32 K words and can be extended up to 128 K words under memory management. The upper 4 K of address space are reserved for I/O devices and the lower 256 words are used for interrupt vectors. The CPU interprets 78 instructions which can be used in 12 different addressing modes.

One of the most important features of the PDP 11 design is the UNIBUS concept: All system components and peripherals are connected to a single bidirectional asynchronous bus, called the UNIBUS. Since all system elements, including the CPU, communicate with each other in identical fashion via the UNIBUS, no distinction can be made between memory access or input/output operations. Although this technique facilitates programming it caused some problems for the monitor implementation.

The HIMOS intercept technique is based on the interrupt mechanism of PDP 11 computers. Interrupts can be caused by

peripheral devices or the CPU itself. A hardware interrupt occurs when a device wishes to indicate that a specific condition has occurred. The relative priority of the request is determined by the processor's priority and the level at which the request is made. If the device's priority is higher than the current processor level the following actions take place:

1. The processor relinquishes control of the bus passing it over to the requesting device.
2. The device sends a unique memory address as base of the interrupt vector. The first word of the vector contains a pointer to the interrupt service routine, while the second location holds the new processor status word (PS).
3. The CPU stores the current PS and PC in temporary registers.
4. The new PC and PS are taken from the interrupt vector. The old PS and PC are pushed in that order onto the current stack and the service routine is initiated.

Upon the RTI command (return from interrupt) these actions are executed in the reverse order.

The following commands in the PDP 11 instruction set can cause an interrupt sequence:

BPT, IOT (Breakpoint and I/O traps):

The first command is used for debugging techniques where breakpoints can be created by substituting the original instruction by the BPT code. IOT can be used to schedule I/O actions.

EMT, TRAP (Emulator trap and trap):

Those two instructions are used for several dispatching techniques. The high byte contains the operation code while the low byte contains a user defined constant in the range from 0 to 377 (octal). In RT-11 the EMT instruction is used as supervisor call.

PIRQ (Programmed interrupt request):

This instruction offers another way for task dispatching. The request is treated in a similar way as device interrupts because the priority level can be specified.

Finally there are 3 important vector locations reserved for processor fault conditions. As soon as one of the conditions occurs, the processor traps to the corresponding vector.

Trap to 4:

Most novice programmers are confronted with this error condition. It indicates the processor "time out" signal. The reasons can be access to nonexisting memory or device locations and odd address in a word instruction.

Trap to 10:

The processor has encountered an illegal operation code. This situation is typical for attempts to execute data instead of program code.

Trap to 24:

Whenever the power drops beyond a certain limit the processor traps to this location to initialize the power fail routine. When power is restored a trap to

the same location will start the power-up routine.

4.2.3 The RT-11 Operating System

Originally it was intended to monitor the PDP 11/45 under the UNIX time sharing system. However, for technical reasons the UNIX system generation was delayed. At the beginning of 1976 the implementation of HIMOS was started under RT-11 (Version 02B-SJ) which is, at this time, the only available operating system.

RT-11 is a single user operating system for computers of the PDP 11 series. On the PDP 11/45 and 11/70 the CPU is permanently kept in kernel mode while executing RT-11 or user tasks scheduled by the system. It requires a minimum of 8 K words of core and can handle a maximum of 32 K words address space. It is designed for program development and real time applications. It supports PAL-11 and MACRO, the PDP 11 assembler languages, FORTRAN IV and BASIC. Several utility programs are available such as the text editor (EDIT), and on line debugging aid (ODT), a file handler (PIP), several device handlers and a batch utility. It can accept device interrupts and service the two most important error traps (trap to 4 and 10).

In the following some of the aspects of RT-11 which are important for monitoring will be discussed. Further information can be obtained from the manuals [38, 42].

When RT-11 is booted it determines the physically available core (the maximum for RT-11 is 28 K words) and loads itself on top of the memory. The resident system part needs approximately 4 K words. The trap vectors and the most

important device vectors are then initialized. When device handlers are brought into core, RT-11 locates them just below the system. Normally, user programs are loaded at address 1000 (this and following addresses and memory values are given in octal if not stated otherwise). The resulting memory layout is given in figure 4.1.

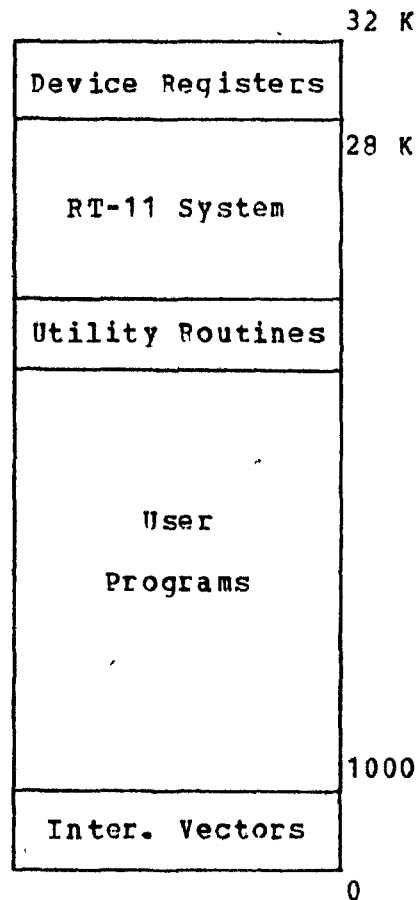


Figure 4.1: Memory Layout under RT-11

An executable program image is called a SAV-module. It always represents core locations from zero to the high limit of the code. This enables a programmer to define certain vector values statically in his program. To make sure that no vital system vectors are erased when a user program is

loaded, the system resets some vectors after a program load. The vectors to be restored are indicated in a protection table within the system. This technique is obviously not the most elegant way of system protection. On the other hand, RT-11 cannot take advantage of a hardware memory protection as this is not a standard item and is only supplied with the memory management unit. Therefore RT-11 can only protect itself at program load time and also against keyboard commands from the console (like Examine or Deposit). This is not the case during run time. A user program has access to any existing address within 32-K. The subject of memory protection has considerably influenced the implementation of HIMOS.

The communication between the operating system and a user program as well as internal RT-11 scheduling is performed via EMT-calls. The use of the EMT instruction is therefore exclusively dedicated to system actions. A large macro library contains all system utility routines, for example, READ and WRITE requests. The macros use EMT-calls to transmit a command via software interrupts to RT-11. For example the EXIT macro call at the dynamic end of a program is translated into an "EMT 350" instruction where the number 350 serves as a function code for RT-11.

Disk and tape storage is blocked into records of 256 (decimal) words. On RK05 cartridge disks this corresponds to one sector. At the beginning of a tape or disk a file directory is located. It contains starting address and lengths of the files on the device. Files are stored sequentially but after several file manipulations there may

be gaps between individual files.

RT-11 has no file protection. Each user can access any information on a device. Only the directory is protected, provided system macros for file manipulation are used.

4.3 Concept of HIMOS

The HIMOS system was specially developed for the system configuration described above. However it can be used to monitor any other computer within the PDP 11 family, provided a second CPU for the control system is available. System dependency is introduced by the software monitor which is designed to operate under RT-11. The hardware monitor part is more flexible although some of its logic has to meet special requirements of the PDP 11/45 architecture.

The PDP 11/45 was chosen to be the host system which is monitored under control of the 11/20. This choice is due to the fact that the bigger 11/45 has more components of interest such as the floating point processor and the memory management unit. Originally it was planned to execute the control software on the recently announced LSI-11 micro computer (manufactured by DEC) but it was not available on time. However, the control program can easily be transferred at any time since the instruction set is compatible.

HIMOS was implemented in three steps. A first version of the software monitor allowed some initial experiments to become familiar with various problems. This version ran like the early software monitors on the host system which was used at the same time to analyze events and to output

the results. Second, the hardware monitor was designed and third the control system was implemented on the 11/20. This approach led to a modular design of the instrumentation system. Each monitor can operate on its own. Only the supervisor CPU is necessary to receive information and to control the operating modes.

4.3.1 Software Monitor Functions

The software monitor (S-MON) is core resident in the host system where it is protected by RT-11. When a specified event occurs it is captured by S-MON which accumulates pertinent information in a buffer before transmitting the complete event record to the control system (CS). The event record is headed by the event code (EC) and can virtually contain an unlimited number of additional words indicating the state of the host system at the time of the event. In a commercial environment S-MON would release the CPU after having gathered the information and would transfer the record concurrently with the continued execution of the interrupted task. In HIMOS however, the record is transmitted and a response from the control system has to be received before S-MON releases the CPU. Although this technique increases the monitor overhead, it is necessary to allow interactive response from the supervisor. If a special event is detected in the control system, the monitor mode may be changed, the hardware monitor initialized etc, before the host system changes its state. This feature opens a wide field of applications for HIMOS which are not possible in most other monitor

implementations. Any hard or software event may be dynamically declared as a special event. The list of these EC's is kept in the control system.

From its design the software monitor uses exclusively the intercept technique. As will be shown later, even the sampling mode is implemented using the same technique. At initialization time the vectors for all interrupts to be recorded are replaced by pointers into the monitor program. The original vector values are saved in S-MON. The choice of the vectors can be interpreted as the first level of data reduction (as a matter of fact it is data elimination). At the same time a list of items has to be specified which are to be recorded upon the occurrence of an event. Every individual event can have its own set containing items of interest. This gives the second level of data reduction.

At this point it is helpful to follow an event through the software monitor (see fig. 4.2). Let us assume "trap to 10" is specified as an event. As soon as the CPU encounters an illegal instruction code, the execution flow is interrupted and control is passed to the location pointed at by memory address 10. As a result the PC will contain the entry point of S-MON. The event code is computed and used in the main dispatcher as an index to call the corresponding event handler (EH). There, the system state, according to a specification list, is stored into a buffer and the complete record is sent to the supervisor system. Assuming that the event analysis did not indicate a special event, the supervisor simply sends an acknowledgement and S-MON will transfer control back to RT-11 as indicated by the original

contents of the interrupt vector.

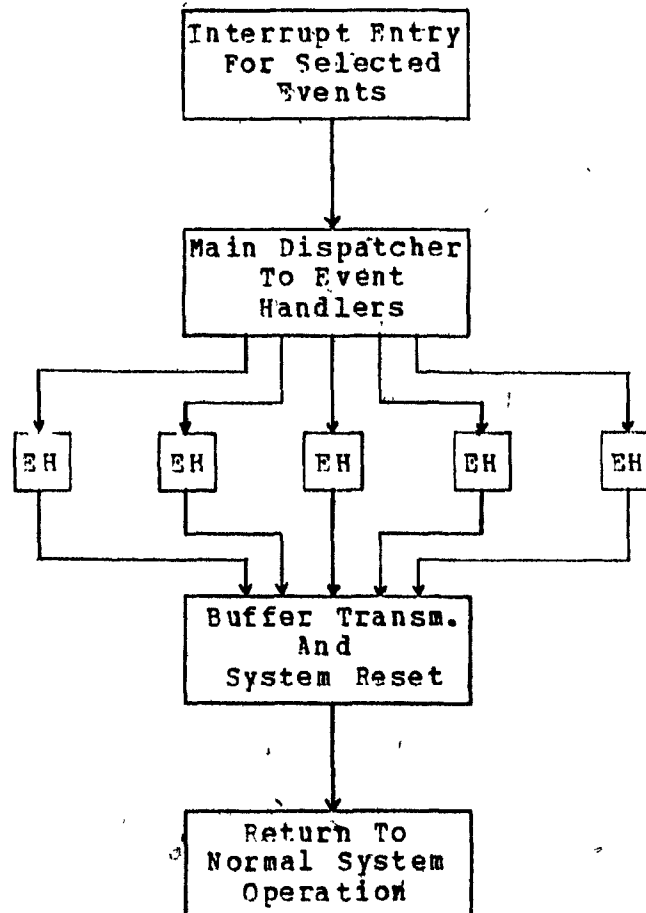


Figure 4.2: Functional Diagram of S-MON

The procedure is completely transparent to the user and to RT-11, as long as the event is acknowledged by the control system. If a special event was detected the control CPU has the possibility to change the system state in the host in which case HIMOS ~~interferes~~ actively with the host performance.

As mentioned before, the sampling mode is handled in a similar way: the 11/45 clock interrupt is intercepted and the event is reported to the supervisor. To transfer

control back to the point where a task was interrupted the necessary information is taken from the system stack. This solution offers the possibility to combine interrupt and sampling techniques during the same monitor session. For example device and CPU activity can be monitored while every second a sample interrupt is created. This can be used to monitor if the system is still in working order or if, for example, the CPU was stopped by a HALT instruction.

4.3.2 Hardware Monitor Design

The design of a hardware device to monitor PDP 11 computers requires some special considerations. Mainly the way how I/O activity is handled causes practical problems. Any I/O operation is performed on the UNIBUS and except for interrupts it cannot be distinguished from a regular memory access. Further difficulties are due to the fact that there are not many meaningful probe points available at the CPU backplane. To monitor, for example, the PS or PC requires complex and dangerous manipulation on the CPU boards. It is desirable that the computer design for next generation systems take that problem into account and bring out sufficient probe points for easy access. However some of the important signals in PDP 11 CPU's are brought to a plug which is used for maintenance purposes. This has simplified access to a subset of important probe points.

A general design of the hardware monitor (H-MON) was given to the Department of Electrical Engineering at McGill where it should have been implemented as part of a course project. Due to problems with the delivery of components the

implementation has not been possible so far. Based on a detailed design together with the engineering drawings, the function of the device will be discussed in a later chapter. Possible applications of the hardware monitor will be considered although results of practical experiments are not available.

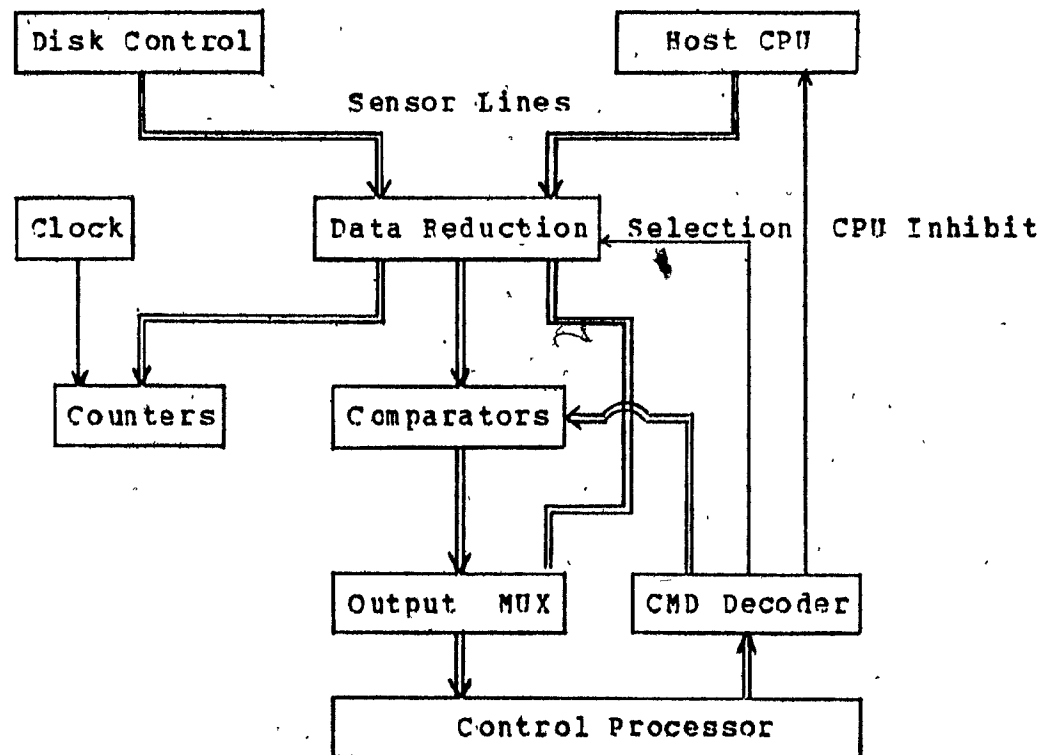


Figure 4.3: Hardware Monitor

According to the initial design, the hardware monitor has several groups of sensors which can be connected to probe points of the host CPU and the disk controller (see fig.4.3). The signals are then supplied to a data reduction logic which is controlled by the supervisor system. Hardware counters can accumulate signals at high speed or they can be used to time events with the internal frequency.

generator. The device operates in different modes which can be changed dynamically by the control system. Results are transmitted over a modified parallel interface for further processing.

4.3.3 Monitor Control and Output

The supervisor system, implemented on the PDP 11/20, has complete control over all monitoring actions. Its command interpreter accepts a large set of instructions which can be grouped into three parts:

- a. Display and recording mode
- b. S-MON control
- c. H-MON control.

(Any references to software features for the control of the hardware monitor imply that the relevant task entries exist in the supervisor system, but the corresponding code has been replaced by a NOP instruction.) To facilitate the definition of a monitoring session, command (CMD) files for each of the three groups can be created for repeated use.

The control system (CS) can interrupt and redefine the operation of S-MON and H-MON at any time either as a function of received data or according to a keyboard command. It analyses event records and stores them by means of an alternating buffering mechanism. Buffers can periodically be dumped on tape or disk. At the same time the CS displays the events in two different ways. For detailed information the complete record for every event is written out in numeric form. As this happens in real time and because the transmission to the terminal operates at

2,400 baud, the host activity has to be slowed down according to the event density and the number of items to be listed per event. In the graphic display mode where the occurrence of an event updates a histogram in real time, the monitor overhead is practically limited by the record transmission between the host and the CS.

Further tasks of the CS are to maintain the system time relative to the monitor initialization and to start and stop the monitor session in a remote way.

It may be instructive to follow up the example of an event occurrence given in section 4.3.1 with regard to the supervisor system (see fig. 4.4). As soon as the event record is received at the S-MON interface it passes through the event filter. The event code (EC) is compared to the list of special events. If the list is empty or if the EC does not match, the filter will send an acknowledgement to S-MON to release the host CPU. The system control will then direct the record to the output handler if real time output is specified. Otherwise the CS operation is finished. A buffer overflow is automatically resolved in the receiver program by switching the buffers and the filled buffer is recorded on mass storage at a later time when the CS is idle. If numerical or graphical display was specified the output handler will perform the necessary data transformation.

In case the event filter finds the EC in its list, the special event sequence is invoked. The display will give a numeric output of the event record (no matter in which mode it was running) and indicate that this is a special event.

During that time the host is suspended and S-MON is in receiving mode, waiting for further instructions, the operator can react in different ways. He may release the host immediately and thus terminate the special event sequence. He also can change the H-MON or S-MON mode according to new requirements for the succeeding analysis, or he decides to terminate the monitor session.

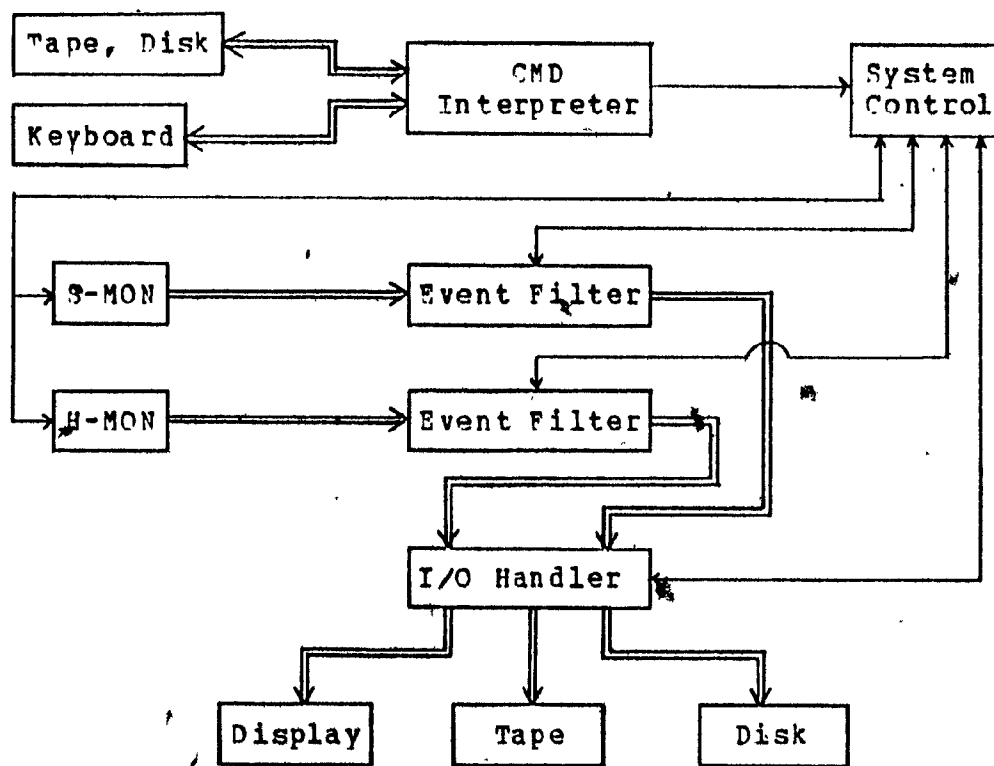


Figure 4.4: Monitor Control System

Other possibilities are left open for future extensions. As the receiving program in S-MON is a flexible, command code driven dispatching routine, other commands to S-MON like "examine location X" or "change value of X to Y" may be added.

4.4 Monitor Implementation

The two chapters on the soft and hardware implementation of HIMOS will give a comprehensive description. Further details on the software can be found in appendix B which contains the complete source listings. Appendix C gives a documentation of the hardware monitor and the communication link.

4.4.1 Software Monitor (S-MON)

S-MON is coded in PDP 11 MACRO assembler and consists of one load module of approximately 2 K words. In order to make S-MON core resident and protect it against user program as well as RT-11 access, the operating system bootstrap had to be modified. The idea is to load the executive not completely on top of the existing memory space but to leave some room for S-MON right above RT-11. In the current configuration RT-11 is located below 24 K. As a result RT-11 assumes that 24 K is the physical limit of available core and will not access higher locations. Furthermore it will abort attempts to modify the top area from the system console. A description of how to change the bootstrap can be found in the software support manual [42].

Another system modification was necessary to preserve the communication channel for the control system. As mentioned earlier, RT-11 loads every executable program starting from physical location 0 and resets afterwards important vector locations between 0 and 1,000. The restoring is performed in accordance with the system protection map. This map is

dynamically changed by S-MON as soon as the monitor is started. The following vectors are protected:

trap vector, loc. 34

clock vector, loc. 104

communication link vector, loc. 450.

User programs are usually terminated by transferring control back to the operating system with the EXIT[™] macro call (EMT 350). If register 0 is cleared at this moment, RT-11 will execute a RESET instruction which disables all interrupts before reenabling the vital interrupts for the system. This however would harm the monitor operations because first, S-MON cannot receive supervisor commands any more and second, if it is in sampling mode, the clock interrupts would be turned off and the timer registers would be reset to 0. For these reasons, S-MON changes the RESET instruction in RT-11 to a NOP. This modification does not distort the performance of RT-11 but it insures full protection for the monitor activities.

4.4.1.1. Event and Data Selection

Software events are defined as interrupts which are intercepted by S-MON. A total of 32 events can be specified which is enough to cover all possible interrupts in the PDP 11/45. At initialization time, S-MON exchanges the content of an interrupt vector with a new PC and PS, and saves the original vector in a monitor table (VECSAV). To identify an interrupt, an event code (EC) is assigned to it. A convenient place for the identification was found in the condition code which resides in the 4 least significant bits

of the PS. As we have 32 possible EC's and only 16 can be specified using the condition code, two different interrupt entries were created in S-MON. This means that the new PC for events in group 0 points to ENTRY0 in S-MON and the PC for group 1 points to ENTRY1 (fig. 4.5). Each group can accept a maximum of 16 different interrupts. The EC is then defined as group index times 16 plus the interrupt index given in PS. This gives a total range of 0 - 37 in octal. Table A.4 in appendix A lists the event codes with the corresponding interrupt and vector locations.

Two mask words (MASK0, MASK1) are used to specify which of the 32 interrupts are to be captured. For example, if bit 3 in MASK1 is set, event 23 will be activated. To define a monitoring session, the mask words can be specified either at the host console when S-MON is started or interactively from the control system. If the command sequence at the beginning of S-MON is skipped, a default configuration will be assumed in which traps to location 4 and 10 (EC 1 and 2) are intercepted.

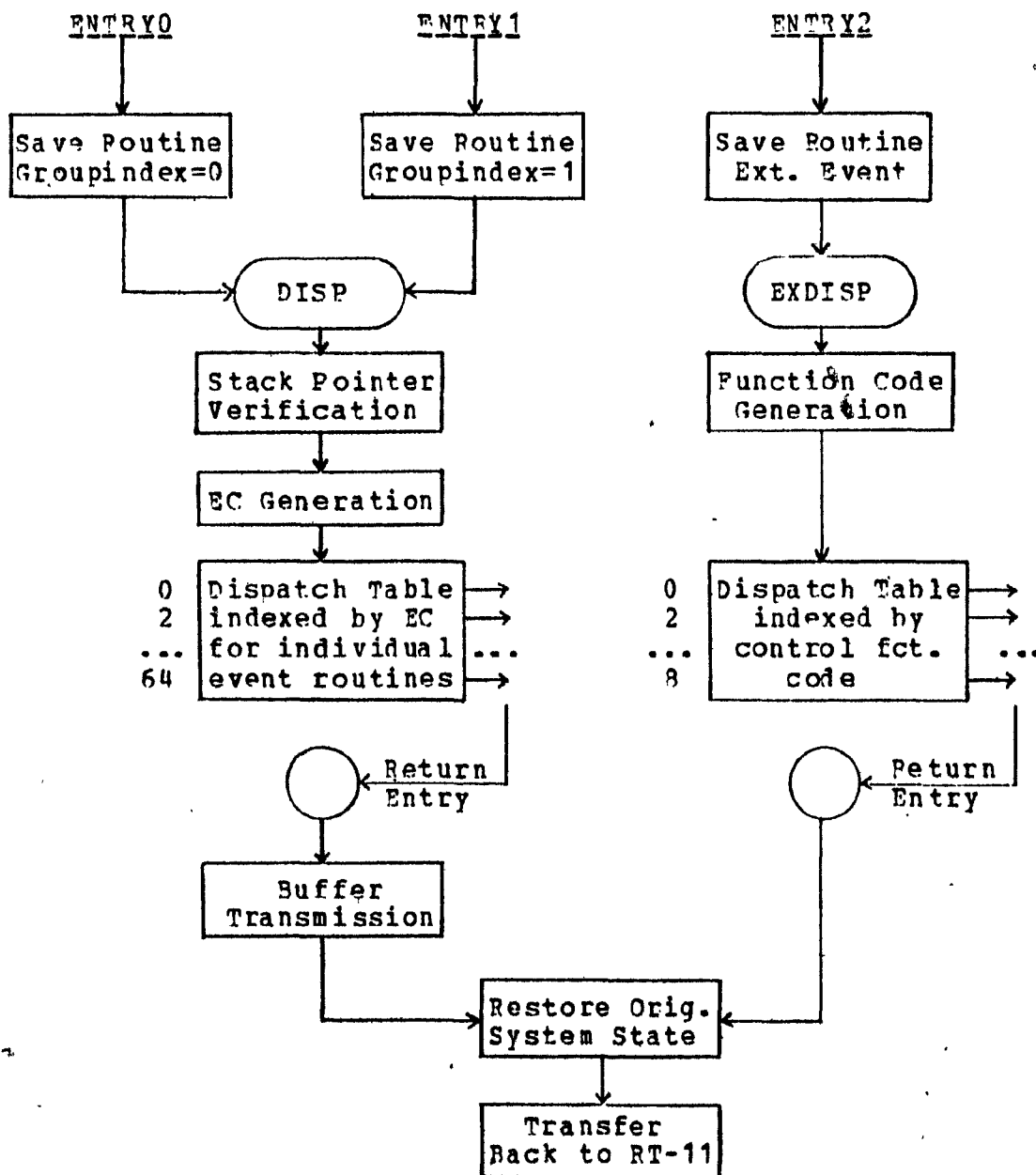


Figure 4.5: Core Resident Part of S-MON.

When an event arrives, the monitor creates a so called event record which contains information about the system state at the time of the interrupt. For every event, different system variables may be of interest and the monitor overhead would be increased considerably if all possible variables were recorded each time. Therefore S-MON

offers the possibility to define individual event records to meet the information requirements for every event. The desired record for each event is defined in the submask (SM) words. The minimum record for all events contains the three most important system parameters, i.e. the PC, PS, and SP. Most events can report additional values such as register content, buffer contents and RT-11 parameters. The exact assignment for the submasks is given in table A.5. For example, SM0 is responsible for the records of event 1 to 5. The low byte covers T4, T10, and the first two values of BPT, the high byte gives the last BPT bit and specifies the records for IOT and power fail. A value is included in the record when the corresponding bit is set.

More information can be reported for the EMT's. As this is a one word instruction [PC-2], the contents of the program counter before it was incremented will be the EMT code along with the function code which serves to distinguish 256 different EMT calls. For example, "104013" shows that an EMT 13 was executed. Because the EMT is used as a supervisor call in RT-11, the error code (EPC) can give valuable information.

The submasks SM2 to SM6 correspond to events in group 1 which are device interrupts. Depending on the device, register and buffer content can be included in the event report. The matching register addresses are contained in the device table of S-MON which is indexed by the EC.

The most important device is the RK-11 disk controller. SM6 is fully dedicated to define the event record for the RK-11 interrupt. It contains the basic processor

information plus the content of 5 disk registers (the register assignment is given in [41]). R3 which originally contained the word count was redefined. This can be justified because normally the word count is 0 when an interrupt occurs. It now contains the block number of the accessed file on the disk which can be obtained by transforming the physical disk address in R5. In the new representation, the disk address can immediately be used to refer to the directory dump where the relevant file name can be found.

Some events like EMT's and disk interrupts occur very frequently and it is desirable to perform preliminary data reduction on the host system before reporting the events to the supervisor. This eliminates unnecessary transmission time and decreases the input load at the control system. The submasks SM1 through SM6 contain one additional bit (bit 0 in each byte) for every event to indicate a preselection. If the bit is set, S-MON will examine one of the values in the event record and compare it to the content of a list of values, the so called select words. If the EC is found the event will be reported, otherwise it is ignored. As shown in table A.5, select words 0 to 3 can hold up to 7 EMT function codes (a function code can be specified in one byte). The end of the table has to be marked by a 0 entry. Select words 4 - 7 are used for TRAP codes. If no trap codes are specified the list can be used for another 7 EMT codes if necessary. The selection for disk events is kept as flexible as possible. Each item in the list specifies first the register number and second the value to be

compared. Thus, the criterion to report a disk event can, for example, be the disk error word or the physical address. TRAP or FMT codes can be extended over select words 11 - 20 if no disk selection is specified. This gives a maximum of 31 FMT codes or 23 TRAP codes which can be listed.

Every event record is headed by the corresponding event code. If specified, the content of the high speed clock buffer can be appended as the last transmitted value. This is important when program segments are to be timed.

All tables in the S-MON data structure are indexed by an even multiple of the EC. In conjunction with the masking technique, this has considerably helped to keep the execution time within reasonable limits.

4.4.1.2 S-MON ROUTINES

When the software monitor is started, it will print an identification on the system console and enter the command interpreter. This routine is responsible for the event and record definition. The execution is controlled by 4 dispatch tables which makes the program fast and easy to understand. The command interpreter (CMDI) is used to define the two mask and the seven submask words and to enter values for the select words. When the index for a mask word is entered, the mask word as it was defined previously is printed and modifications can be specified in the following line. For select words, CMDI distinguishes between SM0-10 and SM11-20. The first group is updated in byte form, i.e. for every word index two new byte values are requested. The second group is word oriented. For every index the two

error checks (e.g. CRC or LRC) could be used.

The speed of the parallel link is only limited by the UNIBUS speed, interrupt latency, and time requirements for the interrupt handler. In HIMOS the transmission rate is approximately 40,000 words/sec. This allows to transmit an average event record of 20 words in 0.5 ms.

HIMOS uses a simple protocol to transmit event records and specification tables over the communication link. The host system terminates an event record by the transmission of EOT and then waits to receive an acknowledgement (ACK) from the control station. When the host tables are updated interactively from the supervisor, the host sends the old table content, terminated by ETX. After the update, the control system issues an enquiry (ENQ) signal before transmitting the new table which is again terminated by ETX. All protocol and function codes are verified in both stations. A wrong function code will dispatch an error routine to generate a corresponding message. An erroneous protocol will cause a retransmission of the last record.

H-MON uses a parallel interface, similar to the ones described above, to transmit data to the control system. For this application, a DMA-channel would provide much better results but the corresponding interfaces were not available.

4.5 Control System Implementation

4.5.1 Program Control Structure

The control system (CS) program consists of 3 object modules: LOG, LOGIO, and LOGGRF (for listings, see appendix B.2). LOG is the main program and the other two modules contain basic I/O routines and subroutines for graphic display on the Tektronix terminal, respectively. To execute the CS program, the three object decks are linked together and loaded at starting address 1000.

After initialization of system parameters and interrupt vectors (S-MON, H-MON, clock), LOG enters the main dispatcher at WAITLP. This is the busy wait state in which the system waits for interrupts and schedules various tasks. During execution of this loop, the program examines the content of a dispatch word (DISPWD) and if one or more bits are set, the corresponding tasks are called (see fig. 4.6). Every time an event record is received, up to 3 bits in the dispatch word can be set in the event record processor (FILTER):

bit 1: record buffer

bit 2: print event record on terminal

bit 3: update graphic output.

The corresponding programs are discussed in chapter 4.5.4. Each of these routines will clear its own dispatch bit before returning control to the wait loop. The S-MON interrupt service routine (RCVREC) sets bit 0 in DISPWD to

schedule FILTER.

The second possibility to leave the loop is given by the reception of a command from the keyboard. This event transfers control to the system command interpreter (SYSCMD) which will analyze the command and schedule the according subroutine. Timer interrupts are directed to the line clock handler (LCLK) where the system time is updated.

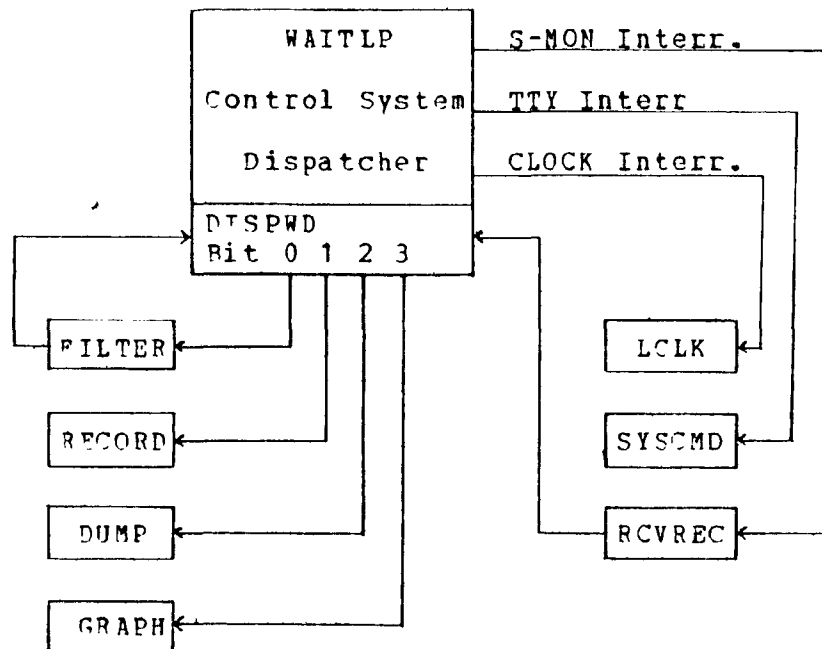


Figure 4.6: Control System Dispatcher

The centralized dispatching technique has helped to keep the program structure modular and easy to follow up. This is especially important in a system where different interrupts can occur at unpredictable times. Further program extensions can be introduced via unused bits in DISPWD without affecting the control structure.

LOG handles three input/output files:

Command File (SPC)

Table File (TAB)

Logging File (LOG).

Default extensions for file names are given in brackets. The command file contains all necessary specifications to define the supervisor system mode. It is created by the system command interpreter which accepts specifications either from the keyboard or from a previously generated command file. The table file is used in a similar way and keeps information regarding the host system mode, i.e. which interrupts are to be intercepted and what information is to be reported for every event. Both files help to simplify the use of HIMOS. Different monitoring sessions can be predefined to standardize frequently used operating modes. To invoke one of those sessions, the operator has only to know a small part of the full instruction set.

The logging file contains the event records as they were received from S-MON. A program (LOGPRT) was written to print this file in a readable form on line printer or data terminals. For performance evaluation which implies a statistical analysis of this file, other programs can be written according to the individual requirements. All HIMOS files can be stored on tape or disk units of the control system.

4.5.2 Command Interpreter and Operating Modes

The control system has three operating modes which specify the form of the event record output. For a detailed online analysis, the records can be displayed in numerical

form at the terminal. The record is headed by the EC and all following values are preceded by a two letter identification, explained in table A.6. Alternatively, the events can be presented in graphical form. A histogram is drawn dynamically in which a maximum of 21 different EC's can be entered along the X-axis. The Y-axis shows the number of occurrences for each event. The initial maximum on the Y-axis is set to 56. Every time one of the columns becomes too big, the picture is rescaled by factor of two. Concurrently with one of the display modes, the system can record incoming events in a logging file on tape or disk. Other display and storing features are explained below along with the corresponding control commands.

The operating mode of the control system can be defined by commands from the keyboard. The command interpreter (SYSCMD) accepts strings of virtually unlimited length but only the first character is significant to specify a command. New input is accepted in the system command mode which is indicated by an asterisk as the response character on the terminal. Commands that are given during other system activity, e.g. event analysis or output on the display, are stored for later processing. The corresponding delay is normally not significant because of the centralized dispatching technique which ensures that pending terminal input is served immediately after the last scheduled task. The first character of an input string is compared to a list of valid commands and the list index is used to dispatch the according task. If the character is not found, an error message is printed, followed by another asterisk on the next

line. In the following, the meaning of all valid system commands as they are listed in appendix A, table A.7 will be discussed. If necessary, further explanations for the corresponding operation mode are given.

W: This command suppresses graphical or numerical terminal output. Recording and counting of event information will be continued according to previous mode specifications. This feature is useful to give the operator some time to observe results on the screen before they are overwritten.

P: Graphic and numeric terminal output are resumed if they were previously suspended by the "Wait" (W) command. Output will continue with the next record after P is given.

E: Once this command is given, event reception is temporarily stopped and the user can enter a list of event codes, called special events. After termination of this input mode, CS will continue to receive events and compare every new EC to the list of special events. If the EC matches, a message announcing a special event along with the complete event record is printed. The "Wait" command is ignored in this case. Simultaneously, the host system is stopped until a "Continue" (C) is entered.

C: If the host was suspended due to a special event, this command will release execution. If the host was not disabled, an error message will be printed.

B: If graphic output was specified, the command will reset the counters of the graphic routine and a coordinate system with the initial scale for histograms will be drawn

on the Tektronix screen. It is recommended to issue this command before the graphic output mode is entered.

A: In graphic mode the screen will be erased and the histogram is drawn again with all columns as they were given before. This feature is necessary because the display cannot be erased partially. Although three quarters of the screen are reserved for command and output text, it can happen that the picture is overwritten.

D: After erasing the screen, a summary count of all events received so far in this session is displayed on the terminal.

S: The command interpreter for the host table update routine is invoked. The command structure is the same as it was described for S-MON. Any host system activity (except S-MON) is interrupted during this procedure until the new event configuration is established.

H: The hardware monitor control word can be defined and a new H-MON mode started. This command has not been implemented so far and does not cause any actions.

L: This command is used to define the logging mode for hard and software events by means of a subcommand interpreter (CSIL). A self explanatory message will ask for further specifications. The operator has the following three possibilities:

OLD: An existing specification file can be used to define the logging mode. The corresponding file name has to be entered and, if desired, the file can be altered.

NEW: A new mode is defined from the terminal. Consequently the operator can specify if the session is to be

recorded and/or displayed in numerical or graphical form. For the recording mode an output file name has to be entered.

UPDATE: The current mode is displayed and for each item as discussed in the last case, a new specification can be entered. The possibility to open a new output file is given.

After the command sequence for one of the three possibilities, the mode specification can be saved on a new command file or in the case of "OLD" it can be rewritten in the unchanged or altered form on the original file. Upon termination of CSIL the new logging mode is active.

While processing the commands, SYSCMD suspends any logging action and in some cases it blocks host activities by disabling its S-MON interface. Although this solution interferes with the operation of the host, it was chosen to insure a well defined state of HIMOS during and after modifications of its operation mode. For commercial applications another method could have been used (e.g. resetting the host to the unmonitored state before processing any commands), but for research and development environment, the technique implemented was proven to be sufficient and necessary when the system state has to be preserved until a new mode is defined.

4.5.3 Event Record Processing

Upon reception of the first interrupt from S-MON which announces the transmission of a new event record, control is passed to the receiving routine (RCVREC). Two alternating buffers of 256 words each are available to store incoming data. A flag indicates the current input buffer and a pair of pointers delimit the last received record within one buffer. As soon as one buffer is filled, RCVREC sets bit 1 in DISPWD to schedule the buffer recording routine (RECORD) if recording mode was specified. Record and buffer output can be processed concurrently with the reception of new records.

After receiving one record RCVREC sets bit 0 in DISPWD to schedule the FILTER and returns control to the main dispatcher. The wait loop will then continue to examine the dispatch word bit by bit in a round robin fashion and invoke the corresponding tasks. If, for example, it was interrupted before processing bit 2, it will first invoke the terminal output routines to process the previous event record before calling FILTER to analyse the just received record. This technique insures that no information is lost. If more than two records arrive at the input before the previous one was processed, reception is disabled until the CS has caught up. This only happens for high event density or when one of the slow terminal output modes is specified. More than two records cannot be buffered prior to the analysis because the system would lose its capability to react in real time upon the occurrence of special events.

If the "Wait" (W) command was given, FILTER will skip the terminal output dispatcher. Otherwise bit 2 or 3 are set according to the mode specification. In case special events were defined, the event table is searched and if the EC was found, a message is printed and the numeric output routine (DUMP) is scheduled via bit 2. The acknowledgement signal to release the host is only transmitted if the event was not found in the table. Otherwise the host keeps waiting for this signal until the "Continue" command (C) is entered from the keyboard.

Finally the current system time is appended to the event record and the count table for the summary report (command D) is updated before FILTER returns control to the wait loop. To reduce the monitor overhead, only one real data reduction is performed in FILTER, namely the special event feature which is especially useful when terminal output was suppressed via the "Wait" command. In this case only the special events are reported in real time while there is still the possibility to record a detailed report of the session on tape or disk for later analysis.

The two terminal output routines DUMP and GRAPH find the record in one of the two buffers according to the pointers which were set as record delimiters in RCVREC. DUMP transforms the record into a readable form and prefixes every value by a two letter identification. Each record is printed on a new line.

Subroutines for the graphic output are found in the object module LOGGRP. The program maintains its own table (ECTAB) to count received events. Every EC which is not yet entered

C in the table will display the new identification along the X-axis, print the initial column increment, and open a new table entry. Upon overflow of one of the event columns, all counts in ECTAB are divided by two and the Y-axis of the display is rescaled. The column increment per event occurrence depends on the current scaling factor.

If no detailed information about host activities is required, the graph mode provides a convenient summary report. The histogram which is updated in real time shows not only what the host is doing but also in what sequence the events happen. It is mainly used to get a first impression of the monitored process and it helps to define the monitoring session for a more detailed analysis.



CHAPTER 5

Application of HIMOS5.1 Execution Analysis under RT-11

One of the design goals of HIMOS was to develop a technique for real time and off line analysis of program execution. There are chiefly two approaches to this problem labelled the active and the passive methods. The latter one consists in the application of standard features provided by the monitor. System events such as I/O activity, executive calls, error traps, etc. can be intercepted. To obtain statistical information, the PC and other important system variables can be sampled. In both cases the target program does not need to be changed.

The active method requires certain changes in the observed program. Check points or hooks are inserted to the source code in order to emit a hook identification to the monitor when they are executed. (This method is similar to the break point technique used in debugging routines). For PDP 11 assembler programs, the TRAP instruction is used to identify a softhook. TRAP works in the same way as the EMT instruction; only it uses a different vector location. The high byte contains the operation code while the low byte holds the user defined hook identification. When, for example, TRAP 123 is executed, an interrupt to location 34 is generated and control is transferred to S-MON which will record event 7 with identification 123.

The TRAP instruction was chosen for several reasons to implement software hooks. It offers the programmer a quick and easy way to identify strategical parts of his program with a one word instruction. Once the hooks are inserted in a given routine, they can be disabled just by replacing the S-MON address in location 34 by an address containing a RTI instruction. Substituting the TRAP code by a NOP instruction eliminates the hook completely without affecting the program performance. Another low level hook technique should be mentioned in case the TRAP instruction is used for other purposes. Interrupts to S-MON can be created by an additional interface. The hook identification can then be transmitted by moving the desired code into the output buffer which is fed back in a closed loop to the interface input buffer from where S-MON can take the information.

To monitor execution of FORTRAN programs, two subroutines (INEMIT, EMIT) were written and can be linked to the main program. Both subroutines are listed in appendix B.3. EMIT finds the hook identification as the first argument of the call in the FORTRAN linkage table and uses it to form the corresponding TRAP instruction. INEMIT has to be called for initialization purposes before EMIT is invoked the first time. This is necessary because FORTRAN alters the TRAP vector for its runtime error message handler. EMIT calls should therefore only be used in fully debugged programs. Where this is not possible, the software hooks have to be implemented by means of the above mentioned alternative.

Software hooks are a valuable tool for many applications. They can be used to time execution over critical parts of a

program as will be shown in chapter 5.2.1. When they are placed on top of the executable code, they can emit a task identification for accounting purposes. If, additionally, event 6 (EMT) with subselection for EMT 350 (EXIT) is monitored, the time between the EMT call and EC 6 gives the total execution time of a job.

Under the RT-11 operating system, S-MON offers different possibilities to measure "productive" CPU time. The busy wait loop of RT-11 can be recognized by alternating calls of EMT 340 and EMT 360. Another approach is given by sampling the PC activity. With regard to CPU activity the memory can be grouped into several regions:

1000 - X	:	User program execution
X - Y	:	I/O handlers
122222 - 124154	:	RT-11 EMT handler
		(if booted for 24K of core)

The values of X and Y depend on the individual program and can be specified via system macro calls. During execution the value of Y can be found in location 50 where it is stored by RT-11 at load time.

In any case, the definition of "productive" CPU time is relative. A true wait state as it is known in other systems does not exist under RT-11. Similar problems arise when I/O time is to be determined. Sampling the I/O handler space in core is one possibility. In chapter 5.2.2 another approach is taken. Disk access time is obtained as the difference between EMT 375 and a disk interrupt. EMT 375 invokes the I/O dispatcher in RT-11 which initiates a disk action. When the disk controller has found the addressed track and

transferred the complete record as it was specified in the word count, it gives an interrupt (EC 30). It was found that the average delay between those two events is in the range of 35 - 45 ms depending on the occupied disk space. This value is in accordance with RK-11 specifications. Undoubtedly the hardware monitor will be very useful to complement these results for more detailed investigations.

Generally, several steps are necessary to monitor a given process. An initial approach is to display global execution characteristics in graphical form on the terminal. Based on this first impression, more detailed measurements can be specified and event records can be directed in numerical form to the terminal or to mass storage for off line evaluation. Figure A.5 shows a histogram obtained during the execution of a test batch stream. Event 5 (power fail) was specified as a special event. It occurred at the end of the batch stream when the 11/45 power was turned off.

5.2 Results of two Monitoring Sessions

The following two paragraphs will give examples of monitoring sessions using the active and the passive methods of execution analysis. In both cases, the machine level representation of the programs is unknown and any observed PC values cannot be matched readily to a functional entity of the code.

5.2.1 FORTPAN Program Execution

A program (SAMGEN) was written to produce J random numbers in the range from 0.00 to 99.99. When the routine is entered, the value of J is requested from the terminal. In the following, the program executes J times a loop in which a random number is generated and immediately written onto disk. The random numbers are obtained from the FORTPAN function "RAN". It should be noted that the program was written to give an example for the possibilities of HIMOS and not to optimize a given problem. A listing of SAMGEN can be found in appendix B.4.

The program was executed twice, each time producing 20 samples. For the first run, SAMGEN was linked with a FORTPAN library which replaced all floating point instructions by software routines. The second time the floating point processor was used which improved the execution time but did not change other characteristics of the execution profile.

In figure A.7 the monitor output for a SAMGEN execution using the floating point processor is listed. The HIMOS command sequence to define this session on the control system is shown in figure A.8. When the command "RUN SAMGEN" is given, RT-11 executes several EMT calls to load the program into core (EMT 374, 375). The first TRAP (EC 7) identifies the EMIT call with argument 0 which can be found in the low byte of CP (content of PC). A delay of approximately 2 seconds follows, during which a message is printed and J read from the terminal. EMIT (1) and EMIT (2)

are produced before and after the first random number generation. During the following 6 EMIT's the output file is opened and at EMIT (3) the first random number is written on disk. These three emitter calls appear another 10 times during the loop execution before EMIT (4) indicates the end of the program. Six EMIT calls are issued to close the output file and enter it into the disk directory before EMIT 350 indicates the RT-11 EXIT call.

For the first run (run A) an execution time of 234 ms (excluding terminal I/O and file handling) was calculated. The second run using FPP (run B) took 193 ms. This difference is due to the fast performance of the floating point processor but it also reflects a certain delay for the disk output. The latter factor was excluded by averaging the time intervals between EMIT (1) and EMIT (2). After subtraction of 2.36 ms per event for HIMOS overhead, the final result shows a random number generation time of 0.270 ms for run A and 0.100 ms for run B. The difference of 170 micro seconds per random number generation can be attributed to the processing time improvement for floating point operations.

5.2.2 FORTRAN Compiler Analysis

In the following example, a short FORTRAN program was compiled under several HIMOS modes to obtain insight into the RT-11 FORTRAN compiler characteristics. The test program HISTO is listed in appendix B.4. It can be used to read a file of sample values from disk and to print a distribution of these values in the form of a horizontal

histogram on the terminal. At the same time, a second file containing the values in ascending order is created.

RT-11 supports a multipass FORTRAN IV compiler. It consists of a root segment and several overlays which are dynamically loaded in core. The core resident part comprises of the root and one overlay, the system stack, device handlers, symbol table, and an internal representation of the source program. During the monitoring sessions it was discovered that the compiler uses no scratch files, only overlays are swapped between core and disk. A high level description of the compiler can be found in RT-11 manuals [52, 53].

Because no source listings of the compiler were available, the analysis was done in a top down fashion. In the first step a graphic performance profile was produced on the Tektronix screen (fig. A.6). The histogram shows five different events:

EC 24 : Timer interrupts

EC 20 : Terminal input

EC 21 : Terminal output

EC 30 : Disk interrupts

EC 6 : RT-11 calls for file handling.

EC 21* is more than twice as high as EC 20, because every typed character is echoed together with one additional filler character. The graph gave an impression of the amount of events and their relation in time.

For the second monitor session it was decided to select only disk interrupts and to record the protocol on disk. Figure A.9 shows the corresponding HIMOS output. The first

(1) 11 events report directory access for the source file HISTO.FOF and the output file HISTO.OBJ, both located in the 4th directory segment (block 16). At system time (ST) 103.18, five blocks of compiler overlay are read in before at ST 103.24 the source program is loaded from disk address (R3) 6531. (All file references were confirmed by comparison with the directory listing). After having read two more overlays, the compiler writes the first block of object code at ST 103.50 on disk. The content of R4 (bus address register) shows that the compiler uses always the same buffer at location 10520 to store one block of produced object code. Compilation is continued by alternating overlay input and object output until ST 106.06 when the last FORTRAN overlay is loaded into core. It is probably used to generate completion messages. Accesses to disk address 116 and 143 show that RT-11 utility routines for terminal I/O are brought into core. Eight directory accesses starting at ST 106.21 close the source and object files.

O Already this simple monitoring session shows several performance bugs in RT-11 and the compiler. Directory access time can be reduced when the search is started within the last directory segment which is still in core. The probability of finding a file entry close to the last accessed one is higher than to find it in the beginning of the directory. A similar technique is found in memory paging algorithms where it is known as the "least recently used" method (see for example [37]). Further I/O time could be saved if the object code buffer size in the FORTRAN

compiler is increased. Corresponding changes in RT-11 and the compiler would reduce the compilation time at least by 12% in this example. The fact that all disk interrupts occurred while the PC was pointing to RT-11 indicates that the compilation process is I/O bound. The program transfers control back to RT-11 for the next I/O request before the previous one is terminated. RT-11 executes a wait loop between locations 127130 and 127124 until the next interrupt arrives.

To obtain more details on the relationship between I/O and computing activity, the compilation was sampled at different frequencies. Figure A.10 shows the HIMOS output for 40 ms samples and fig. A.11 lists the corresponding commands for this session. The PC values show that less than one fifth of execution time is spent in the FORTRAN compiler. Another sampling session with 10 ms intervals confirmed this result. As expected, the compiler is only active immediately after disk interrupts.

The overhead imposed by HIMOS during the sampling sessions amounted to 5% of the total compilation time. It is interesting to see that the monitor overhead did not increase significantly with a higher sampling rate. This can be attributed to the fact that most samples were taken during I/O wait periods.

A last experiment was done to check if compilation time could be reduced by locating the FORTRAN compiler and the object file on the disk closer together. The disk space between the two files was reduced by 1172 blocks which corresponds to 45 tracks. As a result, the compilation was

400 ms quicker. Because the disk arm moved 12 times from one file to the other, the average saving per arm movement is approximately 33 ms.

5.3 Further Possibilities for HIMOS Assisted Research

Apart from the possibilities demonstrated in the last chapter, there are mainly three research topics which have influenced the implementation of HIMOS. The following paragraphs will point out how the monitoring system can be used to assist this research.

During the past few years, several techniques have been developed to prove the correctness of a given high level language program. One method, called the "deductive theory" was presented by Hoare and Lauer [35]. The meaning of a language is defined by axioms together with rules of interference. Theorems are derived from the axioms by means of these rules. The basic axiom of the theory is written as $P\{Q\}R$ where P and R are assertions about the memory state before and after execution of the statement (or program) Q . P states the necessary entry conditions for Q , and R is calculated from P and Q using the theorems of the deductive theory. If Q terminates and if P was true before execution, then R will be true when execution is complete. Given a linear sequence Q of statements without branches or entry points, logical expressions for P and R can be established. When the sequence is executed on the host machine, memory values can be monitored to compute the expression for P . When R is true upon exit from Q , the sequence has been executed correctly according to the programmer's intentions.

If, for example, a programmed loop (L) is observed, two software hooks are sufficient to verify the correctness of execution no matter how many statements are involved. The first hook is located in front of L and causes the verification of P in the control system. If P is not true, the process should be aborted because the termination of the loop cannot be guaranteed. The second hook is inserted immediately after the exit of L. At this point, the evaluation of R will show if the loop was executed correctly. Memory values that are necessary to compute the assertions can be obtained at interrupt time either through the hardware monitor directly from memory or by means of the software monitor from general purpose registers.

This method has the advantage of verifying correctness of programs during execution without producing an excessive amount of data as the trace technique does. It also reduces the necessary preparation time for the calculation of assertions which is required for purely theoretical correctness proofs.

The prevention and detection of deadlocks constitutes another promising application field for monitors. Although HIMOS was implemented for RT-11 which is a single user system under which this problem cannot be encountered, it is felt that a short discussion of this topic is justified by the importance of the problem. A deadlock is a situation in which parts of, or even all system activities are blocked due to an erroneous scheduling decision. A fatal deadlock situation may, for example, arise when the spooling space becomes completely filled with input records for jobs

waiting to execute and with output records for jobs not finished executing. In many systems there is no way to recover the spooling space which is occupied by a partially executed job. Several ad hoc methods have been employed to reduce the probability for deadlocks without giving a guarantee for their complete elimination. Other approaches, mostly based on graph models (see for example [54]), attempt to formalize the problem and to construct detection and recovery algorithms. Hybrid monitors offer a compromise in the detection and prevention of deadlocks. The software part can observe the task arrangement in waiting queues of the system as well as possible relations between those tasks. This may lead to the detection of situations in which tasks are waiting for each other without being able to resolve their requirements. The workload of critical resources can be observed by the hardware monitor. If, for example, the utilization of spooling space becomes too high and the monitor control system obtains at the same time the information that the waiting queue for this resource is growing, an imminent deadlock is indicated. In such situations a warning can be displayed at the operator's console or the monitor can supply a corresponding message directly to the operating system.

Finally, HIMOS can give practical assistance for efficiency analysis of paging algorithms. The memory management unit in the PDP 11/45 divides the full address range of 128 K words into 48 pages of 32 to 4096 words each. When a location outside of the current page limits is accessed a trap to address 250 is generated. When this trap vector is

intercepted by S-MON, the occurrence of page faults can be monitored. In sampling mode, S-MON can examine the content of the page address registers to record the current memory partitions. This information can be complemented with a memory utilization report from H-MON when it is sampling the PC activity.

The examples outlined in this and the previous chapter are far from giving a complete survey of applications for HIMOS. Individual requirements and practical experience with the system should motivate exploration of further possibilities of the monitor.

CHAPTER 6

Evaluation and Conclusion6.1 Evaluation of HIMOS

The monitor has proven to be a reliable tool for performance measurements. Once acquainted with the system, the user can dispense with a large set of control commands. Several command files provide the possibility to predefine various monitoring sessions which can be evoked without profound knowledge of system internal details. All monitor results reflect exactly the host computer activities and HIMOS reports are reproduceable at any time. The system's greatest successes have been the following:

- a. Its centralized control allows the specification of all functions from the terminal at the supervisor system. Erroneous command input or ambiguous instructions (such as two different output modes on the same terminal) are detected and comprehensive error messages are generated. System internal plausibility checks recognize software and hardware errors.
- b. Real time data filtering reduces the amount of recorded and displayed information. Only pertinent data is retained for further analysis.
- c. Its interactivness enables the user to redefine a monitor session according to new requirements while host activities are suspended at a defined state. Host execution can be slowed down to follow up event occurrences in real

time on the display.

- d. Three different output modes of which two can be handled concurrently, satisfy all needs for online and offline analysis. The graphic output produces a dynamic picture of system activities. A modular implementation facilitates future extensions for other integrated evaluation features.
- e. Most failures of the host CPU and any peripherals attached to it can be reported. Analysis of such events is assisted by the accompanying status report.
- f. All necessary information for various applications is supplied, namely:

- Performance improvement
- System profile, activity log
- Problem oriented sessions (debugging)
- Activity reports for simulation model input.

Yet, HIMOS also has some deficiencies, specifically:

- a. The command handler for host system mode specifications accepts a very abbreviated form of input. Some familiarisation is required to define the host tables.
- b. The monitor overhead is relatively high. Creation and transmission of one event record takes from 1.43 ms to 3.14 ms, depending on the length of the record. During this time the host computer is exclusively dedicated to monitor activities. This time requirement is explained by the following 3 reasons:

- 1. To maintain the capacity of interactive response to special events, monitor actions cannot overlap with

host operations to serve user or RT-11 tasks.

2. Event records are encoded into ASCII prior to their transmission.

3. Event records are analyzed in the control system to detect special events before the host system is released for normal execution.

c. Because S-MON was implemented under RT-11 its possibilities are limited to applications within a single user system.

d. The operating system has to be slightly modified to allow undisturbed operation of S-MON. With the exception of the FORTRAN run time error message handler, the changes do not affect any system properties.

The hardware monitor will most probably be implemented by the end of this year. Because of time and equipment limitations, many desirable features were excluded in the final design. Due to the interrupt latency of the PDP 11/20 CPU, the maximum sampling frequency had to be limited to 10 KHz. A DMA link between H-MON and the control system would offer a variety of further applications. As the control software includes already all necessary routine entries to handle the communication to and from H-MON, its adaptation to HIMOS will not create major problems.

6.2 Conclusion

Measuring computer performance during program execution can be compared with taking the blood pressure of a man under effort. Monitoring can actually be compared to

medicine. The only major difference lies in the complexity of the observed systems. In both cases the performance is evaluated and important points for improvement or reasons for failure can be found.

The necessity to monitor activities in large computer systems has been generally recognized. The objectives are not only to increase cost efficiency of an installation but also to ensure reliability and to obtain hints for further efforts in research and development.

Mini computers have conquered an important market place during the past decade. Among other applications they are especially well suited for real time applications in process control. In this environment high reliability is an essential factor and monitors are a valuable tool to insure this objective.

During the implementation of HIMOS, several deficiencies in the design of computer hardware and software became evident. It is hoped that future hardware design will consider probe point requirements for hardware monitors. Especially with regard to the increasing trend of miniaturization and large scale integration there is a danger that many important signals become buried in a single chip. Monitor facilities should be a part of the standard equipment of CPU and peripheral hardware. One step in this direction has, for example, been taken by DEC. The KL10 processor produced by this company incorporates a PDP 11/40 CPU which is partly used to monitor hardware signals.

The same arguments can be cited for the development of operating system software. Information from monitoring

devices is useful to assist operating system decisions. Especially when software reaches the magnitude of OS/360, additional observation tools are indispensable.

Consequently two different approaches can be stated for system development with respect to monitors. The monitor can form an integral part of the system's hard and software. This approach was, for example, taken during the development of "Multics". However, this solution seems to have some dangers. It is quite possible that weak points in the system are only confirmed by corresponding deficiencies in the monitor, especially if the same team is responsible for implementing both parts.

A modular approach using distinct systems seems therefore more promising. It is quite possible that the standard configuration of future computer installations will dedicate one processor exclusively to control and measure the remaining system components. This method was adopted for the implementation of HIMOS. It is believed that HIMOS will constitute a valuable tool for future projects on the mini computer installation at McGill.

"The purpose of measurement is insight, not numbers."

(P.W. Hamming)

APPENDIX A

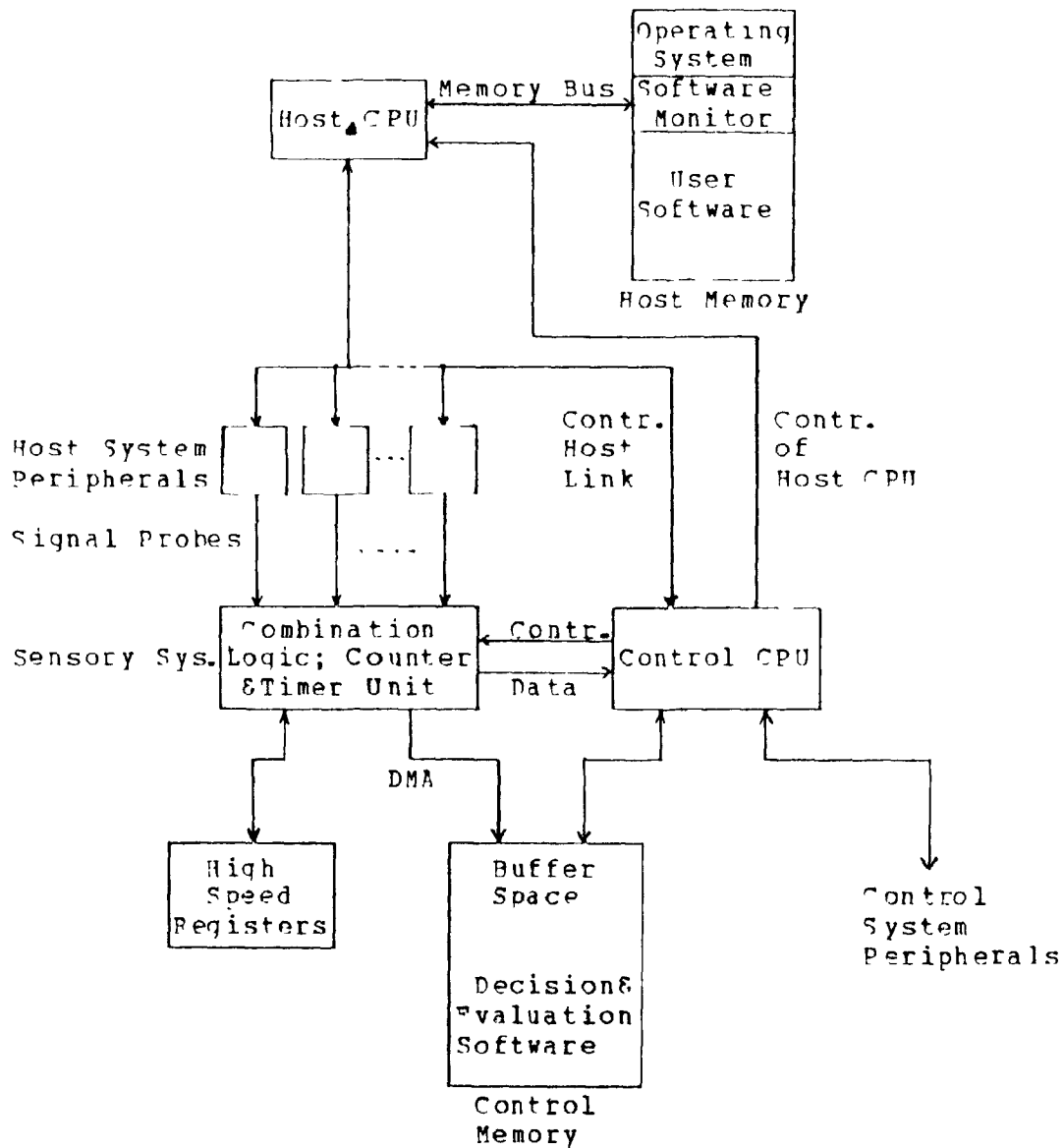
Figures and Tables

Figure A.1: A General Instrumentation System

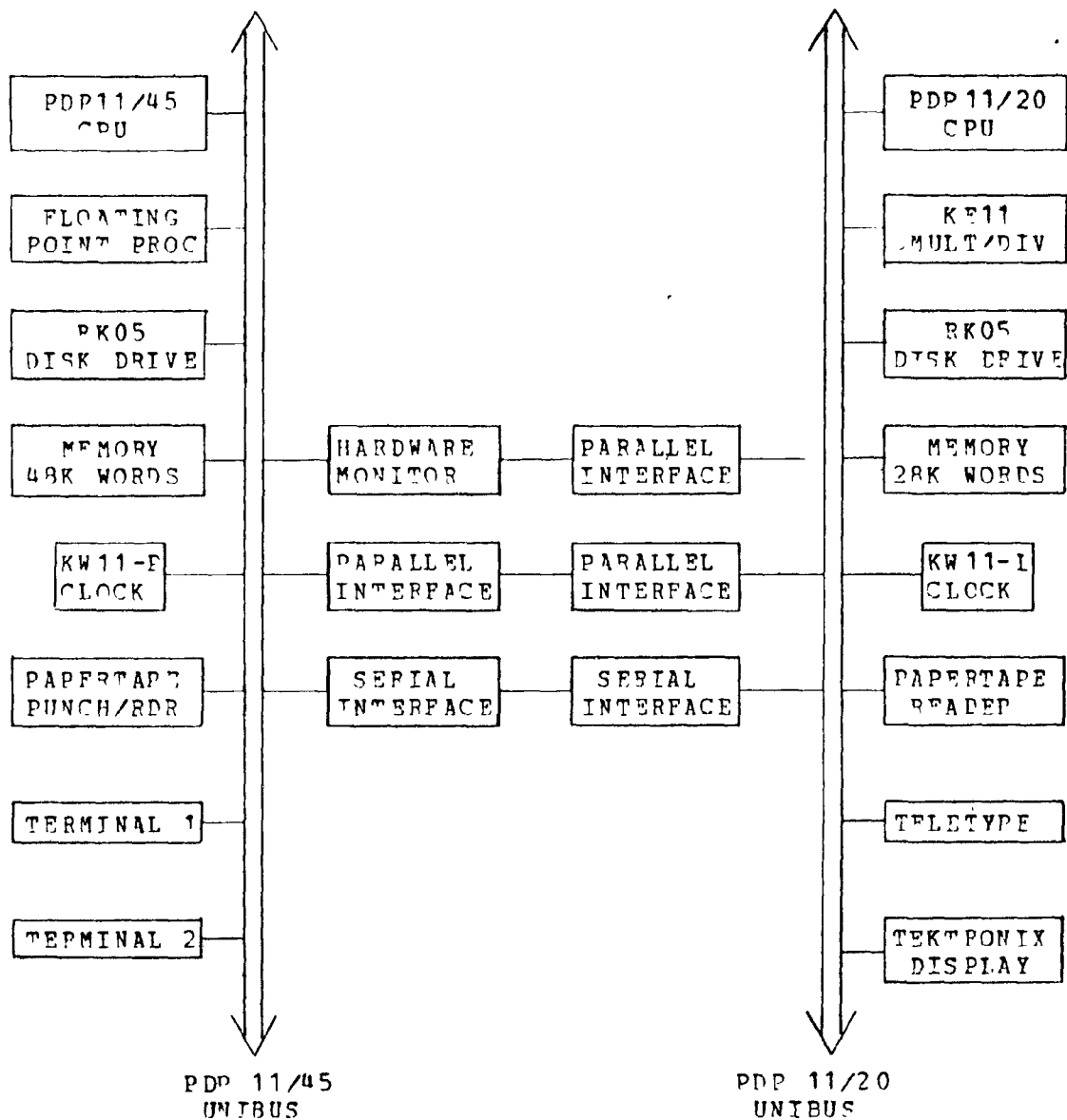


Figure A.2: System Configuration

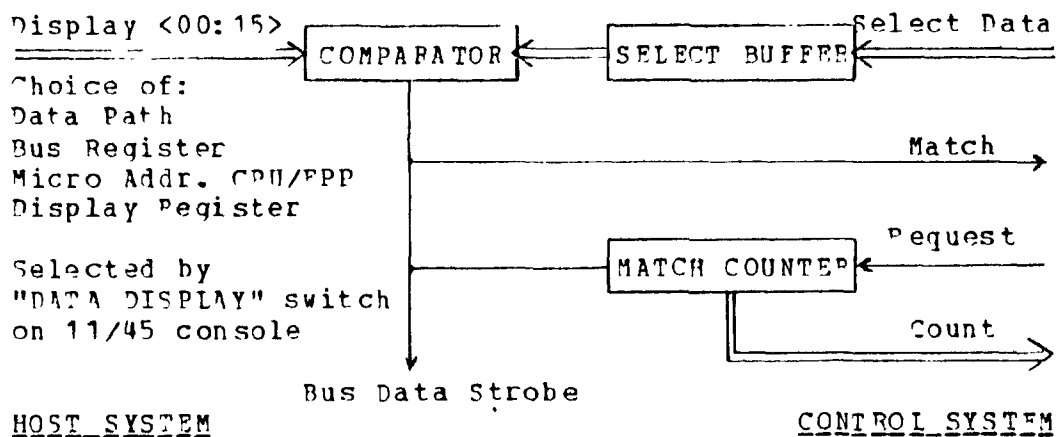


Figure A.3: Display Comparator

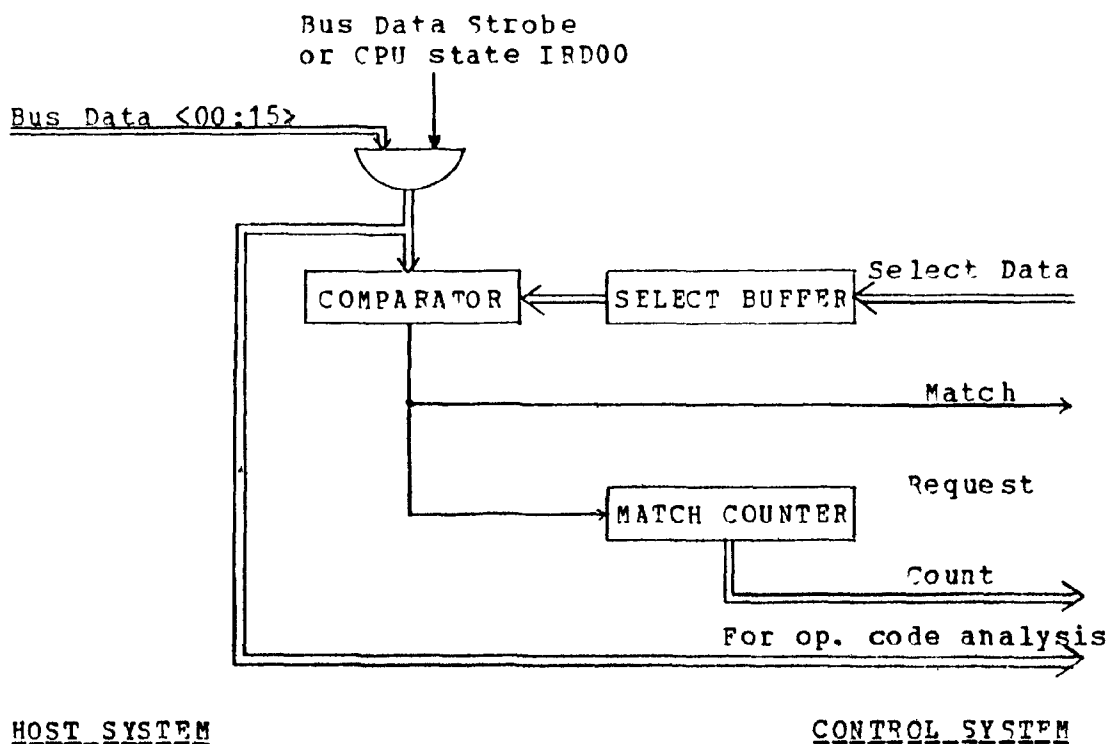


Figure A.4: Operation Code Analysis on UNIBUS Data Lines.

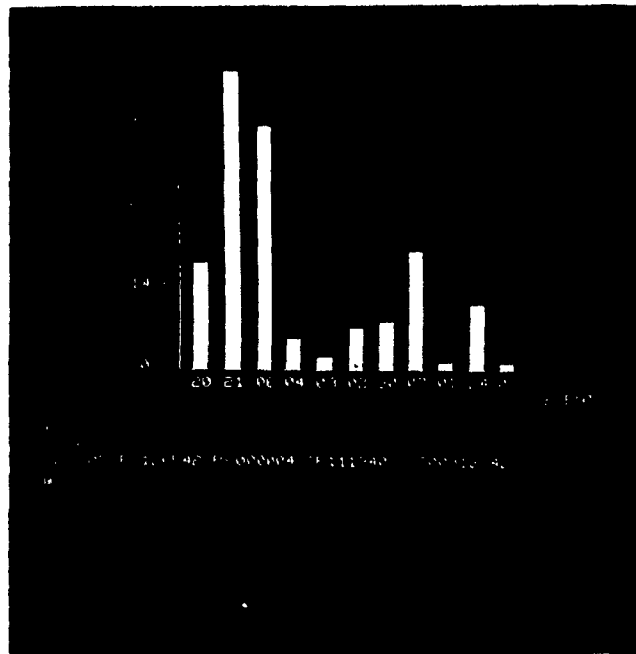


Figure A.5: Performance Profile of a Test Batch Run

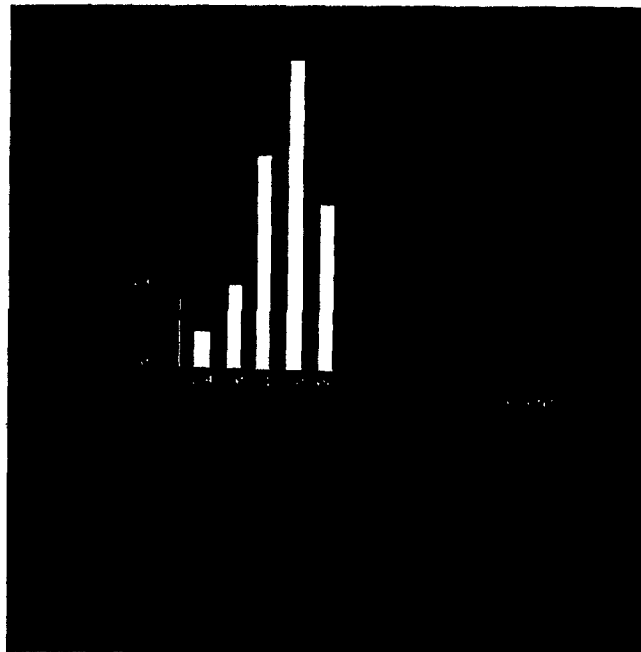


Figure A.6: Performance Profile of a FORTRAN Compilation

HIMOS LOG 14-7-76, 15:43, SAMGEN RUN UNDER FPP . EC6,7

000006	PC106004	CP104374	EW000000	HT152345	ST00275.11
000006	PC115546	CP104375	EW000000	HT152407	ST00275.12
000006	PC115546	CP104375	EW000000	HT154216	ST00275.17
000006	PC115546	CP104375	EW000000	HT155135	ST00275.20
000006	PC115546	CP104375	EW000000	HT156042	ST00275.23
000006	PC115546	CP104375	EW000000	HT156250	ST00275.23
000006	PC013354	CP104375	EW000000	HT162405	ST00275.36
000006	PC013420	CP104375	EW000000	HT162437	ST00275.36
000007	PC001506	CP104400	HT162516	ST00275.37	
000007	PC001506	CP104401	HT030077	ST00277.32	
000007	PC001506	CP104402	HT030130	ST00277.32	
000006	PC005160	CP104375	EW000000	HT034073	ST00277.44
000006	PC005160	CP104375	EW000000	HT035200	ST00277.47
000006	PC005160	CP104375	EW000000	HT036120	ST00277.50
000006	PC005160	CP104375	EW000000	HT037012	ST00277.53
000006	PC005160	CP104375	EW000000	HT037241	ST00277.54
000006	PC005160	CP104375	EW000000	HT037727	ST00277.56
000006	PC005160	CP104375	EW000000	HT040152	ST00277.56
000006	PC005160	CP104375	EW000000	HT040706	ST00277.59
000007	PC001506	CP104403	HT043213	ST00278.06	
000007	PC001506	CP104401	HT043242	ST00278.06	
000007	PC001506	CP104402	HT043275	ST00278.06	
000007	PC001506	CP104403	HT043345	ST00278.06	
000007	PC001506	CP104401	HT043374	ST00278.07	
000007	PC001506	CP104402	HT043427	ST00278.07	
000007	PC001506	CP104403	HT043477	ST00278.07	
000007	PC001506	CP104401	HT043526	ST00278.07	
000007	PC001506	CP104402	HT043561	ST00278.07	
000007	PC001506	CP104403	HT043631	ST00278.07	
000007	PC001506	CP104401	HT043660	ST00278.08	
000007	PC001506	CP104402	HT043713	ST00278.08	
000007	PC001506	CP104403	HT043771	ST00278.08	
000007	PC001506	CP104401	HT044020	ST00278.08	
000007	PC001506	CP104402	HT044052	ST00278.08	
000007	PC001506	CP104403	HT044123	ST00278.09	
000007	PC001506	CP104401	HT044152	ST00278.09	
000007	PC001506	CP104402	HT044204	ST00278.09	
000007	PC001506	CP104403	HT044261	ST00278.09	
000007	PC001506	CP104401	HT044311	ST00278.09	
000007	PC001506	CP104402	HT044342	ST00278.09	
000007	PC001506	CP104403	HT044375	ST00278.10	
000007	PC001506	CP104401	HT044447	ST00278.10	
000007	PC001506	CP104402	HT044507	ST00278.10	
000007	PC001506	CP104403	HT044563	ST00278.10	
000007	PC001506	CP104401	HT044613	ST00278.10	
000007	PC001506	CP104402	HT044645	ST00278.11	
000007	PC001506	CP104403	HT044716	ST00278.11	
000007	PC001506	CP104401	HT044745	ST00278.11	
000007	PC001506	CP104402	HT044777	ST00278.11	
000007	PC001506	CP104403	HT045053	ST00278.11	
000007	PC001506	CP104401	HT045107	ST00278.12	
000007	PC001506	CP104402	HT045171	ST00278.12	
000007	PC001506	CP104403	HT045230	ST00278.12	

Figure A.7: Event Report for "SAMGEN" Execution under FPP (1)

HIMOS LOG 14-7-76, 15:43, SAMGEN RUN UNDER FPP, EC6,7

000007	PC001506	CP104401	HT045260	ST00278	12
000007	PC001506	CP104402	HT045312	ST00278	12
000007	PC001506	CP104403	HT045367	ST00278	13
000007	PC001506	CP104401	HT045417	ST00278	13
000007	PC001506	CP104402	HT045451	ST00278	13
000007	PC001506	CP104403	HT045525	ST00278	13
000007	PC001506	CP104401	HT045555	ST00278	13
000007	PC001506	CP104402	HT045607	ST00278	13
000007	PC001506	CP104403	HT045664	ST00278	14
000007	PC001506	CP104401	HT045714	ST00278	14
000007	PC001506	CP104402	HT046035	ST00278	14
000007	PC001506	CP104403	HT046112	ST00278	15
000007	PC001506	CP104401	HT046142	ST00278	15
000007	PC001506	CP104402	HT046174	ST00278	15
000007	PC001506	CP104403	HT046250	ST00278	15
000007	PC001506	CP104401	HT046300	ST00278	15
000007	PC001506	CP104402	HT046332	ST00278	16
000007	PC001506	CP104403	HT046407	ST00278	16
000007	PC001506	CP104401	HT046437	ST00278	16
000007	PC001506	CP104402	HT046471	ST00278	16
000007	PC001506	CP104403	HT046546	ST00278	17
000007	PC001506	CP104401	HT046713	ST00278	17
000007	PC001506	CP104402	HT046745	ST00278	17
000007	PC001506	CP104403	HT047021	ST00278	17
000007	PC001506	CP104404	HT047051	ST00278	18
000006	PC005160	CP104375	EW000000	HT055734	ST00278. 39
000006	PC005160	CP104375	EW000000	HT056614	ST00278 41
000006	PC005160	CP104375	EW000000	HT057044	ST00278. 42
000006	PC005160	CP104375	EW000000	HT057764	ST00278. 45
000006	PC005160	CP104375	EW000000	HT060670	ST00278. 48
000006	PC005160	CP104375	EW000000	HT061110	ST00278. 48
000006	PC006430	CP104350	EW000000	HT063437	ST00278. 56

Figure A.7: Event Report for "SAMGEN" Execution under FPP (2)

```

*L
LOG MODE DEFINITION
H-LOG OR S-LOG SPECS (H,S) : S
OLD, NEW OR UPDATE (O,N,U) : N
LOG RECORDING (Y,N) : YES
NUMERIC DISPLAY (Y,N) : N
GRAPHIC DISPLAY (Y,N) : N
SAVE FILE (Y,N) : Y
  OUTPUT FILE FOR SPECS
*SES1=
  OUTPUT FILE FOR LOG RECORDING
*LOG5=

*S
HOST TABLE UPDATE
SPEC FILE-OLD OR UPDATE (O,U) : U

MASK UPDATE (Y,N)? : Y
GIVE INDEX OR T
#0
7654321076543210
0000000000001110
0000000011001110
#T

SUBMASK UPDATE (Y,N)? : Y
GIVE INDEX OR T
#1
7654321076543210
0001001000110010
0001001000110011
#T

SELECT WORD UPDATE (Y,N)? : Y
GIVE INDEX OR T
#0
350 :
351 : 343
#1
340 : 374
341 : 375
#T

TIMER FREQUENCY (>3: NO TIME LOG): 1
TIMER COUNT : 0
SAVE FILE (Y,N) : NO
*D

EVENT COUNT:

01:000000 02:000000 03:000000 04:000000 05:000000
06:000027 07:000076 10:000000 11:000000 12:000000
13:000000 14:000000 15:000000 16:000000 17:000000
20:000000 21:000000 22:000000 23:000000 24:000000
25:000000 26:000000 27:000000 30:000000 31:000000
32:000000 33:000000 34:000000 35:000000 36:000000

```

Figure A.8: Command Sequence for HIMOS Session of Fig. A.7

HIMOS LOG 12-7-76, 16 12, FORTRAN COMPIL HISTO FOR, EC30

000030	PC127130	R3	000010	R4	113366	HT155506	ST00102	52
000030	PC127130	R3	000012	R4	113366	HT156426	ST00102	55
000030	PC127130	R3	000014	R4	113366	HT157345	ST00102	58
000030	PC127130	R3	000020	R4	113366	HT157552	ST00102	59
000030	PC127124	R3	000016	R4	113366	HT160265	ST00103	01
000030	PC127130	R3	000016	R4	113366	HT161103	ST00103	03
000030	PC127124	R3	000010	R4	113366	HT161411	ST00103	04
000030	PC127124	R3	000012	R4	113366	HT162331	ST00103	07
000030	PC127130	R3	000014	R4	113366	HT163251	ST00103	10
000030	PC127120	R3	000020	R4	113366	HT163455	ST00103	11
000030	PC127124	R3	000016	R4	113366	HT164170	ST00103	13
000030	PC127130	R3	001007	R4	006400	HT165664	ST00103	18
000030	PC127130	R3	006531	P4	012654	HT167626	ST00103	24
000030	PC127130	P3	001014	R4	006632	HT174304	ST00103	38
000030	PC127130	R3	001021	R4	007010	HT176205	ST00103	43
000030	PC127130	R3	007267	R4	010520	HT000354	ST00103	50
000030	PC127124	R3	001026	R4	007152	HT002356	ST00103	56
000030	PC127130	R3	007270	R4	010520	HT005136	ST00104	05
000030	PC127124	R3	001033	R4	006344	HT007345	ST00104	12
000030	PC127130	R3	007271	R4	010520	HT011717	ST00104	19
000030	PC127130	R3	001040	R4	006502	HT014332	ST00104	27
000030	PC127124	R3	001045	R4	006604	HT016233	ST00104	33
000030	PC127130	R3	001052	R4	007144	HT020123	ST00104	38
000030	PC127124	P3	001060	R4	007360	HT021257	ST00104	42
000030	PC127130	R3	001063	R4	004316	HT023055	ST00104	47
000030	PC127130	R3	001066	R4	004646	HT024036	ST00104	50
000030	PC127130	R3	001073	R4	006406	HT026554	ST00104	59
000030	PC127124	R3	001101	R4	007520	HT030515	ST00105	04
000030	PC127130	R3	001105	P4	005624	HT031537	ST00105	08
000030	PC127130	R3	007272	R4	010520	HT033645	ST00105	14
000030	PC127124	P3	001112	R4	007204	HT036526	ST00105	23
000030	PC127130	P3	007273	R4	010520	HT042676	ST00105	36
000030	PC127124	R3	001120	R4	007276	HT045207	ST00105	43
000030	PC127130	R3	007274	R4	010520	HT047460	ST00105	50
000030	PC127130	R3	007275	R4	010520	HT050336	ST00105	53
000030	PC127130	R3	007276	R4	010520	HT051215	ST00105	56
000030	PC127130	R3	007277	R4	010520	HT052073	ST00105	58
000030	PC127130	P3	001126	P4	007330	HT054507	ST00106	06
000030	PC127124	R3	001002	R4	006500	HT055324	ST00106	08
000030	PC127130	P3	000116	R4	121766	HT057370	ST00106	15
000030	PC127130	P3	000143	P4	121366	HT060657	ST00106	19
000030	PC127130	P3	000016	R4	113366	HT061434	ST00106	21
000030	PC127130	P3	000010	P4	113366	HT061743	ST00106	22
000030	PC127124	P3	000012	R4	113366	HT062662	ST00106	25
000030	PC127130	R3	000014	R4	113366	HT063602	ST00106	28
000030	PC127130	P3	000020	R4	113366	HT064007	ST00106	29
000030	PC127130	R3	000016	R4	113366	HT064522	ST00106	30
000030	PC127130	R3	000016	R4	113366	HT065337	ST00106	33
000030	PC127130	P3	000116	R4	121366	HT067176	ST00106	38

Figure A.9: Disk Activity Report for a FORTRAN Compilation

HIMOS LOG 12-7-76, 18 03, FORTRAN COMPIL EC30, 24(40MS)

```

000024 PC126532 ST00188 39
000030 PC127130 R3 000010 ST00194 01
000030 PC127130 R3 000012 ST00194 04
000024 PC127130 ST00194 05
000030 PC127124 R3 000014 ST00194 07
000024 PC127124 ST00194 08
000030 PC127124 R3 000020 ST00194 08
000030 PC127130 R3 000016 ST00194 10
000024 PC116574 ST00194 10
000030 PC127130 R3 000016 ST00194 12
000024 PC127130 ST00194 12
000030 PC127130 R3 000010 ST00194 13
000024 PC127130 ST00194 15
000030 PC127130 R3 000012 ST00194 16
000024 PC127124 ST00194 17
000030 PC127124 R3 000014 ST00194 19
000024 PC127130 ST00194 20
000030 PC127130 R3 000020 ST00194 20
000030 PC127130 R3 000016 ST00194 22
000024 PC127124 ST00194 22
000024 PC127130 ST00194 24
000030 PC127130 R3 001007 ST00194 27
000024 PC126572 ST00194 27
000024 PC127130 ST00194 29
000024 PC127130 ST00194 32
000030 PC127130 R3 006531 ST00194 33
000024 PC005616 ST00194 34
000024 PC006266 ST00194 36
000024 PC006260 ST00194 39
000024 PC127124 ST00194 41
000024 PC127130 ST00194 44
000024 PC127130 ST00194 46
000030 PC127130 R3 001014 ST00194 47
000024 PC127130 ST00194 48
000024 PC127130 ST00194 51
000030 PC127130 R3 001021 ST00194 53
000024 PC127130 ST00194 53
000024 PC127130 ST00194 56
000024 PC127130 ST00194 58
000030 PC127130 R3 007306 ST00195 00
000024 PC127124 ST00195 00
000024 PC127124 ST00195 03
000024 PC127130 ST00195 05
000024 PC127130 ST00195 08
000030 PC127130 R3 001026 ST00195 08
000024 PC127130 ST00195 10
000024 PC127124 ST00195 12
000030 PC127130 R3 007307 ST00195 14
000024 PC127124 ST00195 15
000024 PC127124 ST00195 17
000024 PC127130 ST00195 20
000030 PC127130 R3 001033 ST00195 21

```

Figure A.10: Sampling Session with 40 ms Periods (1)

HIMOS LOG 12-7-76, 18 03, FORTRAN COMPIL EC30, 24(40MS)

```

000024 PC127130 ST00195 24
000024 PC127130 ST00195 27
000030 PC127124 R3 007310 ST00195 29
000024 PC127130 ST00195 29
000024 PC127124 ST00195 32
000024 PC127130 ST00195 34
000030 PC127130 R3 001040 ST00195 36
000024 PC006350 ST00195 36
000024 PC004422 ST00195 39
000024 PC127124 ST00195 41
000030 PC127130 R3 001045 ST00195 42
000024 PC127124 ST00195 44
000024 PC127130 ST00195 46
000030 PC127130 R3 001052 ST00195 48
000024 PC003556 ST00195 48
000024 PC127130 ST00195 51
000030 PC127130 R3 001060 ST00195 51
000024 PC002642 ST00195 53
000024 PC127124 ST00195 56
000030 PC127124 R3 001063 ST00195 56
000024 PC003426 ST00195 58
000030 PC127130 R3 001066 ST00195 59
000024 PC003720 ST00196 00
000024 PC004152 ST00196 01
000024 PC127130 ST00196 05
000030 PC127124 R3 001073 ST00196 08
000024 PC125612 ST00196 08
000024 PC127130 ST00196 10
000024 PC127130 ST00196 12
000030 PC127124 R3 001101 ST00196 14
000024 PC003356 ST00196 15
000030 PC127130 R3 001105 ST00196 17
000024 PC002426 ST00196 17
000024 PC127130 ST00196 20
000024 PC127124 ST00196 22
000030 PC127130 R3 007311 ST00196 24
000024 PC002676 ST00196 24
000024 PC127130 ST00196 27
000024 PC127130 ST00196 29
000024 PC127124 ST00196 32
000030 PC127124 R3 001112 ST00196 32
000024 PC004700 ST00196 34
000024 PC006204 ST00196 36
000024 PC127124 ST00196 39
000024 PC127124 ST00196 41
000024 PC127130 ST00196 44
000030 PC127130 R3 007312 ST00196 46
000024 PC127130 ST00196 46
000024 PC127130 ST00196 48
000024 PC127130 ST00196 51
000030 PC127130 R3 001120 ST00196 52
000024 PC005340 ST00196 53
000024 PC127130 ST00196 56

```

Figure A.10: Sampling Session with 40 ms Periods (2)

HIMDS LOG 12-7-76, 18 03. FORTRAN COMPIL EC30, 24 (40MS)

```

000030 PC127130 R3 007313 ST00197 00
000024 PC005340 ST00197.00
000030 PC127130 R3 007314 ST00197 01
000024 PC003506 ST00197.01
000030 PC127124 R3 007315 ST00197 05
000024 PC125612 ST00197 05
000024 PC127130 ST00197 08
000030 PC127124 P3 007316 ST00197.08
000024 PC127130 ST00197 10
000024 PC127124 ST00197 12
000024 PC127130 ST00197 15
000030 PC127120 R3 001126 ST00197 15
000024 PC127124 ST00197 17
000030 PC127124 R3 001002 ST00197 17
000024 PC127130 ST00197 20
000024 PC127130 ST00197 22
000030 PC127130 P1 000116 ST00197 24
000024 PC127130 ST00197 24
000024 PC127130 ST00197 27
000030 PC127130 P3 000143 ST00197 28
000024 PC127130 ST00197 29
000030 PC127124 R3 000016 ST00197 30
000030 PC127130 R3 000010 ST00197 31
000024 PC127130 ST00197.32
000030 PC127124 R3 000012 ST00197 34
000024 PC125612 ST00197 34
000024 PC127130 ST00197.36
000030 PC127124 R3 000014 ST00197 37
000030 PC127130 R3 000020 ST00197 38
000024 PC127130 ST00197.39
000030 PC127130 R3 000016 ST00197 40
000024 PC127124 ST00197.41
000030 PC127130 R3 000016 ST00197.42
000024 PC127124 ST00197 44
000024 PC127130 ST00197 46
000030 PC127130 R3 000116 ST00197 48
000024 PC122344 ST00197 48
000024 PC122344 ST00197.51
000024 PC124062 ST00197.53
000024 PC126534 ST00197.56
000024 PC122326 ST00197.58
000024 PC122440 ST00198.00
000024 PC122436 ST00198.03
000024 PC123366 ST00198.05
000024 PC122332 ST00198.08
000024 PC122444 ST00198.10
000024 PC124114 ST00198.12
000024 PC122326 ST00198.15
000024 PC126606 ST00198.17
000024 PC126614 ST00198.20
000024 PC122316 ST00198.22
000024 PC126604 ST00198.24
000024 PC126542 ST00198.27

```

Figure A.10: Sampling Session with 40 ms Periods (3)

```

*L
LOG MODE DEFINITION
H-LOG OR S-LOG SPECS (H,S) : S
OLD, NEW OR UPDATE (O,N,U) : UPD
NO LOG SPECIFIED
LOG RECORDING (Y,N) : N? : Y
NUMERIC DISPLAY (Y,N) : N? : N
GRAPHIC DISPLAY (Y,N) : N? : N
SAVE FILE (Y,N) : N
CONT. OUTPUT ON OLD FILE (Y,N) : N
  OUTPUT FILE FOR LOG RECORDING
*LOG2=

```

```

*S
HOST TABLE UPDATE
SPEC FILE-OLD OR UPDATE (O,U) : U

```

```

MASK UPDATE (Y,N)? : Y
GIVE INDEX OR T

```

```

#1
7654321076543210
0000000000000000
0000000100010000
#T

```

```

SUBMASK UPDATE (Y,N)? : N

```

```

SELECT WORD UPDATE (Y,N)? : N

```

```

INTERR. FREQUENCY: 1
TIMER COUNT : 620
AVE FILE (Y,N) : NO

```

```
"=10KHZ"
```

```
"DIV. BY 400 => 40MS SAMPLE RATE"
```

```

*E
SP.EV. TABLE UPDATE:
GIVE EVENTS (# OR T)

```

```
#1
```

```
#2
```

```
#T
```

```
STOP ON HARD EVENT (Y,N) : N
```

```
*
```

```
SP.EV.:
```

```
000001 PC000007 PS000004 SP007124 ST00331.40
```

```
*C "TO RELEASE HOST SYSTEM"
```

```
*
```

Figure A.11: Command Sequence for HIMOS Session of Fig. A.10

<u>NAME</u>	<u>REF.</u>	<u>YEAR</u>	<u>PROPERTIES</u>
SNUPPR	[21]	1967	Intercept monitor (first phase)
GECOS II	[43]	1968	Event and sample driven; output on printer or tape.
CUE	[24]	1969	Sampling monitor; distributed by Bool and Babbage Inc.
SIPE	[44]	1969	Event driven; output on tape.
PPF	[24]	1970	Sampling monitor; distributed by Bool and Babbage Inc.
SPY	[45]	1970	Sampling monitor; output on paper tape or display; controlled by PDP-9.
MULTICS	[18]	1970	Event driven; graphic display; controlled by PDP-8.
SLACMON	[51]	1970	Sampling or event driven; monitors IBM/360 under OS/MVT.

Table A.1: Survey of Pure Software Monitors.

<u>NAME</u>	<u>REF.</u>	<u>YEAR</u>	<u>PROPERTIES</u>
--	[24]	1958	First hardware monitor by IBM for laboratory use.
POEM	[24]	1963	Program oriented event monitor
SNUPER	[21]	1967	Event driven, preset by host software (second phase).
SPAP	[27]	1967	Accumulating monitor, possibil. to interrupt host system.
SUM	[24]	1969	First commercially distributed hardware monitor.
NEUTRON	[46]	1971	Minicomputer controlled; output on display and tape.
CPM-X	[34]	1972	Event driven; interaction with host software; minicomputer controlled.
DYNAPROBE	[47]	1972	Sampling and event driven; minicomputer controlled.
MICROSUM	[48]	1973	Accumulating monitor; real time display.
DSP1	[49]	1974	Accumulating and sampling; interaction with host software; output on display and tape.

Table A.2: Survey of Hardware and Hybrid Monitors

	<u>HARDWARE</u>	<u>SOFTWARE</u>	<u>HYBRID</u>
Host system degradation	none	medium to high	very low
Level of detail	medium	high	high
Operating sys. dependent	no	yes	less than soft. mon.
Training required	extensive	limited	extensive
Monitoring of op. sys. parameters	very diffic.	easy	easy
Software monit. capability	none	good	very good
Global system profile	good	medium	very good
Interrupt recognition	dep. on system	yes	yes
Device monitoring	yes	difficult	yes
Cost:			
Development	high	medium	high
Maintenance	high	low	high

Table A.3: Comparison of Monitoring Techniques.

<u>EC</u>	<u>Vector</u>	<u>Event</u>
Group 0:		
00	--	not used
01	04	Trap to 4
02	10	Trap to 10
03	14	BPT
04	20	IOT
05	24	Power fail
06	30	EMT
07	34	TRAP
10-17	--	not used
Group 1:		
20	300	Console Keyboard
21	304	Console Printer
22	70	Paper Tape Reader
23	74	Paper Tape Punch
24	104	KW-11P Clock
25-27	--	not used
30	220	RK 11 disk
31	240	PIRQ
32	244	FPP (floating point proc.)
33	250	MMU (memory management)
34-37	--	not used

Table A.4: Event Code Interpretation.

Event Masks:

bit	7	6	5	4	3	2	1	0	
high	-	-	-	-	-	-	-	-	GROUP 0
low	TRAP	EMT	PWFL	IOT	BPT	T10	T4	-	
high	-	-	-	-	MMU	FPP	PIRQ	RK	GROUP 1
low	-	-	-	KW	PP	PR	TP	TK	

Submasks:

bit	7	6	5	4	3	2	1	0	
high	-	SP	PS	PC	<u>SM2</u> SP	PS	PC	SP	PWFL, IOT
low	PS	PC	SP	PS	PC	SP	PS	PC	BPT, T10, T4
high	-	-	-	[PC-2]	<u>SM1</u> SP	PS	PC-2	SEL	TRAP
low	-	-	ERC	[PC-2]	SP	PS	PC-2	SEL	EMT
high	-	-	-	[REG]	<u>SM2</u> SP	PS	PC	SEL	TP
low	-	-	[BUF]	[REG]	SP	PS	PC	SEL	TK
high	-	-	-	[REG]	<u>SM3</u> SP	PS	PC	SEL	DD
low	-	-	[BUF]	[REG]	SP	PS	PC	SEL	PR
high	-	-	-	-	<u>SM4</u> -	-	-	-	
low	-	-	[BUF]	[REG]	SP	PS	PC	SEL	KW-11 CLOCK
<u>SM5</u> - not used -									
high	-	-	-	-	<u>SM6</u> -	-	R5	R4	RK 11 disk
low	R3	R2	R1	R0	SP	PS	PC	SEL	controller

Select Words

- Selw. 0-3 : 7 EMT codes, one 0 as endmarker
- Selw. 4-10: 7 TRAP codes, one 0 as endmarker
- Selw. 11-20: 4 pairs of the form: Reg. number
Reg. value

Table A.5: Event Specifications

<u>ID</u>	<u>EXPLANATION</u>
PC	Program Counter
PS	Processor Status word
SP	Stack Pointer (=R7)
CP	Content of Program counter (=op. code for 1 wrd instr.)
EW	Error Word (=content of #52) for EMT calls
SR	Device Status Register
BF	Device Buffer Content
Rn	Register n of disk controller
HT	Host Time (content of host clock buffer)
ST	Supervisor Time (in sec. relative to start of session)

Table A.6: Value Identification for Numeric Display

<u>CMD</u>	<u>SUBPROG</u>	<u>EXPLANATION</u>
W	WAITUP	Terminal output is disabled
P	PROC	Terminal output is enabled after W
E	SPEVUP	Special event table update mode
C	CONT	Continue after special event occurrence
B	INPICI	Initialize histogram, draw coordinates
A	DRPICI	Erase Screen and draw histogram as it was
D	CNTOUT	Erase screen and print event summary report
S	CSIS	CMD interpreter for S-MON mode definition
H	CSIH	CMD interpreter for H-MON mode def. (not impl.)
L	CSIL	CMD interpreter for logging mode definition
T	TERM	Terminate session and print summary report

Table A.7: Control System Commands

SOFTWARE DOCUMENTATION

1. Software Monitor (S-MON)

S-MON: INTERACTIVE SOFTWARE MON RT-11 MACRO VMO2-09 0011126 PAGE 1

```

1      .TITLE S-MON: INTERACTIVE SOFTWARE MONITOR
2      *****
3
4
5
6
7      * THIS PROGRAM IS PART OF THE MONITOR SYSTEM "MIMOS".
8      * THE FIRST PART IS NONRESIDENT AND ONLY USED FOR INITIALISATION.
9      * THE SECOND PART IS RELOCATED ABOVE RT11 STARTING AT 134000
10     * FROM WHERE SYSTEM ACTIVITIES ARE MONITORED.
11     * S-MON IS A PURE INTERCEPT TYPE MONITOR WHICH CAN ANALYZE
12     * ANY VECTORED EVENT.
13     * PART ONE STARTS AT "START"
14     * PART TWO STARTS AT "STRES"
15     *****
16
17     .MCALL .PRINT,.REGDEF,.EXIT,.OUTDEF,.DEF,.SEND,.RECV,.DEF
18     .MCALL .TTYIN,.TTYOUT
19     .DEF
20     .REGDEF
21     .ASECT
22     .=1000          FLOAD POINT
23     .TITLE CONSTANT DEFINITIONS
24     *****
25
26     M=1              M+1 GROUPS
27     MAXMM=M*4
28     N=16,*(M+1)      N = TOTAL NUMBER OF TESTED VECTORS
29     EDT=4
30     ACK=6
31     B=7*340
32     PSH=17776
33     TTYOUT=17756     TTY OUTPUT STATUS
34     TBUF=17542        TTY BUFFER
35     TSTAT=17540       TTY STATUS REGISTER
36     PLIV=320          VECTOR FOR PARALLEL LINK
37     PLSH=16400        STATUS REG. "-"
38     PLOB=16402        TTYOUT BUF. "-"
39     PLIB=16404        TTYIN BUF. "-"
40     IIS=PLSH          PREDEFINITIONS FOR PAR. LINK
41     IOS=PLSH
42     IIB=PLIB
43     IOB=PLOB
44     IIV=PLIV
45     IOV=PLIV+4
46     CONCH=IIS         STATUS REG. OF COMM. CHANNEL IN
47     CONCH=IOS
48     REGIN=13400        START ADDR. OF RES. ROUTINE
49     FWSK=REGIN        FOR DURING MONITORING
50     DIST=REGIN+STRES  RELOCATION FACTOR
51
52     * RELOCATED VARIABLES
53
54     XBLANK=BLANK+DIST
55     XAST=AST+DIST
56     XCLF=CLF+DIST
57     XREG=REG+DIST     DELINEATER FOR PRINT
58     XSAVE=LUC FOR R0

```

CONSTANT DEFINITIONS.

RT-11 MACRO VM02-09

00111126 PAGE 1+

58	140024	XREG5=REG5+DIST	ISAVE LOC FOR MS
59	137046	XTXTBF=XTXTBUF+DIST	
60	134072	XENTR0=ENTR0+DIST	IENTRY ADDR. FOR GROUP 0
61	134074	XENTR1=ENTR1+DIST	IENTRY ADDR. FOR GROUP 1
62	134006	XENTR2=ENTR2+DIST	IENTRY ADDR. FOR SUPERVISOR
63	140026	XPSSAV=PSW+SAV+DIST	ITD SAVE I NUMBER
64	137012	XVECSV=VECSV+DIST	ISAVE TABLE FOR VECTORS
65	137012	XVECT0=VECT0+DIST	IVECTOR TABLE
66	137170	XMSK0=MASK0+DIST	IFIRST MASK WORD
67	136300	XREPL=REPLACE+DIST	IREPLACEMENT PROCEDURE
68	140030	XDIR=DIR+DIST	ITD DIR, TO SAVC, ISSAY TO ORIG
69	137030	XQUES=QUES+DIST	IT
70	137032	XMPQ=MPQ+DIST	IT
71	137030	XMP5=MP5+DIST	IT
72	137036	XMSF=MSF+DIST	IT
73	137040	XMPQ=MPQ+DIST	IT
74	137042	XPS2=PS2+DIST	IT
75	137220	XSG00=SG00+DIST	ISELECT WORD
76	134742	XDE=DE+DIST	IRELOC. HANDLER ADDRESSES
77	134760	XSG00=SG00+DIST	
78	135070	XSG01=SG01+DIST	
79	135300	XSG10=SG10+DIST	
80	135050	XSG11=SG11+DIST	
81	135760	XSG12=SG12+DIST	
82	134430	XTAB0P=TAB0P+DIST	
83	130010	XHANG=HANG+DIST	
84	137300	XDSPT0=DISPT0+DIST	IMAIN DISPATCH TABLE
85	137400	XDSPT2=DISPT2+DIST	ISUPERVISOR CMD DISPATCH
86	137200	XG100=SL100+DIST	
87	137200	XMS101=MSK101+DIST	
88	137712	XUSFAT=UVSTAT+DIST	
89	140056	XNTI=NTI+DIST	
90	137200	XMS05=MSK05+DIST	
91	134700	XENDR0=ENDR0+DIST	
92	137302	XENDCM=ENDCM+DIST	
93	130026	XCONT=CONT+DIST	
94	137266	XTIMON=TIMEON+DIST	
95	137276	XINT1=INT1+DIST	
96	137300	XINT2=INT2+DIST	
97	137272	XIREG1=IREG1+DIST	
98	137274	XIREG2=IREG2+DIST	
99			
100			
101		.TITLE MACRO DEFINITIONS	
102		*****	
103			
104			
105		.MACRO ,OCOUT	
106		ISX PC,OPRINT	
107		ISX PC,UCPUT	
108		.ENDM	
109			
110		.MACRO ,CHOUT	
111		ISX PC,PRINT	
112		ISX PC,CHPUT	
113		.ENDM	
114			

MACRO DEFINITIONS

RT-11 MACRO VM02-09

00111126 PAGE 1*

115
116
117
118
119
120
121
122
123
124
125
126
127
128
129
130
131
132
133.MACRO .SENDNRD
TST #PLSR
BMI ,=4
MOV R0,#PLOB
.ENDM;TRANSMIT ONE WORD TO SUPPLY.
;WAS PREVIOUS WORD RECEIVED?
;WAIT IF NOT
;SEND WORD IN R0

.MACRO .RCVNRD

;RECEIVE ONE WORD FROM SUPPLY.

TST #PLSR
BPL ,=4
MOV #PLIB,R0
.ENDM;IS THERE A NEW WORD
;WAIT IF NOT
;STORE WORD IN R0.MACRO .BUFOUT
JSR PC,PRINTBF
JSR PC,SENDNF
BIT #2,H4
.ENDM

;TRANSMISSION OF ONE EVENT RECORD

;ENDMARKER SFTZ

```

1
2
3
4
5 001000
6
7 001000
8 001000 105737 177564
9 001000 100375
10 001010 000207
11
12
13 001012
14
15
16 001012 010300
17 001014 005001
18 001016 005003
19 001022 005006
20 001022
21 001026 002700 000200
22 001032 122700 000012
23 001036 001402
24 001040 010006
25 001042 000767
26 001044 005726
27 001046 005716
28 001050 001416
29 001052 122716 000120
30 001056 001410
31 001060 102716 000060
32 001064 011000
33 001066 072003
34 001070 000001
35 001072 002703 000003
36 001076 000762
37 001080 012700 177777
38 001084 000757
39 001086 005703
40 001090 001002
41 001092 052700 177775
42 001096 005726
43 001100 012603
44 001102 000207
45
46
47 001120
48
49
50 001120 010106
51 001126 010206
52 001130 002700 177000
53 001134 012701 000003
54 001140 000000
55
56 001142
57

```

```

TITLE I/O SUBROUTINES - NON RESIDENT
*****
I
I*****
PRINT: IPRINT TEXT STARTING AT (R0)
I*****
I PRINT
PRINT3: TSTB #11100T
RPL PRINT3
RTS PC IRETURN FROM PRINT
I
I*****
OCRD: IHEAD OCTAL NUMBER; LINKAGE: PC
I***** IREQ=3 IF T, -2 IF EMPTY; RESULT IN R1
I
MOV R1, -(SP)
CLR R1 IRESULT BUFFER
CLR R1 ISHIFT COUNTER
CLR -(SP) ISTART MARKER ON TOP
OCRD1: TTYIN
BIC #200, R0
CMPB #12, R0
BEQ OCRD5
MOV R0, -(SP) ISTACK NUMBER
BR OCRD1
OCRD3: TST (SP)+
TST (SP) ITOP OF STACK?
BEQ OCRD5 IYES
CMP #124, (SP) IT?
BEQ OCRD4
SUB #60, (SP) ICONVERT
MOV (SP), R0
ASH R0, R0 ISHIFT NR N TIMES LEFT
BIS R0, R1
ADD R1, R1 IINCR SHIFT CNT
BR OCRD5
OCRD4: MOV #-1, R0 I-3 IF T
BR OCRD5
OCRD5: TST R1
BNE OCRD6
BIS 0177775, R0 I-2 IF EMPTY
TST (SP)+
MOV (SP)+, R3
RTS PC IRETURN FROM OCIN
I
I*****
PRINT3: IPRINT BYTE OCTAL; LINKAGE: PC
I***** IVALUE IN R0
I
MOV R1, -(SP)
MOV R2, -(SP)
BIC #177420, R0 ICLEAR UPPER BYTE
MOV R3, R1 IPRINT 3 DIGITS
BR IYH1
I*****
PRINT3: IPRINT WORD OCTAL; LINKAGE: PC
I***** IVALUE IN R0

```

```

59      MOV R1,-(SP)
60      MOV R2,-(SP)
61      MOV #6,R1
62      MOV R1,R2
63      MOV R0,-(SP)
64      BIC #17770,(SP)
65      CLC
66      ROR R0
67      ASR R0
68      ASR R0
69      DEC R1
70      BNE TYP2
71      MOV (SP)+,R0
72      DEC R2
73      BEO TYP4
74      TST R1
75      BIE TYP4
76      TST R0
77      BEQ TYP5
78      ADD #60,R0
79      ,TTTOUT
80      WAIT
81      INC R1
82      TST R2
83      BNE TYP3
84      MOV (SP)+,R2
85      MOV (SP)+,R1
86      RTS PC
87
88      *****
89      BINARY: JHEAD WORD BINARY; LIMAGE: PC
90      *****
91      J.
92      MOV R3,-(SP)
93      MOV R2,-(SP)
94      CLR R2
95      CLR R3
96      BINARY: ,TTTIN
97      BIC #200,R0
98      CMPB #14,R0
99      BEQ BINARY
100      INC R3
101      ASL R2
102      CMPB #64,R0
103      BEQ BINARY
104      HIS #1,R2
105      BR BINARY
106      BINARY: TST R3
107      BEQ BINARY
108      MOV R2,R1
109      BINARY: ,TTTIN
110      MOV (SP)+,R2
111      MOV (SP)+,R3
112      RTS PC
113
114      *****

```

PRINT 6 DIGITS
ISAVE COUNT

LAST DIGIT?
YES?
ANY TYPED YET?
YES
NO, LEADING 0?
YES, SAT

NUMBER?
NO

RETURN FROM PRINT/WORD

EMPTY WORD
STORE NEW WORD
EMPTY WITH BUFFER

RETURN FROM BINARY

140

I/O SUBROUTINES - NON RESIDENT RT-11 MACRO VM02-59 00:11:26 PAGE 2*

```

115 001330      BINPUT: PRINT WORD BINARY; LINKAGE: PC
116          *****
117          MOV R1, (SP)
118          MOV #BITPOS, R0
119          JSR PC, PRINT
120          MOV (1), R4
121          MOV #BITS, R1
122          BIT #1000000, R4
123          BEO BINPUT2
124          MOV #8, (1)+
125          BR BINPUT3
126          BINPUT2: MOV #80, (1)+
127          BINPUT3: ASL R1
128          CBR R1, #BITS+16
129          BAE BINPUT1
130          MOV #BITS, R0
131          JSR PC, PRINT
132          MOV (SP)+, R1
133          RTS PC
134          *****
135          START: MOV #1000, SP
136          TST #0
137          *****
138          *****

```

RETURN FROM BINPUT

***** ENTRY POINT *****
***** INPUT BUFFER *****

SPECIFICATION OF S-MON MODE RT-11 MACRO VME2-09 02:11:26 PAGE 3

```

1  ;
2  ; TITLE SPECIFICATION OF S-MON MODE
3  ; *****
4  ;
5  ; COMMAND INTERPRETER
6  ; *****
7  ;
8  ; CLR R5 ;SET LOOP SWITCH
9  ;
10 ;
11 ; CHOIR: ;***MASK UPDATE***
12 ;
13 ; MOV AMSK,R5
14 ; JSM PC,PRINT
15 ;
16 ; CM011: ;*** V/N INTERPRETER ***
17 ;
18 ; ;TTYN
19 ; BIC #200,R5
20 ; CMPB #15,R5
21 ; BLO CM012
22 ; MOV R0,R1
23 ; RR CM011
24 ; ;TTYN
25 ; CM012: CMPB #114,R1
26 ; CM013: BNE CM014
27 ; JMP CM013(5)
28 ; CM014: JMB #131,R1
29 ; BLO CM015
30 ; MOV AC,USER,R0
31 ; JSM PC,PRINT
32 ; JMB #172,R1
33 ; CM015: ;*** INDEX INPUT ***
34 ;
35 ; MOV #REQ,R0
36 ; JSM PC,PRINT
37 ; MOV #CM015,R0
38 ; JSM PC,PRINT
39 ; JST R0
40 ; RLT CM017
41 ; JMB #CM015(5)
42 ; CM017: JMB #CM013(5)
43 ;
44 ; CM018: ;*** MASK I/O ***
45 ;
46 ; ;INDEFXX
47 ; ASL R1
48 ; AND #MASK,R1
49 ; JSM PC,PRINT
50 ; JMB #MASK
51 ; JSM PC,BIND
52 ; RR CM010
53 ; CM019: ;*** SUBMASK I/O ***
54 ;
55 ; ;INDEFXX
56 ; ASL R1
57 ; AND #MASK,R1
58 ; JSM PC,PRINT
59 ; JMB #SUBMASK

```

SPECIFICATION OF S-MON MODE

RT-11 MACRO VM02-00

0011126 PAGE 3+

```

58 001610 004767 177514
59 001614 004767 177422
60 001622 004743
61
62 001622
63
64 001622 005725
65 001624 012700 002565
66 001634 004767 177144
67 001634 004767 177576
68
69 001640
70
71 001644 005725
72 001646 012700 002621
73 001646 004767 177120
74 001654 002167 177568
75
76 001656
77
78 001650 006301
79 001660 001022
80 001662 002701 000020
81 001666 002424
82 001670 012700 000002
83
84 001674
85
86 001674 005000
87 001676 110200 006300
88 001702 004767 177216
89 001706 012700 002665
90 001712 004767 177062
91 001716 004767 177070
92 001722 005100
93 001724 100402
94 001726 110162 006300
95 001732 002202
96 001736 007321
97 001736 000674
98
99 001740
100
101 001740 010200 006300
102 001744 004767 177172
103 001750 012700 002665
104 001754 004767 177020
105 001760 004767 177026
106 001764 005100
107 001766 100600
108 001770 010162 006322
109 001774 000675
110
111 001776
112
113 001776 005267 004342
114 002002 005267 004342

```

```

JSR PC,BINPRT      ;PRINT MASK
JSR PC,BINRU        ;READ NEW MASK
RM CMB16            ;FIND NEXT INDEX
;
CMD110: ;*** SUBMASK UPDATE ***
;
TST (S)+            ;INCR. LOOP SWITCH
CMD111: MOV #SMASKIN,R0
JSR PC,PRINT
JMP CMD11            ;READ Y/N
;
CMD112: ;*** SELECT WORD UPDATE ***
;
TST (S)+            ;INCR. LOOP SWITCH
CMD113: MOV #SELWIN,R0
JSR PC,PRINT
JMP CMD11            ;READ Y/N
;
CMD114: ;*** SELECT WORD I/O ***
;
ASL R1              ;INDEX=2
MOV R1,R2            ;SAVE INDEX
CMP #20,R1           ;INDEX>12?
BLT CMD116           ;YES-LOAD WORD I/O
MOV R2,R3            ;PROCESS 2 BYTES
;
CMD115: ;*** BYTE I/O ***
;
CLR R0
MOV SLDR60(2),R0     ;GET OLD BYTE VALUE
JSR PC,PRINTBYT      ;PRINT IT
MOV #GAP,R0
JSR PC,PRINT          ;PRINT 1
JSR PC,GETCH          ;GET NEW VALUE
TST R0
BHI CMD15A            ;NO NEW VALUE
MOV R1,SLDR60(2)      ;STORE IT
CMD15A: INC R2
SUB R3,CMD115         ;SECOND BYTE
RM CMB16
;
CMD116: ;*** WORD I/O ***
;
MOV SLDR60(2),R0     ;GET OLD WORD VALUE
JSR PC,PRINTWD       ;PRINT IT
MOV #GAP,R0
JSR PC,PRINT          ;PRINT 1
JSR PC,GETCH          ;GET NEW VALUE
TST R0
BHI CMD16            ;NO NEW VALUE
MOV R1,SLDR60(2)      ;STORE IT
RM CMB16
;
CMD117: ;*** TIMER CONTROL ***
;
CLR TIM0
CLR TIM0H            ;NO TIME LOGGING
;

```

SPECIFICATION OF S-MON MODE

RT-11 MACRO VM02-09

00:11:26 PAGE 34

```

115 002026 005067 004334 CLR TREG1
116 002012 002767 000020 004230 BIT X20,MASK1 TIME INTERRUPT?
117 002020 001420 SEQ CM0119 MASK-ASK FOR TIMING
118
119
120
121 002022 012700 002701 MOV #TIMST1,R0 GET FREQUENCY
122 002020 004767 176706 JSR PC,PRINT
123 002032 004767 176754 JSR PC,UCMD READ 0-3
124 002030 000321 ASL R1
125 002043 002701 000010 B1S #10,R1 STOP REP. MODE DOWN
126 002044 002167 000276 B1S #1,TREG1
127 002050 005267 000266 INC TIMEON 1000. INTERM. AFTER EACH
128 002050 002700 002777 CM0118: MOV #TIMCNT,R0 GET COUNT
129 002060 004767 176714 JSR PC,PRINT
130 002064 004767 176722 JSR PC,UCMD
131 002070 0014167 000254 MOV R1,TREG2 SET CNT
132 002074 002421 BR CM0120 TEMP OF INTR. TIMING
133
134 002070 CM0119: *** TIME LOG MODE ***
135
136 002076 012700 002731 MOV #TIMST2,R0 GET FREQUENCY
137 002102 004767 176672 JSR PC,PRINT
138 002100 004767 176720 JSR PC,UCMD
139 002112 002701 000005 CMP #3,R1
140 002110 100410 BMT CM0120
141 002120 000321 ASL R1
142 002122 002701 000031 B1S #31,R1
143 002120 000167 000214 B1S #1,TREG2 START REP. MODE UP
144 002132 005267 000206 INC TIMELOG SET TIME LOG MODE
145 002130 002706 BR CM0118 LOG, GET CNT
146
147 002140 CM0120: *** TRAP INIT ***
148
149 002140 012737 140036 000034 MOV #VRTII,0034
150 002140 005037 000036 CLR #036 SET HTI ADDR. IN VECTOR
151
152 002152 CM0121: *** VECTOR PROTECTION ***
153
154 002152 013700 000054 MOV #054,X0
155 002150 002760 177777 000306 R1S #17777,524(2)
156 002160 002760 177777 000332 R1S #17777,532(2)
157 002170 002760 177777 000302 R1S #17777,540(2)
158 002200 012737 000240 146762 MOV #240,00146762
159
160

```

TITLE INITIATION PROGRAM

```

TITLE INITIALISATION PROGRAM
*****

MOV #10000,R1
SUB #8000,R1
MOV #1000,R1
JSR PC,PRINT
MOV R1,R0
JSR PC,PRTEND

; POWERL: CLR #24          ;SET POW.FAIL VECT
                ;SET SUPER INTERR. VECT.
MOV #FINT2,#11V
MOV #BETI,#11V+2
HIS #1000,#11S
TST TIMEON
BEQ INIT1
MOV TIME1,#104
MOV TIME2,#106
MOV #ATIME,F,TIMPTR

INITIAL1: MOV #BYTES,R#
PRELOC: MOV #EQUIN,M1
        MOV #EQUIN,R2
ALIGN2: MOV (#M1),R1)+
        RME ACN2
        CLR FADIN
        JSR PC,EXCH
        MOV TREGI,#10UF
        MOV HREGI,#1STAT
        RLC #CONCHO
        HIS #CONCHO
        NOP
        NOP
        NOP
        MOV #AMP
        TST TIMEON
        BEQ FIN
        HIS #INI,#1STAT
        EXIT
FIN:

```

NON RESIDENT DATA SECTION

```

.MLIST HEX
TXT1: .ASCII/JFFSET /
      .BYTE 200
TXT2: .BYTE 15,12
      .ASCII/HIGH LIMIT /
      .BYTE 200
BITPOS: .ASCIZ 17654321876543210/
BITS: .ASCIZ /ACCMPCACCMPCCMPCC/
MSKIN: .ASCII <1><1>/M/S/A UPDAIF {Y,N}? : /<20>
INCR: .ASCII /GIVE INDEFA CM 1/<60>
SMKIN: .ASCII <1><1>/SUMMAS UPDATE {Y,N}? : /<20>

```

INITIALISATION PROGRAM RT-11 MACRO VM92-09 0011126 PAGE 4

```

50 002621 015 SELWIN: .ASCII <15><12>/SELECT AOWD UPDATE (Y,N)? 1 /<278>
51 002601 012 CROSS: .ASCII <15><12>/<284>
52 002605 008 GAP: .ASCII / : /<20>
53 002671 077 114 CHDERR: .ASCII /TILL CHD?/
54 002741 015 111 TIMST1: .ASCII <15><12>/INTERM. FREQUENCY /<288>
55 002751 012 124 TIMST2: .ASCII <15><12>/INTERM FREQUENCY (3: NO TIME LOG): /<278>
56 002771 124 115 TIMCNT: .ASCII /TIMPER COUNT : /<280>
57 003024 001022 EVEN
58 003026 001040 CHD11: CHD110
59 003030 001050 CHD12: CHD112
60 003034 001070 CHD17: CHD117
61 003038 001080 CHD18: CHD12:
62 003042 001090 CHD19: CHD11:
63 003046 001100 CHD113: CHD13:
64 003050 001110 CHD110: CHD110
65 003054 001120 CHD112: CHD112
66 003058 001130 CHD117: CHD117
67 003062 001140 CHD19: CHD19
68 003066 001150 CHD110: CHD110
69 003070 001160 CHD110: CHD110
70 003074 001170 CHD110: CHD110
71 003078 001180 CHD110: CHD110
72 003082 001190 CHD110: CHD110
73 003086 001200 CHD110: CHD110
74 003090 001210 CHD110: CHD110
75 003094 001220 CHD110: CHD110
76 003098 001230 CHD110: CHD110
77 003102 001240 CHD110: CHD110
78 003106 001250 CHD110: CHD110
79 003110 001260 CHD110: CHD110
80 003114 001270 CHD110: CHD110
81 003118 001280 CHD110: CHD110

```

```

.LIST BEX
.NLIST ME

```

RESIDENT ROUTINES

RT-11 MACRO VM02-09

0011126 PAGE 5

```

1  .TITLE RESIDENT ROUTINES
2  *****
3  ;
4  STRESS: NOP ;START CF RESIDENT ROUTINE
5  ;
6  ;
7  ;INTERCEPT ENTRY FOR GROUP 0
8  ;
9  ;
10 12 003056 013737 17776 140026 ENTRY0: MOV RPSW,PCXPSAV
11 003064 010767 004020 MOV PC,PC1A00
12 003170 000026 BR SAVE
13 003372 012720 000000 MOV R3,R0
14 003070 000002 BR DISP
15
16
17
18
19 20 003100 013737 17776 140026 ENTRY1: MOV RPSW,PCXPSAV
21 003106 010767 003376 MOV PC,PC1A00
22 003112 000415 BR SAVE
23 003114 012720 000001 MOV R1,R0
24 003120 000031 BR DISP
25
26
27
28
29
30 003122 013737 17776 140026 ENTRY2: MOV RPSW,PCXPSAV
31 003130 002737 000120 BIC R12,R0,PCXCHT
32 003136 010767 003376 MOV PC,PC1A00
33 003142 000001 BR SAVE
34 003140 000020 BR EXDISP
35
36
37
38
39 003140 000200 SAVE: NOP
40 003150 010507 MOV R5,REGS
41 003154 012725 MOV R0,REGS,R5
42 003160 012725 MOV R0,R5)++
43 003164 010225 MOV R1,(R5)++
44 003166 010225 MOV R2,(R5)++
45 003170 010225 MOV R3,(R5)++
46 003172 002767 ADD R0,PC1A00
47 003170 000177 JMP PC1A00
48

```

```

;GO BEYOND OR INSTR
;SIMULATE RTS

```

```

;SAVE GPR'S

```

```

;IGNOR DISP. FOR EXT EV.

```

```

;SAVE PSW
;DISABLE SUP. INTER.
;NO SIMULATE JSR

```

```

;SAVE PSW WITH I-LOOP
;NO SIMULATE JSR
;INDICATE GROUP INDEX

```

```

;INTERCEPT ENTRY FOR GROUP 1

```

```

;SAVE PSW WITH I-LOOP
;NO SIMULATE JSR
;INDICATE GROUP INDEX

```

EVENT HANDLER DISPATCH RT-11 MACRO VM02-09 0011126 PAGE 0

```

1
2
3
4 003204 000240
5 003208 005004
6
7
8 003210 010601
9 003212 002701 100000
10 003216 000127 000500
11 003222 100006
12 003224 052704 000001
13 003230 012701 001000
14 003234 010106
15 003236 005011
16 003240 010607 003602
17 003244 012706 130000
18
19
20 003250 072027 000004
21 003254 002737 177760 100026
22 003262 003700 100026
23 003266 010005
24 003270 000240
25 003272 010037 177570
26
27
28 003276 006300
29 003300 000770 137304
30 003304
31
32
33 003314 000240
34 003316 006305
35 003320 000305
36 003322 010500
37 003324 016701 003556
38 003330 102701 000000
39 003334 010167 003546
40 003340 016021 137412
41 003344 005720
42 003346 010011 137412
43 003352 000767 000350
44 003356 016706 005524
45 003362 000240
46 003364 105767 002530
47 003370 001405
48 003372 105067 002522
49 003376 012704 003376
50 003402 100350
51 003406 000002
52

```

```

; TITLE EVENT HANDLER DISPATCH
; *****
DISP: NOP
CLR R0 ; FLAG REGISTER

;
MOV SP,R1
BIC #100000,R1 ; CLEAR SIGN BIT
CMP R1,#0000 ; EVENT STK <= 0000
RPL SVSTK ; NO-PROTECT VITAL VECTORS
BIS #1,R0 ; FLAG SP MODIFICATION
MOV #1000,R1 ; YES - SET IT TO 1000
MOV R1,SP ; SAVE IT AS OLDSK
CLR (R1) ; SIMULATE EVENT PC=0
SVSTK: MOV SP,OLDSK ; SAVE STACK
MOV #NEWSK,SP ; INST NEW STACK

;
ASH #4,R0 ; SHIFT GROUP NO. J 4 LEFT
BIC #177760,#XPSSAV ; MASK I-NUMBER
ADD #XPSSAV,R0 ; COMPUTE EVENT CODE=16*J+I
MOV R0,R5 ; SAVE R0 DURING I/O
NOP
MOV R0,#177570 ; TEST

;
ASL R0 ; MULTI EV CODE BY 2
JSR PC,#XOSPT0(0) ; CALL SUBGROUP HANDLER
; BUFOUT

;
RESTOR: NOP
ASL R5
ASL R5
MOV R5,R0 ; 16* EVENT CODE
MOV OLDSK,R1 ; PUSH STACK AND
SUB #4,R1 ; PUSH FOR ADDITIONAL VECTOR
MOV R1,OLDSK ; SAVE STACK AND FOR RTI
MOV XVELEV(R0),(R1)+ ; MOVE SAVED VECTOR ON STACK
TST (R1)+
MOV XVELEV(R0),(R1)
JSR PC,UNSAVE
MOV OLDSK,SP ; INFT SP
NOP
TST0 TRAP
REQ RESTRI
CLR0 TRAP
MOV #,R0 ; SET R0=0 FOR EXIT
ENT 300
RESTRI: RTI ; RETURN TO SYSTEM
;

```

```

1
2
3
4
5 003400
6
7
8 003400 005004
9 003410 010067 003472
10 003410 012706 130000
11 003420
12 003430 012737 000777 177570
13 003440 004770 137400
14 003440 004767 000256
15 003450 016706 003432
16 003450 002737 000100 164000
17 003460 000002
18
19
20 003460
21
22
23 003460
24 003470 004770 137400
25 003500 000207
26
27
28 003504
29
30 003500 012737 000007 177570
31 003510 005237 140030
32 003510 004767 001552
33 003520 012701 137170
34 003520 012702 137300
35 003530 012100
36 003534
37 003540 000102
38 003550 001370
39 003550 012737 000070 177570
40 003560
41 003570 002700 000006
42 003570 001401
43 003600 000000
44 003600 012701 137170
45 003600 012702 137300
46 003610 012737 000700 177570
47 003620
48 003630 010021
49 003640 000102
50 003650 001370
51 003660 012737 000700 177570
52 003660 002720 000003
53 003650 001401
54 003650 000000
55 003660 005057 140030
56 003660 005737 137266
57 003660 001400

```

.TITLE CONTROL COMMAND HANDLER

 EXDISP: DISPATCHER FOR SUPERVISOR COMMANDS

 CLR R4
 MOV SP,OLDSTK ;SAVE STACK
 MOV #NEWSTK,SP ;SET NEW STACK
 ;RCVNRD ;GET FUNCTION CODE INTO R0
 MOV #777, #177570 ;TEST
 JSR PC, #XDSPT2(0) ;DISPATCH SUPERVISOR COMMAND
 JSR PC, UNSAVE
 MOV OLDSTK, SP
 BIS #100, #CONCHI ;ENABLE SUP. INTERN.
 RTI ;RETURN TO SYSTEM

 HANG: WAITLOOP FOR FURTHER SUP. CMDS

 ;RCVNRD ;WAIT FOR A CMD
 JSR PC, #XDSPT2(6) ;DISPATCH NEXT CMD
 RTS PC ;RETURN FROM HANG, REL. HOST

 TABUPD: UPDATE OF MONITOR TABLES

 MOV #7, #177570 ;TEST
 INC #XDIR ;SET EXCHANGE BACK+ROS
 JSR PC, EXCH ;INSET ORIGINAL VECTORS
 MOV #XDSK0, R1 ;FROM OF SPEC TABLE
 MOV #XENDCM+2, R2 ;POSITION OF "
 TABUP1: MOV (R1), #X ;SEND SPEC TO SUPER
 ;RCVNRD
 CMP R1, #0 ;FINISHED?
 BNE TABUP1 ;NO
 MOV #7, #177570 ;TEST
 ;RCVNRD ;WAIT FOR ACK
 CMP #0, #0 ;WHEN THIS IS 1, THE NEW
 ;BND TABUP2 ;TABLE IS DEFINED AT SUPER
 HALT ;PROTOCOL ERROR
 TABUP2: MOV #XDSK0, R1
 MOV #XENDCM+2, R2
 MOV #7, #177570 ;TEST
 ;RCVNRD ;RECEIVE NEW TABLE
 MOV #, (R1) ;FROM SUPERVISOR
 CMP R1, #0 ;FINISHED?
 BNE TABUP3 ;NO
 MOV #7, #177570 ;TEST
 CMP #0, #0 ;HAS LAST EXCH
 REQ TABUP4
 HALT ;PROTOCOL ERROR
 TABUP4: CLR #XDIR ;EXCHANGE VECTORS ACCORDING
 TST #XDIR ;TO INTENT. MODE?
 BND TABUP5 ;NO

```

58 003670 013737 137276 000104      MOV #XTINT1,#104      YES-SET DUMMY VECTOR
59 003676 013737 137300 000106      MOV #XTINT2,#106
60 003700 004767 001364      TABUPS: JSR PC,EXCH      YES-SET DUMMY VECTOR
61 003710 013737 137274 172542      MOV #XTNEG2,#ATBUF    YES-SET DUMMY VECTOR
62 003716 013737 137272 172540      MOV #XTNEG1,#ATSTAT   YES-SET DUMMY VECTOR
63 003724 002207      RIS PC
64
65
66
67 003726 002207      UNSAVE: NOP
68 003730 012705 140012      MOV #XREG4,R5      PRESTORE GPR'S
69 003734 012500      MOV (R5)+,R0
70 003736 012501      MOV (R5)+,R1
71 003740 012502      MOV (R5)+,R2
72 003742 012503      MOV (R5)+,R3
73 003744 012504      MOV (R5)+,R4
74 003746 016705 003126      MOV REG5,R5
75 003752 002207      RIS PC
76
77
78 003754      *****
79      ENDRUN: PRESET SYSTEM; TERMINATE MONITORING SESSION
80      *****
81
82
83 003754 012706 134000      MOV #NEWSTK,SP
84 003760 012737 000001 140010      MOV #1,#XDIW      SAVE TO ORIGINAL
85 003766 004767 001302      JSR PC,EXCH      FOR THE EXCHANGE
86 003772 005037 172540      CLR #ATSTAT      CLEAR IIFW
87 003776 005037 164000      CLR #CCHCHI      CLEAR CHANNEL
88 003782 004767 177720      JSR PC,UNSAVE
89 003786 016706 003074      MOV OLDSTA,SP
90
91      CLR R0      FINIT SYSTEM AT .EXIT
92      .EXIT
93

```

```

      .TITLE EVENT HANDLERS
      *****
;
;
;*****
DE:    DISPATCH ERROR, WRONG EV CODE
;*****

      JSR PC,ECOUT          ISEND EVENT CODE
      MOV #QUES,R0         ISEND QUESTION MARK
      .CMOUT
      RTS PC                IRETURN FROM DE
;
;*****
SG00:  ?EV1-EV5: TRAP 4,TRAP 12, HPT, IOI, P&PL
;*****
;
      JSR PC,ECOUT          ISEND EVENT CODE
      MOV MSK005,R1        IGET SPEC MASK
      MOV #1,R2            ISTARTING WITH EV 1
SG001: CMP R2,R5           ILOGA FOR EV CODE
      BEQ SG002
      ASH #-3,R1           ISHIFT 3 RIGHT
      INC R2              IIO PLACE THE SPEC BITS FOR
                        ITHE EVENT RIGHT BOUNDFU
SG002: BIT #1,R1           IPC?
      BEQ SG003           INO
      JSR R4,PCOUT
      .WORK 0
SG003: BIT #2,R1           IPS?
      BEQ SG004           INO
      JSR PC,PSOUT
      BIT #4,R1           ISP?
      BEQ SG005           INO
      JSR PC,SPOUT
      JSR PC,TIHOOT
SG005: CMP R5,R3           Iwas IT EV.1 OR 2?
      BEQ SG006           INO
      INCB TRAP           ISET FLAG
SG006: RTS PC             IRETURN FROM SG00
;
;*****
SG01:  ?EV6-EV7 : EMT, TRAP - WITH SUB SELECTION
;*****
;
      MOV MSK007,R1        IGET MASK
      MOV #X5000,R2        IGET ADDR. OF 1ST SEL BYTE
      CMP #0,R5            ITERMINE IF 0 OR 7
      BEQ SG011
      S&AB R1              IORING EV7 MASK IN LOW BYTE
      ADD #0,,R2           I11 R SEL BYTES PER EV 1
SG011: BIT #1,R1           ISELECTION?
      BNE SG012
      JSR PC,SG01A
      RTS PC              IEND -GOTO REPORT PROCESSING
;
;
SG012: ?SUBSELECTION
;
      MOV SOLUSIX,R5        IGET EMT/TRAP CODE

```

```

58 004210 005743      TST -(3)
59 004212 011303      MOV (3),R3
60 004214 000240      NOP
61 004216 124312      CMPEB R3,(2)      ;EV INSTR.=SEL INSTR.?
62 004220 001003      BNE SG010      JND
63 004222 004767 000012 JSR PC,SG01A      ;YES-GOTO REPORT PROCESSING
64 004226 000207      RTS PC      ;RETURN FROM SG01
65 004230 005202      INC R2      ;GET NEXT SEL BYTE
66 004232 105712      TSTB (2)      ;SELECTION?
67 004234 001367      BNE SG011      ;YES-GOTO SUBSECTOR
68 004236 000207      RTS PC      ;RETURN FROM SG01
69
70 004240
71
72 004242 004767 000602 JSR PC,ECOUT      ;SEND EVENT CODE
73 004244 032701 000002 BIT #2,R1      IMC?
74 004246 001403      BEO SG01A1      JND
75 004250 004467 000610 JSR R4,PCOUT      ;FROM PC-2
76 004252 000002      .WORD 2      ;SP?
77 004254 032701 000004 SG01A1: BIT #0,R1      JND
78 004256 001402      REQ SG01A2      ;PC?
79 004260 004767 000620 JSR PC,PCOUT      JND
80 004262 032701 000010 SG01A2: BIT #10,R1      ;SP?
81 004264 001402      BEO SG01A3      JND
82 004266 004767 000630 JSR PC,PCOUT
83 004268 032701 000020 SG01A3: BIT #20,R1      ;PC-2?
84 004270 001403      REQ SG01A4      JND
85 004272 004467 000640 JSR R4,CPCOUT      ;FROM PC-2
86 004274 000002      .WORD 2      ;SP?
87 004276 032701 000040 SG01A4: BIT #40,R1      ;PC-2?
88 004278 001410      REQ SG01A5      JND
89 004280 012740 137042 MOV #XMS2,R0      ;FS-SEND (S2)
90 004282 004332      .CNDUT
91 004284 013700 000052 MOV #X52,R0
92 004286 004342      .CCOUT
93 004288 004767 000676 SG01A5: JSR PC,TIMOUT
94 004290 000207      RTS PC      ;RETURN FROM SG01A
95
96
97 004354
98
99
100 004356 010500      MOV R5,R4      ;GET EC
101 004358 042700 177771 BIC #177771,R0      ;MASK INDEX 0,2,4,6
102 004360 016001 137204 MOV #X5101(0),R1      ;GET MASK ADDR
103 004362 032705 000001 BIT #1,R5      ;EC 000?
104 004364 001401      BEO SG100      JND
105 004366 002301      S-LAB R1      ;YES
106 004368 010503      MOV R5,R3      ;GET EC*2
107 004370 000603      ASL R3
108 004372 004767 000440 JSR PC,ECOUT      ;SEND EVENT CODE
109 004374 032701 000002 BIT #2,R1      IMC?
110 004376 001403      BEO SG101      JND
111 004378 004467 000446 JSR R4,PCOUT
112 004380 000000      .WORD 0      ;PC OF NEXT INSTR.
113 004382 032701 000004 SG101: BIT #4,R1      ;SP?
114 004384 001402      BEO SG102      JND

```

```

115 004430 034767 000456
116 004434 032701 000010
117 004440 001402
118 000442 034767 000472
119 004440 032701 000020
120 004452 001406
121 004454 017302 137712
122 004460 034467 000544
123 004464 040 200
124 004466 030300
125 004470 032701 000040
126 004474 001412
127 004476 016302 137712
128 004502 005722
129 004504 011202
130 004506 004467 000516
131 004512 047 200
132 004514 027000
133 004516 004767 000526
134 004522 002207
135
136
137 004524
138
139
140 004524 016701 001540
141 004530 010146
142 004532 032701 000001
143 004536 001114
144 004540 012601
145 004542 004767 000300
146 004546 006001
147 004550 032701 000001
148 004554 001403
149 004556 004467 000304
150 004562 000300
151 004564 006001
152 004566 032701 000001
153 004572 001402
154 004574 004767 000312
155 004600 006001
156 004602 032701 000001
157 004606 001402
158 004610 004767 000324
159 004614 012703 000660
160 004620 016302
161 004622 006305
162 004624 016305 137712
163 004630 006201
164 004632 032701 000001
165 004636 001412
166 004640 117367 000015
167 004644 022703 000063
168 004650 001416
169 004652 011502
170 004654 004467 000350
171 004660 050 000

```

```

JSP PC,PSOUT
SG102: BIT #10,R1          JSP?
      BEQ SG103          JNO
      JSR PC,SPOUT
SG103: BIT #20,R1          J(PEF)?
      BEQ SG104          JNO
      MOV AXSTAT(3),R2    JGET (REG)
      JSR R0,REGOUT
      BYTE 46,200         IDENT: '4'
      WORD 0
SG104: BIT #40,R1          J(QUE)?
      BEQ SG105          JNO
      MOV XDSTAT(3),R2    JGET BUF, ADDR.
      TST (2)+
      MOV (2),R2
      JSR R0,REGOUT
      BYTE 47,200         IDENT: '5'
      WORD 0
SG105: JSR PC,TIMEOUT      RETURN FROM SG10
      RTS PC
/
/*****
SG111: JEV301 RK DISK1 WITH SUB SELECTOR
/*****
ISEL: PC, PS, SP, REG0...REG5,
/
      MOV MSK110,R1        JGET MASK
      MOV R1,-(SP)
      BIT #1,R1            JSUB SELECTION?
      MNE SG117            JYES-GOTO SUB SELECTOR
SG110: MOV (SP)+,R1
      JSR PC,ECOUT         JSEND EVENT CODE
      ASR R1
      BIT #1,R1            JPC?
      BEQ SG111           JNO
      JSR R0,PCOUT
      WORD 0               JPC OF NEXT INSTRUCTION
SG111: ASR R1
      BIT #1,R1            JPS?
      BEQ SG112           JNO
      JSR PC,PSOUT
SG112: ASR R1
      BIT #1,R1            JSP?
      BEQ SG113           JNO
      JSR PC,SPOUT
SG113: MOV #0,R5
      MOV R5,-(SP)
      ASL R5
      MOV XDSTAT(5),R5    JADDR. OF FIRST REG.
SG114: ASR R1
      BIT #1,R1
      BEQ SG116
      MOV R5,SG115+1
      CMP #63,R5
      BEQ SG114A
      MOV (5),R2
SG114C: JSR R0,REGOUT
SG115: BYTE 00,0

```

EVENT HANDLERS RI-11 MACRO VM02-09

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```

172 004662 040 200
173 004664 005725
174 004666 005233
175 004670 022774 000066
176 004674 001355
177 004676 004767 020546
178 004702 012675
179 004764 000227
180
181 004706
182
183
184 004706 010146
185 004710 010500
186 004712 022020
187 004714 011670
188 004716 010071
189 004720 072127 177773
190 004724 042701 177400
191 004730 005371
192 004732 006371
193 004734 032170 000020
194 004740 001471
195 004742 005201
196 004744 070127 000014
197 004750 042102 177760
198 004754 000001
199 004756 062701 000030
200 004762 010172
201 004764 012601
202 004766 000732
203
204
205
206 004770
207
208 004770 010500
209 004772 006500
210 004774 010000 157712
211 004776 012701 137200
212 004778 012702 000000
213 004780 012103
214 004782 006503
215 004784 000003
216 004786 011304
217 004788 042120
218 004790 021304
219 004792 006120
220 004794 017213
221 004796 012601
222 004798 000227
223
224
225 004834
226
227 004834 004767 000006
228 004836 004767 000006
229 004838 000207
230

```

```

SG116: BYTE 00,200
TST (5)+
INC R3
CMP #64,R3
BNE SG114
JSR PC,TIMEOUT
MOV (SP)+,R5
RTS PC
;RETURN FROM SG11

SG114: ;*** COMPUTE BLOCK NO. (LAST BLK+1 OF ACCESSFD FILE)***
; ((CYL-1)*2+SECT)*14+SECT+39
;
MOV R1,-(SP)
MOV R5,R0
;ADDR. OF REG.3
CMP (0)+,(0)+
;ADDR. OF REG.5 IN R0
MOV (0),R0
;GET DISK ADDR.
MOV R0,R1
ASH #5,R1
;SHIFT CYL ADDR. RIGHT BOUNDED
BIC #177000,R1
DEC R1
;CYL-1
ASL R1
;2
BIT #20,R0
;SURFACE?
REQ SG114B
;0
INC R1
;ADD 1 TRACK
MUL #14,R1
;14
RIC #177760,R0
;MASK SECTOR
ADD R0,R1
;ADD SECTOR
ADD #5,R1
;150 FOR DIRECTORY OF 2 TRACKS
MOV R1,R2
;READY TO SEND STUFF OUT
MOV (SP)+,R1
BR SG114C
;GO BACK FOR OUTPUT
;END OF BLOCK NUMBER CALCULATION
;
;
SG117: ;***SUB SELECTION FOR EV30***
;
MOV R5,R0
ASL R0
MOV X0STAT(0),R0
;ADDR. OF FIRST REG.
MOV #X0STAT,R1
;ADDR. OF FIRST SEL. WORD
MOV #0,R2
;0
SG118: MOV (1)+,R3
;LEFT REG NO.
ASL R3
;REG. ADDR.
ADD R0,R3
;REG. COUNT.
MOV (3),R4
;END ALL "000".
BIC (1)+,R4
CMP (1),R4
BNE SG11D
;GOTO REPORT PROCESSING
SCB #2,SG11B
;END IF 4 LINES
MOV (SP)+,R1
;END EV REPORT PRODUCED-CORR. STACK
RTS PC
;RETURN FROM SG11

;*****
SG12: JEV31-EV33: PIRQ,PP,MMU
;*****
JSR PC,ECCOUT
JSR PC,TIMEOUT
RTS PC

```

```

1
2
3
4
5 005046
6
7 005046 012700 136440
8 005052
9 005050 010500
10 005064
11 005064 000227
12
13
14 005066
15
16 005066 012700 137032
17 005072
18 005076 017700 002000
19 005102 102400
20 005104
21 005110 000224
22
23
24 005112
25
26 005112 012700 137030
27 005116
28 005122 016700 001700
29 005126 005720
30 005134 011000
31 005132
32 005136 000207
33
34
35 005140
36
37 005140 012700 137030
38 005144
39 005150 016700 001732
40 005154
41 005160 000207
42
43
44 005162
45
46 005162 012700 137040
47 005166
48 005172 017700 001710
49 005176 102400
50 005200 032700 000001
51 005204 001400
52 005206 012700 137020
53 005210
54 005216 000003
55 005220 011000
56 005222
57 005226 000204

```

```

      .TITLE BUFFER ROUTINES
      *****
      I
      I
      I*****
      ECOUT:  IOUTPUT OF EVENT CODE; LINKAGE: PC
      I*****
      MOV #XCRLF,R0
      .CHOUT
      MOV R0,R0
      .OCOUT
      RTS PC
      I
      I*****
      PCOUT:  IOUTPUT OF PC-K; K=1ST ARG.; LINKAGE: R4
      I*****
      MOV #XMP,R0
      .CHOUT
      MOV #OLQSTK,R0
      SUB (4),R0
      .OCOUT
      RTS R4
      I
      I*****
      PSOUT:  IOUTPUT OF EVENT PS; LINKAGE: PC
      I*****
      MOV #XMP,R0
      .CHOUT
      MOV #OLQSTK,R0
      TST (0)
      MOV (0),R0
      .OCOUT
      RTS PC
      I
      I*****
      SPOUT:  IOUTPUT OF EVENT SP; LINKAGE: PC
      I*****
      MOV #XMP,R0
      .CHOUT
      MOV #OLQSTK,R0
      .OCOUT
      RTS PC
      I
      I*****
      CPCOUT: IOUTPUT OF EVENT (PC-K); LINKAGE: R4
      I*****
      MOV #XMP,R0
      .CHOUT
      MOV #OLQSTK,R0
      SUB (4),R0
      BIT #1,R0
      BEQ CPCOUT
      MOV #XMP,R0
      .CHOUT
      MOV #XMP,R0
      .OCOUT
      RTS R4
      I
      I*****
      CPCOUT: IOUTPUT OF EVENT (PC-K); LINKAGE: R4
      I*****
      MOV #XMP,R0
      .CHOUT
      MOV #OLQSTK,R0
      SUB (4),R0
      BIT #1,R0
      BEQ CPCOUT
      MOV #XMP,R0
      .CHOUT
      MOV #XMP,R0
      .OCOUT
      RTS R4

```

BUFFER ROUTINES RT-11 MACRO VM02-09

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58
59
60 005230
61
62
63 005230 010400
64 005232
65 005230 022424
66 005240 010200
67 005242
68 005240 000200
69
70
71 005250
72
73 005250 005767 001070
74 005254 001406
75 005250 013702 172544
76 005262 004467 177742
77 005260 052 200 200
78 005271 200
79 005272 000207
80

```

*****
REGOUT: REGISTER OUTPUT; LINKAGE: RA
***** INWORD IN R2; IDENTIFIER IN (R4)
      MOV R4,R8          TADDR OF IDENTIFIER
      LCHOUT
      CMP (R1), (R3)+     SKIP ARGUMENTS
      MOV R2,R0          GET VALUE OF WORD
      LCHOUT
      RTS R4              RETURN FROM REGOUT
*****
TIMOUT: SEND TIME FROM KW11-P
*****
      TST TIMLG          TIME LOG?
      BEQ TIMOUT1        NO
      MOV #172544,R2     GET CNT
      JSR RA,REGOUT      SEND II
      BYTE 52,200,200,200  IDENTIF. I
*****
TIMOUT1: RTS PC          RETURN FROM TIMOUT

```

```

1
2
3
4
5 005274 005002
6
7 005276 012700 137170
8 005302 000200
9 005304 011000
10 005306 005001
11 005310 032700 000001
12 005314 001402
13 005316 004707 000032
14 005322 022701 000017
15 005326 001403
16 005330 005201
17 005332 006200
18 005334 000705
19 005336 022702 000004
20 005342 001403
21 005344 002702 000004
22 005350 000752
23 005352 000207
24
25
26 005354
27
28
29
30 005354 010203
31 005356 006303
32 005360 006303
33 005362 000103
34 005364 006303
35 005366 010304
36 005370 002704 137612
37 005374 006303
38 005376 010305
39 005400 002705 137412
40 005402 011403
41 005406 005737 140030
42 005412 001015
43 005414 012325
44 005416 011315
45 005420 005743
46 005422 012704 137170
47 005426 000200
48 005430 005724
49 005432 011423
50 005434 012704 000340
51 005440 000104
52 005442 010413
53 005444 000207
54 005446 012523
55 005450 011513
56 005452 000207
57

.TITLE REPLACEMENT PROCEDURE
*****
*****
*****
EXCH: CLR R2          IR2=0-J (GROUP INDEX 0-M)
*****
*****
REPL1: MOV #XMSK0,R0    IGET DEVICE MASK J
      ADD R2,R0
      MOV (0),R0
      CLR R1
      ICOUNT 0-15
      BEQ REPL3          INEXT DEVICE SET?
      JSR PC,REPLAC      END
      CNP #15,R1         IYES DO EXCHANGE
      BEQ REPL4          IEND OF MASK?
      INC R1             IYES
      ASR R0             IEND
      OR REPL2           ISHIFT MASK
      CMP #MAX0M,R2      IALL GROUPS DONE?
      BEQ REPL5          IYES - DONE
      ADD #4,R2          IINCR MASK OFFSET.
      BR REPL1           INEXT GROUP.
      RTS PC            IRETURN FROM EXCH
*****
*****
REPLAC: JSUBROUTINE TO REPLACE VECTOR I (=R1)
*****
      IIN GROUP J (=R2/4). THE DIRECTION OF THE
      REPLACEMENT IS SPECIFIED IN DIRECT
*****
      MOV R2,RT          IALC. OFFSET IN VECTOR
      ASL R3              I(J+16+I)*2+VECTAB
      ASL R3
      ADD R1,R3
      ASL R3
      MOV R3,R4          I=2 FOR WORD
      OR #VECT0,R4
      ASL R3
      MOV R3,R5          IUFFS. IN VECTSAV=(J+4+1)*4
      ADD #VECSV,R5
      MOV (4),R3
      TST #XQIN          IADDR. OF VECTOR
      BNE SAVBAK         IWHICH WAY?
      MOV (3)+(5)+
      MOV (3),(5)
      TST -(3)
      MOV #XMSK0,R4
      OR R2,R4           IGET PC(J)
      TST (4)+           IADDR. OF REPLACEMENT PC
      MOV (4),(3)+
      MOV #H7,R4
      ADD R1,R4
      MOV R4,(3)
      RTS PC             IREPL.PS(J)
      SAVBAK: MOV (5)+(3)+
      MOV (5),(3)
      RTS PC             IREST. SAVED VECT. TO ORIG.
      IRETURN FROM REPLAC

```

```

58 005454
59
60 005454 010346
61 005456 012703 000010
62 005462 012190
63 005464 004767 000030
64 005470 077304
65 005472 012700 136440
66 005476 004767 000112
67 005522 004767 000162
68 005526 077215
69 005530 012603
70 005532 000207
71 005534 000000
72 005536 000 200
73
74 005520 010046
75 005522 010146
76 005524 010206
77 005526 012702 000006
78 005532 010401
79 005534 002701 177770
80 005540 002701 000060
81 005544 110146
82 005546 000000
83 005550 000200
84 005552 000200
85 005554 077212
86 005556 012702 000006
87 005562 112677 000454
88 005566 005267 000450
89 005572 077205
90 005574 012700 136442
91 005620
92 005604 012602
93 005606 012601
94 005610 012600
95 005612 000207
96
97
98 005614 121027 000200
99 005620 001422
100 005622 121027 000000
101 005626 001405
102 005630 112077 000406
103 005634 005267 000402
104 005640 000765
105 005642 112777 000015 000372
106 005650 005267 000366
107 005654 112777 000012 000360
108 005662 005267 000354
109 005666 000207
110
111
112
113
114

```

```

DUMP: JUMPS N LINES (N=R2) OF 8 WORDS
FROM STARTING ADDR IN R1
MOV R3, -(SP)
DUM1: MOV #10, R3 76 WORDS
DUM2: MOV (R1)+, R0
JSR PC, OCPUT
SOB R3, DUM2
MOV #XCHLF, R0
JSR PC, CHPUT 1PRINT CRLF
JSR PC, SFN00F
SOB R2, DUM1
MOV (SP)+, R3
RTS PC
CRLF: .WORD 0
BLANK: .BYTE 40, 200
OCPUT: MOV R0, -(SP)
MOV R1, -(SP)
MOV R2, -(SP)
MOV #6, R2 76 DIGITS
MOV R0, R1
OCPUT1: RLC #177770, R1 1MASK RIGHTMOST DIGIT
ADD #60, R1 1CONVERT TO ASCII
MOV R1, -(SP)
ROR R0
ASR R0
ASR R0
SOB R2, OCPUT1
MOV #6, R2
OCPUT2: MOV R0, (SP)+, #TXTPTR 1PUT NUMBER IN BUFFER
INC TXTPTR
SOB R2, OCPUT2
MOV #XBLANK, R0 1PUT A BLANK
CHOUT
MOV (SP)+, R2
MOV (SP)+, R1
MOV (SP)+, R0
RTS PC 1RETURN FROM OCPUT
CHPUT: CHPB (R0), #200 1ETX
BEQ CHPUT2 1YES
CHPB (R0), #0 1ETX WITH CRLF?
BEQ CHPUT1
MOV R0, (R0)+, #TXTPTR 1PUT CHAR
INC TXTPTR
BR CHPUT
CHPUT1: MOV R0, #15, #TXTPTR 1PUT CRLF
INC TXTPTR
MOV R0, #12, #TXTPTR
INC TXTPTR
CHPUT2: RTS PC 1RETURN FROM CHOUT

```

```

TITLE TRANSMISSION ROUTINE
*****

```

```

115 005670 010046 SENDBF: MOV R0,-(SP)
116 005672 010146 MOV R1,-(SP)
117 005674 012737 052525 177570 MOV #52525,#177570
118 005702 042737 000100 164000 BIC #100,#COMCHI
119 005710 012701 137046 MOV #TXT0F,R1
120 005714 020167 000322 CMP R1,TXTPTR
121 005720 0010P2 BNE SEND00
122 005722 000167 000124 JMP SEND03
123 005726 032767 000001 000306 SEND00: BIT #1,TXTPTR
124 005734 0010P5 BEQ SEND01
125 005736 112777 000001 000276 MOVB #1,TXTPTR
126 005744 005267 000272 INC TXTPTR
127 005750 012100 SEND01: MOV (R1)+,R0
128 005752 ,SHOWRD
129 005764 020167 000252 CMP R1,TXTPTR
130 005770 001367 BNE SEND01
131 005772 032737 000001 177570 BIT #1,#177570
132 006000 0014P2 BEQ SEND02
133 006002 052700 000002 BIC #2,R0
134 006004 012700 000000 SEND02: MOV #EOT,R0
135 006012 ,SHOWRD
136 006024 ,RCVWRD
137 006030 012737 177777 177570 MOV #177777,#177570
138 006044 012767 137046 000170 MOV #TXT0F,TXTPTR
139 006052 012601 SEND03: MOV (SP)+,R1
140 006054 012600 MOV (SP)+,R0
141 006056 052737 000100 164000 BIC #100,#COMCHI
142 006060 000240 NOP
143 006066 000240 NOP
144 006070 000240 NOP
145 006072 000207 RTS PC
146

```

```

TEST
GET START ADDR. OF BUFFER
ANY TEXT TO SEND?
YES
NO-EXIT
NEXT FREE=000?
NO
YES, PUT FILLER IN HIGH BYTE
AND SET EVEN BOUNDARY
SEND A WORD
END OF LINE?
NO
CHECK SWH
NOT SET
MARK END
YES
WAIT FOR ACK
TEST
RESET TEXT POINTER
ENABLE SUPER INTERRUPT.
RETURN FROM SENDBF

```


RESIDENTY DATA SECTION RT-11 MACRO VM02-09 00111126 PAGE 11+

56 006336	000	000	SLI103: .BYTE 0,0 ;17
57 006340	000	000	.BYTE 0,0 ;20
58 006342	000000		TIME01: .WORD 0
59 006344	000000		TIME1: .WORD 0
60 006346	000000		TREG1: .WORD 0
61 006350	000000		TREG2: .WORD 0
62 006352	10036		TINT1: .XRTI
63 006354	000340		TINT2: 0K7
64 006356	000003		ENDCOM: .WORD 3
65			

TIME INTERR. TO DUMMY MTI
TIME INTERR. AT PRI0. 7
END OF COMM. AREA

[illegible]

```

?
DISPT0: XOE
XSG00
XSG01
XSG0P
XSC00
XSG00
XSG01
XSG01
XOE
XOE
XOE
XOE
XOE
XOE
XOE
XOE
DISPT1: XSG1A
XSG1A
XSG1A
XSG1A
XSG1A
XSG1B
XSG1P
XSG11
XSG12
XOE
XOE
XOE
XOE
```

```

JEV 3, DISPATCH ERROR
JEV 1, TRAP 4
JEV 2, TRAP 12
JEV 5, BPT
JEV 4, IOT
JEV 5, PWFL
JEV 6, EMT
JEV 7, TRAP
JEV 10,
;
;
;
;
;
;
JEV 20, TX
JEV 21, TP
JEV 22, PR
JEV 23, PP
JEV 24,
JEV 25,
JEV 26,
JEV 27,
JEV 30, RK
JEV 31, PIR2
;
;
;

```

1
2 006454 134742
3 006456 134430
4 006460 134410
5 006462 134700
6 006464 134426

7
8 006466
9
10
11
12
13
14 006666 000000
15 006670 000004
16 006672 000010
17 006674 000014
18 006676 000020
19 006700 000024
20 006702 000030
21 006704 000034
22 006726

23
24
25
26 006726 000300
27 006730 000304
28 006732 000070
29 006734 000074
30 006736 000104
31 006740 000000
32 006742 000004
33 006744 000008
34 006746 000220
35 006750 000200
36 006752 000244
37 006754 000250
38 006756 000000
39 006760 000200
40 006762 000000
41 006764 000000
42
43
44

45 006766 007026
46 007028 177560
47 007030 177564
48 007032 177550
49 007034 177554
50 007036 172500
51 007040 000000
52 007042 000000
53 007044 000000
54 007046 177400
55 007050 177772
56 007052 000000
57 007054 177572

DISPT2: XOE
XTABUP
XNANG
XENDRU
XCONT

100 DISP. ERROR
102 TABLE UPDATE
104 SUSPEND HOST
106 TERMINATE SESSION
110 CONTINUE AFTER HAND

VECSAVE: BLANK N=2
VECT. ADDR.
-J- -I- -VECT-

*** GROUP 0 ***

VECTAB: 0 10 2 NOT IMPLEMENTED
4 10 1 TRAP 4
10 10 2 TRAP 10
14 10 3 GPT
20 10 4 IOT
24 10 5 POF
30 10 6 EMT
34 10 7 TRAP
-+16. PRESEVED FOR 8 VECT. IN GRP 0

*** GROUP 1 ***

300 11 0 TR
304 11 1 IP
70 11 2 PR
74 11 3 PP
100 11 4 KWI-P
0 11 5
0 11 6
0 11 7
220 11 10 KWI
240 11 11 PING
244 11 12 FPP
250 11 13 HHU
0 11 14
0 11 15
0 11 16
0 11 17

DEVICE STATUS REGISTER TABLE

OVSTAT: -+32. 170 STATUS FOR TRAPS IN GRP 0
177560 11 0 TKS
177564 11 1 TPO
177550 11 2 PHS
177554 11 3 PPS
172500 11 4 KWI-P
0 11 5
0 11 6
0 11 7
177000 11 10 RNCS
177772 11 11 PING
0 11 12 FLT ERR
177572 11 13 HHU 594

LOG:: INTERACTIVE SYSTEM MONITOR RT-11 MACRO VM02-09 09:15:55. PAGE 1

```

1      .TITLE LOG:: INTERACTIVE SYSTEM MONITOR
2      *****
3      *
4      * CONTROL AND LOGGING PROGRAM FOR MONITOR SYSTEM "RT-11"
5      * LOG IS THE MAIN PROGRAM AND HAS TO BE LINKED TO
6      * THE OBJECT MODULES LOGIO AND LOGDEF.
7      * ALL ENTRIES AND ROUTINES FOR THE HARDWARE MONITOR ARE
8      * FORESEEN, THE CORRESPONDING CODE IS REPLACED BY WORDS.
9      *
10     *****
11     *
12     .GLOBAL TTYOUT,OPRINT,OPEN,PRTOUT,OPENP
13     .GLOBAL BITPRT,BITPOS,BITS,DECODE,IOERR,IOERS
14     .GLOBAL IOERR1,PUTRNO,CHDIR,CHDIR1,IOERR2
15     .GLOBAL CMTOUT,PRTT,PRTNR,TTYNO,CCLK
16     .GLOBAL LDI,CHI,CHIR,CTI,CTI1,CTI2,CTI3
17     .GLOBAL HANG,TRICL,SRP,SPF,DISP,DISP1
18     .GLOBAL ESEMP,SEMP,SEMP1,SEMP2,SEMP3,SEMP4
19     .MCALL .PRINT,.REDEF,.EXIT,.OUTREF,.OFF,.SEMP,.RECV,.DEF
20     .MCALL .TTYIN,.TTYOUT,.WRITE,.V2,.TTYIN,.WRITE,.READY,.READY
21     .MCALL .CLOSE,.CLOSE,.CLOSE,.WAIT,.FETCH,.ENTER,.RELEASE
22     .V2..
23     .HFDDEF
24     .ASECT
25     .=1640
26     IPS=117564
27     FUF=04
28     ACK=06
29     RW=352
30     BP=340
31     PSW=17774
32     HANG=4
33     TBUF=172542
34     TSTAT=172540
35     PLIV=320
36     PLB=16400
37     PLB=16402
38     PLB=16404
39     IIS=PLB
40     IOS=PLB
41     IIR=PLB
42     IGR=PLB
43     IIV=PLB
44     IOV=PLB
45     COMCH=IIS
46     COMCH=IOS
47     TTYOUT=117564
48     MAXCH=22
49     MAXEV=10
50     LSPFIL=MAXEV
51     HSPFIL=ENDCOM-MASK+2
52

```

```

***** I.P. STATUS
***** CHT BUFFER
***** STATUS REGISTER
***** FOR PARALLEL LINK
***** REG. -
***** HUF. -
***** HUF. -
***** FOR PAR. LINK

```

```

***** STATUS REG. IF CMT. CHANNEL IN
***** TTY OUTPUT STATUS

```

```

***** LENGTH OF LOG SPACE

```

Light copy

```

1      .TITLE INIT AND WAITLP
2      .*****
3
4      .LIST MP
5
6      .MACRO ,SNDWRD
7      TST *+PLSR
8      HLT ,=4
9      MOV R2,*+PLSR
10     .ENOM
11
12     .MACRO ,RCVWRD
13     TST *+PLSR
14     HLT ,=4
15     MOV *+PLSR,R2
16     .ENOM
17
18     START: MOV *,SP
19            CLR R5
20            MOV *+RUF1,SBPTH
21            MOV *+STAR,R4
22            JSR PC,PRINT
23            R1C *+PC,*+105
24            MOV *+RUFEC,*+11V
25            MOV *+R6,*+11V+2
26            JSR PC,INIT
27
28     .*****
29
30     WAITLP: ;SYSTEM WAIT LOOP AND TASK DISPATCHER
31
32     .*****
33
34     001050 012706 001002
35     001051 005005
36     001052 012707 011130 010112
37     001053 012708 007240
38     001054 000767 000000
39     001055 002737 007100 164000
40     001056 012737 005600 000520
41     001057 012737 000520 000522
42     001058 000767 000000
43
44     001052
45
46     001052 005037 177776
47     001053
48     001054 103402
49     001055 000767 001062
50
51     001066 005767 012170
52     001072 001767
53     001074 005000
54     001076 012705 000401
55     001102 130567 012160
56     001106 001406
57     001110 010000
58     001112 010506
59     001114 000774 012720
60     001120 012605
61     001122 012604
62     001124 000201
63     001126 005505
64     001130 105150
65     001132 000729
66     001134 000762
67
68     .CLP *+PSH
69     .TTRM
70     RCS WAIT1
71     JSR PC,SYSCMD
72
73     .*** SCAN DISPATCH WORD AND DISK SCHEDULED TASK ***
74
75     WAIT1: TST *+DISPD
76            BEQ WAITLP
77            CLR R4
78            MOV *+R1,R5
79            R1C *+R5,*+DISPD
80            BEQ WAIT1
81            MOV R4,*+R5
82            MOV R5,*+R5
83            JSR PC,*+DISPD(4)
84            MOV *+R5,*+R5
85            MOV *+R5,*+R5
86
87     WAIT2: CLC
88            ASL R5
89            RCS WAITLP
90            TST *+R5
91            BEQ WAIT1
92
93     .PRIG/AC-BITS CLEAN
94     .IF THERE SOME ITV INPUT?
95     .NO
96     .YES-ANALYZE AND DISPATCH END
97
98     .IF THERE ANYTHING THERE?
99     .IF THERE WAITING
100    .YES
101    .IF THERE LOWER BYTE OF DISPATCH
102    .IF THERE
103    .NO
104    .IF THERE REGS
105    .IF THERE DISPATCH TASK
106
107    .IF THERE PLEASE
108    .END OF DISPATCH
109    .IF THERE TO DISPATCH
110    .IF THERE NEXT HIT

```

UTILITY TASKS

```

1  TITLE UTILITY TASKS
2  *****
3
4
5
6  0001130      005037      177546
7  0011130      013767      000180      012200
8  0021132      013767      000180      000180
9  0031150      012737      000000G      000180
10 0041150      012700      010600
11 0051102      005720
12 0061102      012720      030060
13 0071100      012720
14 0081170      012720      030060
15 0091174      112720      000060
16 0101200      115720
17 0111202      012710      030060
18 0121205      012703      013162
19 0131212      012701      000041
20 0141210      005220
21 0151200      005301
22 0161222      001375
23 0171220      005737      164000
24 0181230      005737      000180      177546
25 0191250      002737      000180      100000
26 0201200      002737      000180
27 021200      002007
28
29
30 0001200
31
32 001200      010046
33 0021200      010146
34 0031202      010346
35 0041200      010446
36 0051250      012501
37 0061260      022703      000120
38 0071260      125045
39 0081200      012700      010640
40 0091272      002767      000000G
41 0101270      000032
42 0011300
43 0021330
44 0031310      042754
45 0041310      162700
46 0051320      130019
47 0061322      001413
48 0071324      022027      000014
49 0081330      100010
50 0091332      020574
51 0101330      010400
52 0011300      000060G
53 0021300      000060
54 0031304      105721
55 0041300      000060
56 0051350      001405
57 0061352      000060
58 0071350      001551

```

UTILITY TASKS RT-11 MACRO MH02-09 00:15:55 PAGE 3+

58 001356	105737	177564	
59 001362	100375		
60 001364	012600		
61 001366	012603		
62 001370	012601		
63 001372	012600		
64 001374	000205		
65			

TSTR 00TPS
 BPL -d
 PR183: MOV (SP)+,R4
 MOV (SP)+,R3
 MOV (SP)+,R1
 MOV (SP)+,R0
 R15 R5

RETURN FROM PRINTB

CMDI FOR HOST MODE SPECS RT-11 MACRO VM02-09 00:15:55 PAGE 4

FILE CMDI FOR HOST MORE SPES

```

00000000
CMDDIPT: JCOMMAND INTERPRETER FOR HONorable UPDATE
00000000

```

```
MOV R2,=(SP)
MOV R1,=(SP)
MOV R2,=(SP)
MOV R3,=(SP)
MOV R5,=(SP)
CLR R5
ISET LOOP $JITPC
```

CMD10: ***** UPDATE ***

ADVISORY #MSXIA, RR
JSR PC, 44141

 COMMAND: P/N INJENKRETEH *****

AREAD Y DW NO WITH LEADING 64MM462

SECRET

SECRET

1991

WILLY ECKHARD
FUTSEYCH

Y JNUH 15711 JMI

INDEX

IVES
IVES-DISPATCH FOR VACHE BUTCH
IVES-DISPATCH FOR LEAT WORTH

INDEX 4

ANNU. OF MASS.

Wm. H. H. H. H.

1987, 1994, 1998

```

58
59 001560
60
61 001560 000301
62 001570 000301 011016
63 001574 000301 00000000
64 001570 000301 00000000
65 001574 000301
66
67 001576
68
69 001580 000301
70 001580 000301 00000000
71 001584 000301 00000000
72 001580 000301 177576
73
74 001584
75
76 001584 000301
77 001580 000301 00000000
78 001584 000301 00000000
79 001580 000301 177576
80
81 001584
82
83 001584 000301
84 001584 000301
85 001584 000301 00000000
86 001584 000301
87 001584 000301 00000000
88
89 001584
90
91 001584 000301
92 001584 000301 011042
93 001584 000301 00000000
94 001584 000301 00000000
95 001584 000301 00000000
96 001584 000301 00000000
97 001584 000301
98 001584 000301
99 001584 000301 011042
100 001584 000301
101 001584 000301
102 001584 000301
103 001584 000301
104
105 001584
106
107 001584 000301 011042
108 001584 000301 00000000
109 001584 000301 00000000
110 001584 000301 00000000
111 001584 000301 00000000
112 001584 000301
113 001584 000301
114 001584 000301

```

```

1
CMOI19: **** SUBMASK I/O ****
1
    ASL R1                INDEX+2
    ADD MASK025,R1        INCR. OF SUBMASK
    JSR PC,BINPR1         PRINT MASK
    JSR PC,BINPR1         INCR. NEW MASK
    BR CMOI6              INCR. TEXT INDEX

CMOI10: **** SUBMASK UPDATE ****
1
    TST (5)+              INCR. LOOP SWITCH
    MOV MASKIN,R0
    JSR PC,PRINT
    JMP CMOI11            INCR. Y/N

CMOI12: **** SELECT WORD UPDATE ****
1
    TST (5)+              INCR. LOOP SWITCH
    MOV MASKIN,R0
    JSR PC,PRINT
    JMP CMOI11            INCR. Y/N

CMOI14: **** SELECT WORD I/O ****
1
    ASL R1-               INDEX+2
    MOV R1,R2             SAVE INDEX
    CMP R2,R1             INDEX+1
    BLT CMOI10            INDEX+1
    MOV R2,R3             INCR. 2 BYTES

CMOI15: **** BYTE I/O ****
1
    CLR R2
    MOV SL0000(2),R0      GET OLD BYTE VALUE
    JSR PC,BINPR1         PRINT IT
    MOV R0,R2
    JSR PC,PRINT          PRINT :
    JSR PC,BINPR1         GET NEW VALUE
    TST R2
    BHI CMOI15A           IF NEW VALUE
    MOV R1,SL0000(2)      STORE IT

CMOI15A: INC R2           SECOND BYTE
    DEC R3
    BFE CMOI15
    BR CMOI6

CMOI16: **** WORD I/O ****
1
    MOV SL0000(2),R0      GET OLD WORD VALUE
    JSR PC,PRINTAND
    MOV R0,R2
    JSR PC,PRINT          PRINT IT
    JSR PC,BINPR1         GET NEW VALUE
    TST R2
    BHI CMOI16           IF NEW VALUE
    MOV R1,SL0000(2)      STORE IT

```

170

Flight copy

```

1
2
3
4
5
6 006214
7
8
9 006214 135267 005044
10 006220 042767 000002 005040
11 006220 145767 005062
12 006232 041443
13 006234 145767 002665
14 006240 001405
15 006242 012703 012134
16 006246 145467 002053
17 006252 000002
18 006254 012703 011134
19 006260
20 006324 103004
21 006326 004567 000000
22 006332 000
23 006334 000472
24 006336 005267 004726
25 006342 145067 004716
26 006346 145767 005001
27 006352 001405
28 006354 012746 000300
29 006360 012746 004370
30 006364 020167 177240
31 006370 000207
32
33
34 006372
35
36
37 006372 145267 004666
38 006376 042767 000074 004662
39 006404 016701 022522
40 006410 022167 022514
41 006414 140412
42 006416 166701 002596
43 006422 010103
44 006424 016767 002500 000004
45 006432 024567 172610
46 006436 000000
47 006440 012424
48 006442 012702 013134
49 006446 166702 002456
50 006452 010203
51 006454 016767 022450 000004
52 006462 024567 172560
53 006466 000002
54 006470 162701 011134
55 006474 014103
56 006476 012767 011134 000004
57 006504 044567 172536

```

```

*TITLE LOG OUTPUT ROUTINES
*****

```

```

*****
RECORD: OUTPUT OF A BUFFER ONTO MASS STORAGE
*****

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```

LOG OUTPUT ROUTINES

RT-11 MACRO VM02-09

08:15:55 PAGE 104

```

58 006510 000000
59 006512 105267 004546
60 006514 105767 004631
61 006522 001406
62 006524 012746 000300
63 006526 012746 006540
64 006534 000167 177072
65 006540 000207
66
67
68
69
70 006542
71
72 006542 105267 004516
73 006546 002767 000010 004512
74 006554 016700 002350
75 006560 005120
76 006562 004767 000000
77 006566 002701 177400
78 006572 010100
79 006574 004767 000000
80 006580 105067 004460
81 006584 105767 004543
82 006586 001406
83 006592 012746 000300
84 006596 012746 006526
85 006602 000167 177002
86 006620 000207
87
88

```

```

DUMP1A: .WORD 0
DUMP3: CLR0 CH0SY
      TST0 OVH0N
      REU DUMP4
      MOV #006,-(SP)
      MOV #DUMP4,-(SP)
      JMP PLVHOR
DUMP4: RTS.PC

```

```

      WAS THERE DATA OVERRUN
      AND
      YES-SET UP NEW STACK
      FROM RTI IN PLVREC
      AND JUMP THERE
      RETURN FROM DUMP

```

```

DUMP4: RTS.PC
      .
      .
      .
      *****
GRAPH1: DISPLAY DATA GRAPHICALLY
      *****

```

```

      INC0 CH0SY
      RIC #14,019000
      MOV SREC1,R0
      TST R0
      JSR PC,DECODE
      RIC #17,010,R1
      MOV R1,R0
      JSR PL,CPA001
      CLR0 CH0SY
      TST0 OVH0N
      REU GRAPH10
      MOV #006,-(SP)
      MOV #GRAPH12,-(SP)
      JMP PLVHOR

```

```

      INDICATE T/O BUSY
      ALLK DISPATCH HIT
      GET-ET

```

```

      INTO R1

```

```

      BRAM IN MODE 1

```

```

      WAS THERE DATA OVERRUN
      AND
      YES-SET STACK FOR RTI

```

```

GRAPH10: RTS.PC
      .
      .

```

```

      RETURN FROM GRAPH

```



```

115
116
117
118 010730 002326
119 010732 002334
120 010734 002506
121 010736 003214
122 010740 003216
123 010742 005114
124 010744 005166
125 010746 005122
126 010750 005374
127 010752 005476
128 010754 000220
129
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134 010756 010516
135 010760 010520
136 010762 010523
137 010764 010526
138 010766 010531
139 010770 010534
140 010772 010537
141 010774 010542
142 010776 010545
143 011000 010516
144 011002 010507
145 011004 010552
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151 011006 000216
152 011010 000200
153 011012 000002
154 011014 000200
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158 011016 077777
159 011020 010662
160 011022 010706
161 011024 010576
162 011026 031016
163 011030 031062
164 011032 001602
165 011034 000000
166 011036 000000
167 011040 000020
168
169
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```

DISPATCH TABLE FOR SYSCMD INDX

DISPT	WAITUP	INDX
SETPLG	100	
CSIS	102	
CSIM	104	
CSIL	106	
PNOC	108	
TERM	110	
CONT	112	
EWAS	114	
SPEYUP	116	
ENTOUT	118	

IDENTIFIERS FOR NUM DUMP

NAME	INDX
IDADD: TDER	10
TOPC	12 PC
TDPS	14 PS
TOSP	16 SP
TDCP	18 CP
TDEW	20 EW
TDSK	22 SW
TDRF	24 RF
TDR	26 R
TDER	28
TDRF	30 RF
TDRS	32 RS

BEGINNING OF COMMUNICATION AREA

MASK0: WORD 0000000000000000

PC0: WORD 0

MASK1: WORD 0000000000000000

PC1: WORD 0

SUPMASK INDX

MASK	WORD	INDX	MASK	WORD	INDX
MASK005	0011111111111111	10	MASK PC,PS,SP FOR EV 1-5		
MASK067	0000000000000000	11	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK101	0000000000000000	12	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK123	0000000000000000	13	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK145	0000000000000000	14	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK167	0000000000000000	15	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK189	0000000000000000	16	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK211	0000000000000000	17	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK233	0000000000000000	18	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		
MASK255	0000000000000000	19	MASK SEL,PL,PS,SP,(PL-2),(PL-1),EV6,7		

SELECT WORDS FOR CONDITIONAL EVENT REPORT

INDEX

DATA SECTION

RT-11 MACRO VM02-09

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```

229 013345 000 SM0TEV: .BYTE 0 ITRIGGER EV FOR MODE 0.
230 013346 000 SM0DEW: .BYTE 0 IDISPLAY, MODE 0
231 013347 000 SM1EV: .BYTE 1 IEVENT FOR Y IN MODE 1
232 232 .EVEN
233 013350 000200 .WORD 0
234 234
235 235 END OF LOG SPEC FILE
236 236
237 013352 000 SUPDFL: .BYTE 0 I1: FILE UPDATE, FILE STILL OPEN
238 013353 000 OVRUN: .BYTE 0 I1: S-LOGGING DATA OVERRUN
239 239 .EVEN
240 013354 000300 LCKSV1: .WORD 0 I1: SAVE AREA FOR KMI VECT.
241 013355 000 DSKHND: .BLKW 000 I1: DISK HANDLER SPACE
242 016010 074503 TTYHND: .BLKW 000 I1: TTY HANDLER SPACE
243 016010 074503 EXT1: .HAUSD "SPC" IDEF.EXT. FOR LOG SPEC FILES
244 016020 074503 .HAUSD "SPC"
245 245 .BLKW 2
246 016020 046537 EXT2: .HAUSD "LOG" IDEF.EXT. FOR LOG,REC,OUTPUT
247 016030 046537 .HAUSD "LOG"
248 248 .BLKW 2
249 016030 076452 EXT3: .HAUSD "TAM" IDEF.EXT. FOR HOST TABLE
250 016040 076452 .HAUSD "TAM"
251 251 .BLKW 2
252 016040 120200 OUTSPEC: .BLKW 10 I1: LOCK FOR C11SPC
253 016100 120200 T1SPEC: .HAUSD "TT" I1: LOCK FOR TTY DEF.
254 016100 052450 021420 .HAUSD "DUMMY, EXT"
255 255 .EVEN
256 256 .TITLE LOG - HIMS CONTROL PRGRM.
257 001000 .END START

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192

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58 000106 012700 000000G CNTDT3: MOV #CNILX,R0      IPRINT IV
59 000112 004767 000272     JSR PC,PRINT
60 000116 012400     MOV (4)+,R0      IGET COUNT
61 000120 004767 000352     JSR PC,UPRINT
62 000124 005302     DEC R2
63 000126 001354     BNE CNTDT2
64 000130 012700 000300G     MOV #CNILF,R0
65 000134 004767 000250     JSR PC,PRINT
66 000140 000145     BR CNTDT1
67 000142 012700 000000G CNTDT4: MOV #CNILF,R0
68 000146 004767 000236     JSR PC,PRINT
69 000152 052737 000100 000000G R15 #100, #COMCHI
70 000160 012604     MOV (SP)+,R0
71 000162 012603     MOV (SP)+,R3
72 000164 012602     MOV (SP)+,R2
73 000166 012601     MOV (SP)+,R1
74 000170 012600     MOV (SP)+,R0
75 000172 000207     RTS PC      IRETURN FROM CNTDT
76
77
78 000174     *****
79     TIMLOG: IAPPEND SYS-TIME TO A RECORD
80     *****
80 000174 010146     MOV R1,-(SP)
81 000176 010246     MOV R2,-(SP)
82 000220 105767 000000G     TSTB HANG      ILOGGING ACTIVE?
83 000204 001603     BNE TMLUG5     AND
84 000200 012701 000000G     MOV #TIFL0,R1
85 000212 016702 000000G     MOV #PTIR,R2
86 000210 005302     DEC R2      ISTACKUP OVER BOT
87 000220 000227 000000G TMLUG1: CMP R2,#SHUF2     IEND OF BUFFER 1?
88 000224 001003     BNE TMLUG2     AND
89 000226 052767 000002 000000G     HIS #2,DISPWO     IYES-DISPATCH BUFFER RECOMMING
90 000234 000227 000000G TMLUG2: CMP R2,#SHUF     IEND OF BUFFER 2?
91 000240 001415     BNE TMLUG4     IYES-APP ARGUND
92 000242 112122     TMLUG3: MOV3 (1)+,(2)+     ILOGIC IT
93 000244 121127 000204     CMPS (1),#200     IEND OF TFIELD?
94 000250 001363     BNE TMLUG1     AND-EXT
95 000252 052702 000001G     HIT #1,#2     IYES-PC1 DASY?
96 000256 001401     BNE TMLUG3A     AND
97 000260 105322     CLRH (2)+     IYES-ADJUST TO EVEN BOUNDARY
98 000262 010267 000000G TMLUG3A: MOV R2,#BPTH     IADJUST RECORD END POINTERS
99 000266 010267 000000G     MOV R2,#BEC2
100 000272 000243     BR TMLUG5
101 000274 052767 000002 000000G TMLUG4: HIS #2,DISPWO     IYES-DISPATCH
102 000302 012702 000000G     MOV #SHUF1,R2     IYES-DISPATCH BUFFER RECOMMING
103 000310 105267 000000G     INCB SHUF54     IAPP ARGUND
104 000312 000153     BR TMLUG4     IYES-ITC BACK TO BUF 1
105 000314 012602     TMLUG5: MOV (SP)+,R2
106 000316 012601     MOV (SP)+,R1
107 000320 000207     HIS PC      IRETURN FROM TMLUG
108
109
110 000322     *****
111     LCLAT: ILTRF CLOCK INTERRUPT HANDLER
112     *****
112 000322 010046     MOV R0,-(SP)
113 000324 012700 000011G     MOV #TIFL0+11,R0
114 000330 105213     IGET (R0)

```

LOGIO::I/O ROUTINE FOR LOG RT-11 MACRO VM02-09 00:10:51 PAGE 10

ADDRESS	DATA	INSTRUCTIONS	REMARKS
000072	122710	CMPS #72,(0)	POWERFLOW?
	021001	REQ LCLK1	YES
	020021	RM LCLK5	NO-NONE
000060	112710	MOV #40,(0)	
	115200	INCB -(0)	POWERFLOW?
000066	122710	CMPS #80,(0)	YES
	021401	REQ LCLK2	NO-NONE
000060	112710	RM LCLK5	
000060	105700	MOV #01,(0)	IS-IF REWIND
	000402	TSTB -(RM)	
000060	000402	RM LCLK4	
	112710	MOV #00,(0)	
	105200	INCB -(0)	POWERFLOW?
000072	122710	CMPS #72,(0)	YES
	122710	REQ LCLK3	
	021002	MOV (SP)+,R0	RETURN FROM LCL
	000202	RTT	
	000400		
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286	001126	005703	BINP001: JST R3	EMPTY WORD
287	001130	001401	REG BINR04	STORE LEFT WORD
288	001132	010211	MOV R2,(1)	IS BIT HIGH? HUFFEN
289	001134		BINR04: ITYTH	
290	001144	012602	MOV (SP)+,R2	
291	001142	012603	MOV (SP)+,R3	
292	001144	000207	RTS PC	RETURN FROM STAMP
293				
294			*****	
295	001140		BINP001: PRINT WORD BINR04; LINKAGE: PC	
296			*****	
297			*****	
298	001136	010136	MOV R1,-(SP)	
299	001150	012700	MOV 001105,RJ	IDENTIFY BIT IDENTIFIER STAMP
300	001154	004767	JSR PC,P2,R1	
301	001150	011100	MOV (1),R0	STAMP TO BE PRINTED
302	001160	012701	MOV R1,R0	
303	001160	000000	BINP001: BIT 012602,R0	LINKER LEFTMOST BIT
304	001172	001403	REG BINP02	
305	001174	112721	MOV R2,R1,(1)+	
306	001170	000002	RTS PC	
307	001202	112721	BINP001: MOV R2,R1,(1)+	
308	001220	000300	BINP001: JST R0	
309	001210	000307	CHP R1,R2,R3,R4	
310	001214	001364	MOV R1,R1	
311	001216	001700	MOV R1,R0	
312	001222	004767	JSR PC,P2,R1	
313	001226	001601	MOV (SP)+,R1	
314	001230	000207	RTS PC	RETURN FROM BINP01
315			*****	
316			*****	
317			*****	
318	001232		DECODE: CONVERT ASCII CHARACTERS TO BIT LINKAGE: PC	
319			*****	
320			*****	
321			*****	
322	001232	010000	MOV R0,-(SP)	
323	001234	010346	MOV R3,-(SP)	
324	001230	000001	CLR R1	PREPARE BUFFER
325	001240	000303	CLR R3	SHIFT OUT
326	001242	000506	CLR R0	STAMP TO BE STACK
327	001204	102710	MOV R0,-(SP)	STORE CHARACTER
328	001250	101007	CHP R3,R0,R1	CHARACTER IN ALPHABETIC
329	001254	100425	RTI DECODE	
330	001250	122710	CHP R2,R1,(R)	
331	001262	100402	RTI DECODE	
332	001264	112006	MOV R0,-(SP)	
333	001266	100066	RTI DECODE	
334	001270	000716	DECODE: JST (SP)	IS STAMP OF STACK?
335	001272	001416	REG DECODE	YES
336	001274	102716	MOV R2,R1	PREPARE BUFFER
337	001300	011000	MOV R0,-(SP)	
338	001302	010030	MOV R0,-(SP)	
339	001306	010537	MOV R0,-(SP)	
340	001312	010700	MOV R0,-(SP)	
341	001316	000001	RTS PC	
342	001320	000703	RTS PC	SHIFT 1 TIMES LEFT
343	001324	000720	RTS PC	SHIFT 1 TIMES LEFT

```

344 001326 000760
345 001330 005726
346 001332 012603
347 001334 012600
348 001336 000207
349
351
352
353 001340
354
355
356 001340 010046
357 001342 010146
358 001344 002715 000000
359 001346 112567 000000
360 001348 005205
361 001350 013700 000052
362 001362 012701 000000
363 001364 004767 000010
364 001372 012700 000000
365 001376 004767 177006
366 001402 012601
367 001404 012600
368 001406 000205
369
370
371 001410
372
373
374 001410 010246
375 001412 012702 000000
376 001416 010006
377 001420 004716 177170
378 001424 002716 000000
379 001430 000241
380 001432 006000
381 001434 006200
382 001436 006200
383 001440 005302
384 001442 001365
385 001444 012702 000000
386 001450 112621
387 001452 005302
388 001454 001375
389 001456 012602
390 001460 000207
391
392 001462
393
394
395 001462 010146
396 001464 010246
397 001466 005001
398 001470 005002
399 001472
400 001476 002700 000000
401 001502 122700 000000

```

```

      BR DECOD2
DECOD3: TST (SP)+
        MOV (SP)+,R3
        MOV (SP)+,R0
        RTS PC          ;RETURN FROM DECOD

      I
      I*****
      IGERH: ICREATE TEXT FOR I/O ERRORS ON CHANNEL I
      I***** ICH=NO=1ST ARG., LINKAGE=RS
      I
        MOV R0,-(SP)
        MOV R1,-(SP)
        ADD #0,(S)
        MOVH (S)+,IOER2    ;GET CH=NO
        INC R5             ;CORRECT LINKAGE REG.
        MOV #IOER3,R1      ;GET I/O ERROR WORD
        JSR PC,PUTWHD      ;STARTING ADDR. OF BUFFER
        MOV #IOER1,R4
        JSR PC,PRINT       ;PRINT ERROR MESSAGE
        MOV (SP)+,R1
        MOV (SP)+,R2
        RTS R5             ;RETURN FROM IGERH

      I
      I*****
      PUTWHD: IPUT WM CONV. TO 6 DIGIT ASCII IN LOC
      I***** ISTARTING AT W1; LINKAGE=PC
      I
        MOV R2,-(SP)
        MOV #0,R2          ;6 DIGITS
        MOV R0,-(SP)       ;STACK THEM DOWN
        RLC #17770,(SP)
        ADD #0,(SP)
        CLC
        ROR R0
        ASH R0
        ASH R0
        DEC R2
        BNE PUTW01
        MOV #0,R2          ;WORD REVERSED SEQUENCE
        PUTW02: MOVH (SP)+,(1)+
        DEC R2
        RNE PUTW02
        MOV (SP)+,R2
        RTS PC             ;RETURN FROM PUTWHD

      I*****
      CHDIN: ICMU INPUT; LINKAGE=RS
      I***** IRESULT IN R0;-1=ERR, 0=OK, 1=CHD(I)
      I
        MOV R1,-(SP)
        MOV R2,-(SP)
        CLR R1
        CLR R2
        CHDIN1: ITEXT
        BIC #200,R0
        CFP0 #00,R0

```

Light copy

LOGIO:11/0 ROUTINE FOR LOG RT-11 MACRO VM02-09 00114:51 PAGE 1*

```

402 001500 001771      REG CHOIN1
403 001510 122700      CHPB #15,R0
404 001514 001502      BNE CHOIN2
405 001510 005202      INC R2
406 001520 000412      BK CHOIN4
407 001522 122715      CHDIN2: CHPB #0,(S)
408 001526 001500      REG CHOIN3
409 001530 005201      INC R1
410 001532 120325      CHPB #0,(S)+
411 001534 001404      REG CHOIN4
412 001536 000771      BK CHOIN2
413 001540 010791      CHDIN3: MOV #1,R1
414 001542 005205      CHDIN4: TSTB (S)
415 001544 000404      BK CHOIN5
416 001546 105715      INC R5
417 001548 001422      BK CHOIN6
418 001550 000774      INC R5
419 001552 005205      CHDIN5: BIT #1,R5
420 001554 000774      INC R5
421 001556 005205      CHDIN6: INC R5
422 001558 005205      TST R2
423 001560 005702      BNE CHOIN4
424 001562 001701      TTYIN
425 001564 005205      INC R5
426 001566 005702      INC R5
427 001568 122710      CHPB #15,R0
428 001570 001371      BNE CHOIN7
429 001572 001407      TTYIN
430 001574 002700      INC R1,R2
431 001576 001602      MOV (SP)+,R2
432 001578 012001      MOV (SP)+,R1
433 001580 005700      TST R0
434 001582 100401      BMT CHOIN9
435 001584 000205      R15 R5
436 001586 012700      CHDIN9: MOV #0,SP,R0,R2
437 001588 000767      JSR PC,PRIN1
438 001590 012700      MOV #1,R2
439 001592 000205      R15 R5
440      RETUM FROM CHOIN
441      ENH

```

```

1
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8
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10
11
12
13
14
15
16 000003
17 000004
18 000005
19
20
21
22
23
24
25 000006
26
27 000007
28 000008
29
30
31 000009
32
33 000010
34 000011 105737 177564
35 000012 105737
36 000013 000014
37
38
39 000015
40
41 000016 010007 001656
42 000017 010008 001706
43 000018 010009
44 000019 010010
45 000020 010011
46 000021 010012
47 000022 010013
48 000023 010014
49
50
51 000024
52
53 000025 010015 001706
54 000026 010016
55 000027 010017
56 000028 010018
57 000029 010019

```

```

      .TITLE LOGGRF: GRAPHICS FOR LOGGING
      *****
      *
      * ROUTINES FOR GRAPHIC DISPLAY OF 3-40K EVENTS IN "H1-05"
      * A HISTOGRAM IS DRAWN ON THE TEXTMODEX 40X25 SCREEN.
      * EVENTS ARE ENTERED ALONG THE X-AXIS, THEIR FREQUENCY IS
      * PLOTTED ON THE Y-AXIS.
      * UPON OVERFLOW, THE PICTURE IS RESC-LED BY FACTOR 2.
      *
      *****
      *
      .MCALL .REGDEF,.TTYIN,.TTYOUT,.DEF
      .GLOBAL DECOUT,DCOUTV,ATEXT,PRINT
      .GLOBAL GRIM1,PICT1,DRAW1
      .DEFDEF
      .DEF
      STACK=SP

      .TITLE PLOTTING I/O ROUTINES
      *****
      *
      *****
      CHIN:  INPAD AN 8 BIT CHAR FROM TEXTMODEX
      *****
      *
      .TTYIN
      RTS R5

      *****
      CHOUT:  IPUT AN 8 BIT CHAR. TO TEXTMODEX
      *****
      *
      .TTYOUT
      TSTB #TP0
      RFL ,+4
      RIS R5

      *****
      SAVE:  ;SAVE REGISTERS
      *****
      *
      MOV RC,SAVER
      MOV #SAVER+2,R0
      MOV R1,(R)+
      MOV R2,(R)+
      MOV R3,(R)+
      MOV R4,(R)+
      MOV R5,(R)+
      RIS PC

      *****
      UNSAVE: ;RESTORE REGISTERS
      *****
      *
      MOV #SAVER+2,R7
      MOV (R)+,R1
      MOV (R)+,R2
      MOV (R)+,R3
      MOV (R)+,R4

```

```

58 000062 012005      MOV (R0)+,R5
59 000064 016700 001614  MOV SAVAR,R0
60 000070 000207      RTS PC
61
62
63
64
65
66
67
68
69
70
71
72
73
74
75
76
77
78
79
80
81
82 000072 010367 001622  TPL0T:  MOV      R0,TPTMOD
83 000076 010366      MOV      R0,-(SP)
84 000100 001425      BEQ      TPTDV
85 000102 100010      BPL      TPTNRM
86 000104 012700 000037  MOV      #037,R0
87 000110 004567 177072  JSR      R0,CROUT
88
89 000114 012700 000035  TPTDV:  MOV      #035,R0
90 000120 004567 177062  JSR      R0,CROUT
91
92 000124 012500      TPTNRM: MOV      (R0)+,R0
93 000126 100001      BPL      TPT10
94 000130 005000      CLW      R0
95 000132 020027 002000  TPT10:  CMP      R0,#124
96 000136 100402      BHI      TPT12
97 000140 012700 001777  MOV      #125,R0
98 000144 010007 001552  TPT12:  MOV      R0,TPT1
99
100 000150 012500      MOV      (R0)+,R0
101 000152 100001      BPL      TPT14
102 000154 005000      CLW      R0
103 000156 020027 001415  TPT14:  CMP      R0,#781
104 000162 100402      BHI      TPT16
105 000164 012700 001410  MOV      #120,R0
106 000170 010007 001530  TPT16:  MOV      R0,TPT1
107
108 000174 016700 001524  TPTPNT: MOV      TPTV,R0
109 000200 006100      ROL      R0
110 000202 006100      ROL      R0
111 000204 006100      ROL      R0
112 000206 006100      SHAR     R0
113 000210 002700 177740  BIC      #177740,R0
114 000214 002700 000040  BLS      #000040,R0

```

MOV (R0)+,R5
MOV SAVAR,R0
RTS PC

TPL0T

THIS ROUTINE IS CALLED TO PLOT
IN VECTOR, OR POINT
PLOT MODE DEPENDING ON THE
VALUE OF REG R AS DESCRIBED BELOW.

IF
R=0 INITIALIZE AND DARK VECTOR TO X,Y

R=1 BRIGHT VECTOR TO X,Y

R=2 POINT PLOT TO X,Y

CHECK REG 2, TPTMOD
SAVE R0 ON STACK
JUMP IF INIT. AND DARK VECT
JUMP IF NORMAL VECTOR

OUTPUT 4 GS TO INITIALIZE

MOV X COORD TO REG 2
JUMP IF GEN V

IF REG SET TO 0
CHECK FOR 0 SCREEN
JUMP IF IN RANGE
SET TO EDGE IF TOO HIGH
SAVE X VALUE

GET Y COORD
JUMP IF GEN V

CHECK FOR TOO LARGE Y
JUMP IF IN RANGE
MOVE TO EDGE OF SCREEN
SAVE Y VALUE

GET Y VALUE
MOVE UPPER 5 BITS
TO UPPER BYTE

SHAR UPPER AND LOWER BYTE
THAR OFF EXTRA
SET LIGHT Y TAG

PLOTTING I/O ROUTINES

115 000220	000567	177562	JSR	WS,CHOUT	WS,CHOUT HI Y
116 000224	001600	001474	MOV	TPY,R0	SET Y CORNO
117 000232	002700	177740	BIC	#177740,R0	MASK IN LOW 5 BITS
118 000234	002700	002140	BIS	#002140,R0	AND SET LOW 5 BITS
119 000240	000567	177562	JSR	WS,CHOUT	SWAP OUT LOW 5 BITS
120					
121 000244	001600	001452	MOV	TPY,R0	SET X CORNO
122 000250	000100		ROL	R0	AND AUGUST LIKE Y
123 000252	000100		ROL	R0	
124 000254	000100		ROL	R0	
125 000256	000300		SWAB		SWITCH BYTES
126 000260	002700	177740	BIC	#177740,R0	MASK OFF L2P4
127 000264	002700	000040	BIS	#000040,R0	SET HI 4 L2P4
128 000270	000567	177512	JSR	WS,CHOUT	OUTPUT HI 4 BYTES
129 000274	001600	001422	MOV	TPY,R0	SET X CORNO
130 000300	002700	177740	BIC	#177740,R0	LEAVE ONLY LOW 5 BITS
131 000304	002700	000100	BIS	#000100,R0	SET IN LOW 5 BITS
132 000310	000567	177472	JSR	WS,CHOUT	OUTPUT LOW 5 BITS
133 000314	005167	001400	TST	TPY,R0	CHECK FOR POINT PLOT
134 000320	100003		CLH		RETURN
135 000322	005467	001372	CLH	TPY,R0	CLEAR AND RETURN VECTOR
136 000326	000722		BR	TPY-T	
137					
138 000330	012600		MOV	(SP)+,R0	RESTORE MW AND EXIT
139 000332	000205		RTS	R0	RETURN FROM PLOT
140					

```

GRAPH ROUTINES      RT-11 MACRO VM02-09      00109:04 PAGE 2

1
2
3
4
5 000334
6
7
8 000334      010000
9 000330      010100
10 000340      010200
11 000342      010300
12 000344      012703
13 000350      012700
14 000354      012701
15 000360      012702
16 000364      010100
17 000366      005023
18 000370      002701
19 000374      005302
20 000376      001372
21
22 000400      012707
23 000406      005067
24 000412      012767
25 000420      012700
26 000424      012701
27 000430      012702
28 000434      010220
29 000436      102702
30 000442      005301
31 000444      001373
32 000446      005010
33 000450      012603
34 000452      012602
35 000454      012601
36 000456      012600
37 000460      000207
38
39
40 000462
41
42
43 000462      000767
44 000466      012700
45 000470      004767
46 000476      005003
47 000480      012700
48 000484      004767
49 000486      012701
50 000488      016300
51 000490      004767
52 000492      005723
53 000494      020327
54 000496      001411
55 000500      012700
56 000504      004767
57 000506      016300

.TITLE GRAPH ROUTINES
*****
GRINM1: INITIALIZE FOR GRAPH MODE 1
*****
MOV R0,-(SP)
MOV R1,-(SP)
MOV R2,-(SP)
MOV R3,-(SP)
MOV R4PT,R3
MOV R4PT,M4
MOV R4PT,R1
MOV R20,R2
MOV R1,(R)+
CLR (R)+
ADD #42,(R)
DEC R2
BNE GRINM1

MOV R0,DIVISX INITIALLY 6 LINES PER EVENT
CLR TIFM 12 COLS. IN GRAPH
MOV M4,WIDTH
MOV VTEXT,M4
MOV M4,M1
MOV M5,M2
MOV R2,(R)+
SUB #14,(R)
DEC R1
RUE GTHLS
CLR (R)
MOV (SP),R3
MOV (SP),R2
MOV (SP),R1
MOV (SP),M4
RTS PC
RETURN FROM GRINM1 J

INITIALIZE GRAPH FOR MODE 17 LINKAGE? PC
JURN PICTURE AS IN XPT AND YPT

JSK PL,SAVE
MOV PSCHER,R1
JSK PL,PRIU
CLR R3
MOV R04,R2
JSK PC,PRINT
MOV M6,M1
MOV VTEXT(R3),M3
JSK PL,XLOCUT
TST (R).
CNP M3,M1
REQ PIC115
MOV MXT1,M4
JSK PC,PRINT
MOV VTEXT(R3),M3

PSET XPT TO STARTING VAL.
:POR COL. AND CLEAR YPT

PSET VALUES FOR Y-AXIS
Y VALUE DESCRIPTION
PLOTTER ALWAYS 3
PFTURN FROM GRINM1 J

FINISHED?
YES
FIELD LENGTH FOR DECUIT
PRINT TEST FOR Y-AXIS

```

GRAPH ROUTINES RT-11 MACRO VM82-89 00109104 PAGE 2*

```

50 000550 000767 0000000  JSR PC,DECODE
51 000554 000763 0000000  BR PICT14
52 000558 012100 0001200  PRINT TEXT FOR X-AXIS
53 000562 000767 0000000  JSR PC,PRINT
54 000566 016703 0001374  MOV ITEM,M3
55 000570 005703 0000000  CLR R2
56 000574 005703 0000000  IFINISHED?
57 000578 001410 0000000  JFCS
58 000582 016200 0001474  MOV ECTAR(P),R0
59 000586 000547 0000000  JSR P5,OCOUTV
60 000590 000002 0000000  *WORD 2
61 000594 005743 0000000  TST -(3)
62 000598 005743 0000000  TST (2)+
63 000602 005743 0000000  BR PICT14
64 000606 012700 0001495  MOV *CLRF,R0
65 000610 005767 0000000  JSR PC,PRINT
66 000614 012700 0001731  MOV *B63,M0
67 000618 005767 0000000  JSR PC,PRINT
68 000622 012700 0000000  MOV *TEXT,M0
69 000626 005767 0000000  JSR PC,PRINT
70 000630 005767 0000000  CLR R0
71 000634 005767 0000000  JSR M5,TPLOT
72 000638 005767 0000000  *WORD 170
73 000642 005767 0000000  *WORD 170
74 000646 005767 0000000  INC R0
75 000650 005767 0000000  JSR M5,TPLOT
76 000654 005767 0000000  *WORD 170
77 000658 005767 0000000  *WORD 329
78 000662 005767 0000000  JSR M5,TPLOT
79 000666 005767 0000000  *WORD 1220
80 000670 005767 0000000  *WORD 329
81 000674 005767 0000000  CLR R1
82 000678 005767 0000000  PICT11: CPP ITEM,M1
83 000682 005767 0000000  BR PICT13
84 000686 005767 0000000  MOV *PT(1),M3
85 000690 005767 0000000  MOV *PT(1),PICT1A
86 000694 005767 0000000  MOV *PT(1),PICT1B
87 000698 005767 0000000  MOV *PT(1),PICT1C
88 000702 005767 0000000  MOV *PT(1),PICT1D
89 000706 005767 0000000  MOV *PT(1),PICT1E
90 000710 005767 0000000  MOV *PT(1),PICT1F
91 000714 005767 0000000  MOV *PT(1),PICT1G
92 000718 005767 0000000  MOV *PT(1),PICT1H
93 000722 005767 0000000  MOV *PT(1),PICT1I
94 000726 005767 0000000  MOV *PT(1),PICT1J
95 000730 005767 0000000  MOV *PT(1),PICT1K
96 000734 005767 0000000  MOV *PT(1),PICT1L
97 000738 005767 0000000  MOV *PT(1),PICT1M
98 000742 005767 0000000  MOV *PT(1),PICT1N
99 000746 005767 0000000  MOV *PT(1),PICT1O
100 000750 005767 0000000  JSR M5,TPLOT
101 000754 005767 0000000  PICT1A: *WORD 0
102 000758 005767 0000000  PICT1B: *WORD 0
103 000762 005767 0000000  PICT1C: *WORD 0
104 000766 005767 0000000  PICT1D: *WORD 0
105 000770 005767 0000000  PICT1E: *WORD 0
106 000774 005767 0000000  PICT1F: *WORD 0
107 000778 005767 0000000  PICT1G: *WORD 0
108 000782 005767 0000000  PICT1H: *WORD 0
109 000786 005767 0000000  PICT1I: *WORD 0
110 000790 005767 0000000  PICT1J: *WORD 0
111 000794 005767 0000000  PICT1K: *WORD 0
112 000798 005767 0000000  PICT1L: *WORD 0
113 000802 005767 0000000  PICT1M: *WORD 0
114 000806 005767 0000000  PICT1N: *WORD 0
115 000810 005767 0000000  PICT1O: *WORD 0
116 000814 005767 0000000  JSR M5,TPLOT

```

PRINT TEXT FOR X-AXIS

JNR OF EV

IFINISHED?

JFCS

PRINT EC

LENGTH 2

PRINT NAME OF X-AXIS

ISET GRAPH MODE

ISFT LIGHT VECTOR

IURAN Y-AXIS

IURAN X-AXIS

ICNT = OF EXISTING EVENTS

CALL EV COL, DRAWN.

IYFS

IURAN TEXT

IURAN OF LINES

ISET ARGUMENTS

IPOSITION VECTOR

IINC X-COORD.

IDEC LINE CNT

ISET LINE

ISET EV CNT

ISET COL

ISET VECTOR TO RETURN LOC.

LINE	ADDRESS	INSTR	OPERAND	COMMENT
115	001032	000000		
116	001034	000362		
117	001036	012708		
118	001042	004767		
119	001046	176774		
120	001052	000207		
121				
122	001034			
123				
124				
125				
126	001054	010046		
127	001056	010176		
128	001068	010246		
129	001062	010546		
130	001064	005041		
131	001066	026701		
132	001072	001405		
133	001074	020261		
134	001100	001144		
135	001102	005721		
136	001104	007770		
137	001106	007127		
138	001112	100466		
139	001116	012700		
140	001120	004767		
141	001124	000167		
142				
143	001130	007767		
144	001136	010761		
145				
146	001142	016102		
147	001146	102702		
148	001152	010267		
149	001156	005800		
150	001160	004567		
151	001164	000200		
152	001168	000061		
153	001174	012700		
154	001174	000267		
155	001240	016100		
156	001234	004567		
157	001210	000302		
158				
159	001212	026127		
160	001220	100464		
161	001272	016102		
162	001226	016167		
163	001234	016167		
164	001282	006767		
165	001254	012767		
166	001256	012767		
167	001264	006167		
168	001272	006167		
169	001304	005700		
170	001306	004567		
171	001308	000362		

DATA SECTION RT-11 MACRO VM82-09 00:09:04 PAGE 3+

```

002111 000 105 126 XTEXT: .ASCII / EVENT /<200>
26 002112 000 124
002115 105 000
002120 000
002123 000
27 002124 012 015 GAP: .ASCII: <12><15> /<200>
002127 000 000
002130 000 000
002133 000 000
002136 000 000
002139 000 000
002142 000 000
002145 000 000
28 002146 015 012 CRLF: .BYTE 15,12
29 002147 037 200 ALPHA: .BYTE 17,200
30 .EVEN
31 .TITLE LOGGRF
32 .END
33
000001

```

3. LOGPRT - Printing Routine for Offline Output of HIMOS Reports

209

```

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8
9
10
11
12
13
14 000000
15 000000
16 177564
17 177564
18 000066
19 000000
20 001000
21 001000 012706 001000
22 001004
23 001016
24 001034 103761
25 001036 012705 000066
26 001042 004767 000114
27 001000 005002
28
29 001050
30 001114 103004
31 001116 005737 000052
32 001122 001411
33 001124 000725
34 001126 012703 001000
35 001132 012701 001104
36 001136 004767 000106
37 001142 005272
38 001144 007741
39
40 001146
41 001160
42
43
44 001162
45
46 001162 012702 000160
47 001166
48 001172 002700 000200
49 001176 110022
50 001200 122700 000012
51 001204 001370
52 001206 105012
53 001210 012700 000156
54 001214 004767 000002
55 001220 000207
56
57 001222

          .TITLE LOGPRT : HIMOS OUTPUT INTERPRETER
          *****
          *
          * THIS PROGRAM PROCESSES "HIMOS" LOG FILES
          * THE FILES ARE TRANSFORMED INTO READABLE FORM AND
          * PRINTED ON THE SYSTEM TERMINAL, ONE RECORD PER LINE.
          * THE FIRST VALUE IS THE EC, THE REMAINING ITEMS ARE
          * EXPLAINED BY A LEADING 2 LETTER ID. CODE.
          *
          *****
          *
          .MCALL .PRINT,.HFGDEF,.EXIT,.HFA2,...VP..
          .MCALL .CSIGEN,.CLOSE,.TTYOUT,.TTYIN,.EXIT
          .VE..
          .HFGDEF
          TTYOUT=177560
          TTYIN=177560
          LINES=54.
          .ASCT
          .E1000
          START: MOV #1,SP
                  .CLOSE #3
                  .CSIGEN #DSKHHND,#EXT,#0
                  RCS START          OPEN ERROR,TRY AGAIN
                  MOV #LINES,#5      SET LINE CNT
                  JSR PC,GFILIN
                  CLR R2              BLOCK CNT
          *
          RDBLK: .HFA2 #AREA,#3,#BUFF,#256,,R2  READ ONE BLOCK
                  RCC R00K          INPAD IN
                  TST #R2            EOD?
                  BEQ TERM          EIG-TERMINATE
                  RR START          INQ-HANDY ERROR
          RDOCK: MOV #512,,R3        512 BYTES
                  MOV #BUFF,R1      STARTING AT BUFF
                  JSR PC,HPRINTB     PRINT AND DECODE
                  INC R2             INPAD CNT
                  PR RDBLK
          *
          TERM:  .CLOSE #3
                  .EXIT
          *
          *****
          GETLIN: READ HEADING LINE
          *****
                  MOV #LINHHF,R2
          GETLI:  .TTYIN
                  RLC #256,R0
                  MOVB #R,(2)+
                  CHPB #12,R0
                  BNE GETLI
                  CLRB (2)
                  MOV #CLF,R0
                  JSR PC,HPRINT
                  RTS PC              RETURN FROM GETLIN
          *****
          PRINT: .PRINT TEXT STARTING AT (R0)
    
```

LOGPRT : MINOS OUTPUT INTERPRET RT-11 MACRO VM02-09 02:10:25 PAGE 14

```

54 001222 010146      *****      MOV R1, -(SP)
55 001223 010201      PRINT0:      MOV R3, R1
56 001224 112102      MOV R1, (R3)+, R0
57 001225 177564      TTYOUT
58 001226 105737      TSTB #RTTYOUT
59 001227 100375      BPL -4
60 001228 105737      TSTB R0
61 001229 000200      BEQ PRINT1
62 001230 122733      CMPB #200, R4
63 001231 001411      BEQ PRINT2
64 001232 001411      RM PRINT4
65 001233 000015      PRINT1: MOV R15, R3
66 001234 112740      TTYOUT
67 001235 000012      MOV R12, R0
68 001236 105737      TTYOUT
69 001237 140375      PRINT2: TSTB #RTTYOUT
70 001238 012601      BPL -4
71 001239 000227      MOV (SP)+, R1
72 001240 000227      RTS PC
73 001241 000227      *****
74 001242 000227      PRINT0: PRINT R3 CHARS, STARTING AT (R1)
75 001243 000227      *****
76 001244 000227      PRINT1:
77 001245 000227      *****
78 001246 000227      PRINT2:
79 001247 000227      *****
80 001248 000227      PRINT3:
81 001249 000227      *****
82 001250 000227      PRINT4:
83 001251 000227      *****
84 001252 000227      PRINT5:
85 001253 000227      *****
86 001254 000227      PRINT6:
87 001255 000227      *****
88 001256 000227      PRINT7:
89 001257 000227      *****
90 001258 000227      PRINT8:
91 001259 000227      *****
92 001260 000227      PRINT9:
93 001261 000227      *****
94 001262 000227      PRINT10:
95 001263 000227      *****
96 001264 000227      PRINT11:
97 001265 000227      *****
98 001266 000227      PRINT12:
99 001267 000227      *****
100 001268 000227      PRINT13:
101 001269 000227      *****
102 001270 000227      PRINT14:
103 001271 000227      *****
104 001272 000227      PRINT15:
105 001273 000227      *****
106 001274 000227      PRINT16:
107 001275 000227      *****
108 001276 000227      PRINT17:
109 001277 000227      *****
110 001278 000227      PRINT18:
111 001279 000227      *****
112 001280 000227      PRINT19:
113 001281 000227      *****
114 001282 000227      PRINT20:
115 001283 000227      *****
116 001284 000227      PRINT21:
117 001285 000227      *****
118 001286 000227      PRINT22:
119 001287 000227      *****
120 001288 000227      PRINT23:
121 001289 000227      *****
122 001290 000227      PRINT24:
123 001291 000227      *****
124 001292 000227      PRINT25:
125 001293 000227      *****
126 001294 000227      PRINT26:
127 001295 000227      *****
128 001296 000227      PRINT27:
129 001297 000227      *****
130 001298 000227      PRINT28:
131 001299 000227      *****
132 001300 000227      PRINT29:
133 001301 000227      *****
134 001302 000227      PRINT30:
135 001303 000227      *****
136 001304 000227      PRINT31:
137 001305 000227      *****
138 001306 000227      PRINT32:
139 001307 000227      *****
140 001308 000227      PRINT33:
141 001309 000227      *****
142 001310 000227      PRINT34:
143 001311 000227      *****
144 001312 000227      PRINT35:
145 001313 000227      *****
146 001314 000227      PRINT36:
147 001315 000227      *****
148 001316 000227      PRINT37:
149 001317 000227      *****
150 001318 000227      PRINT38:
151 001319 000227      *****
152 001320 000227      PRINT39:
153 001321 000227      *****
154 001322 000227      PRINT40:
155 001323 000227      *****
156 001324 000227      PRINT41:
157 001325 000227      *****
158 001326 000227      PRINT42:
159 001327 000227      *****
160 001328 000227      PRINT43:
161 001329 000227      *****
162 001330 000227      PRINT44:
163 001331 000227      *****
164 001332 000227      PRINT45:
165 001333 000227      *****
166 001334 000227      PRINT46:
167 001335 000227      *****
168 001336 000227      PRINT47:
169 001337 000227      *****
170 001338 000227      PRINT48:
171 001339 000227      *****
172 001340 000227      PRINT49:
173 001341 000227      *****
174 001342 000227      PRINT50:
175 001343 000227      *****
176 001344 000227      PRINT51:
177 001345 000227      *****
178 001346 000227      PRINT52:
179 001347 000227      *****
180 001348 000227      PRINT53:
181 001349 000227      *****
182 001350 000227      PRINT54:
183 001351 000227      *****
184 001352 000227      PRINT55:
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186 001354 000227      PRINT56:
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188 001356 000227      PRINT57:
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190 001358 000227      PRINT58:
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192 001360 000227      PRINT59:
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194 001362 000227      PRINT60:
195 001363 000227      *****
196 001364 000227      PRINT61:
197 001365 000227      *****
198 001366 000227      PRINT62:
199 001367 000227      *****
200 001368 000227      PRINT63:
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202 001370 000227      PRINT64:
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204 001372 000227      PRINT65:
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218 001386 000227      PRINT72:
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226 001394 000227      PRINT76:
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230 001398 000227      PRINT78:
231 001399 000227      *****
232 001400 000227      PRINT79:
233 001401 000227      *****
234 001402 000227      PRINT80:
235 001403 000227      *****
236 001404 000227      PRINT81:
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238 001406 000227      PRINT82:
239 001407 000227      *****
240 001408 000227      PRINT83:
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242 001410 000227      PRINT84:
243 001411 000227      *****
244 001412 000227      PRINT85:
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246 001414 000227      PRINT86:
247 001415 000227      *****
248 001416 000227      PRINT87:
249 001417 000227      *****
250 001418 000227      PRINT88:
251 001419 000227      *****
252 001420 000227      PRINT89:
253 001421 000227      *****
254 001422 000227      PRINT90:
255 001423 000227      *****
256 001424 000227      PRINT91:
257 001425 000227      *****
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264 001432 000227      PRINT95:
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266 001434 000227      PRINT96:
267 001435 000227      *****
268 001436 000227      PRINT97:
269 001437 000227      *****
270 001438 000227      PRINT98:
271 001439 000227      *****
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273 001441 000227      *****
274 001442 000227      PRINT100:
275 001443 000227      *****
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277 001445 000227      *****
278 001446 000227      PRINT102:
279 001447 000227      *****
280 001448 000227      PRINT103:
281 001449 000227      *****
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283 001451 000227      *****
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291 001459 000227      *****
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532 001700 000227      PRINT229:
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550 001718 000227      PRINT238:
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562 001730 000227      PRINT244:
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564 001732 000227      PRINT245:
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566 001734 000227      PRINT246:
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598 001766 000227      PRINT262:
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602 001770 000227      PRINT264:
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606 001774 000227      PRINT266:
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610 001778 000227      PRINT268:
611 001779 000227      *****
612 001780 000227      PRINT269:
613 001781 000227      *****
614 001782 000227      PRINT270
```

LOGPR1 : MINOS OUTPUT INTERPRET RT-11 MACRO VM02-09 00110:25 PAGE 1*

```

115 001450 012600
116 001452 000207
117
118 001454
119 001114 000537
120
121 001124
122 001144
123 001144
124 001144
125 001151 012
126 001151 012
127 001151 012
128 001151 012
129 001151 012
130 001151 012
131 001151 012
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149 001151 012
150 001151 012
151 001151 012
152 001151 012
153 001151 012
154 001151 012
155 001151 012

```

MOV (SP)+,R0
RTS PC

DSKIND: .BLKW 000
EXT: .MADSO "LUG"
.BLKW 3
AREA: .BLKW 10
BUFF: .BLKW 256
HEAD: .BYTE 15,12,12,12,12
.BYTE 12,12,12,12,12

CRLF: .WORD 0
LINBUF: .BLKW 80

NAME INDEX
IDADD: IDER 10 *
IDPC 12 PL
IDPS 14 PS
IDSP 16 SP
IDCP 18 CP
IDFW 112 FW
IDSM 114 SM
IDMF 116 MF
IDNR 120 NR
IDGR 122 GR
IDMT 124 MT
IDST 126 ST

IDER: .BYTE 52,200
IDPC: .ASCII /PC/<200>
IDPS: .ASCII /PS/<200>
IDSP: .ASCII /SP/<200>
IDCP: .ASCII /CP/<200>
IDFW: .ASCII /FW/<200>
IDSM: .ASCII /SM/<200>
IDMF: .ASCII /MF/<200>
IDNR: .ASCII /NR/<200>
IDGR: .ASCII /GR/<200>
IDMT: .ASCII /MT/<200>
IDST: .ASCII /ST/<200>
.END START

RETURN FROM PRINTER

5. Program Examples for Monitor Sessions

5.1 HISTO - Program for Compiler Analysis

```

RT-11 FORTRAN IV      VB1B-08      PAGE 001

C PROGRAM TO PRINT A DISTRIBUTION HISTOGRAM
C SAMPLE VALUES ARE READ FROM THE INPUT FILE PTM1.DAT
C THE SORTED ARRAY IS STORED IN FTR2.DAT

      LOGICAL*1 STAR(80)
      REAL ARRAY(200)
      INTEGER HIST(10)
      DATA STAR/80*0/
      DO 10 I=1,200
        READ (1,20)END=100) ARRAY(I)
      20 FORMAT (F5.2)
      10 CONTINUE
      100 ISIZE=1
      DO 120 J=1,ISIZE-1
        DO 110 K=J+1,ISIZE
          IF (ARRAY(J).GE.ARRAY(K)) GO TO 110
          TMP=ARRAY(J)
          ARRAY(J)=ARRAY(K)
          ARRAY(K)=TMP
        110 CONTINUE
      120 CONTINUE
      DO 125 K=4,ISIZE
        125 WRITE(2,20)ARRAY(K)
      DO 126 K=1,10
        126 HIST(K)=0
      DO 152 K=1,ISIZE
        N=IFIX(ARRAY(K)/10+1)
        HIST(N)=HIST(N)+1
      130 WRITE(1,15)
      135 FORMAT (1X,'DISTRIBUTION IN INTERVALS OF 10.00:',/)
      DO 150 N=1,100,10
        J=N-10
        WRITE (7,100)HIST(K/10),I,K
      140 FORMAT (/,1X,13,'VALUES, RANGE: ',13,' TO ',13,5)
      IF (HIST(K/10).EQ.0) GO TO 150
      WRITE(7,140) (STAR(K),N=1,HIST(K/10))
      145 FORMAT (1H*,24,82A1)
      150 CONTINUE
      155 WRITE (7,140) ISIZE
      160 FORMAT (/,1X,'TOTAL NUMBER OF SAMPLES= ',13)
      CALL EXIT
      END
0001
0002
0003
0004
0005
0006
0007
0008
0009
0010
0011
0012
0013
0014
0015
0016
0017
0018
0019
0020
0021
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0032
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0034
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0036
0037
0038
0039
0040

```

5.2 SAMGEN - Random Number Generator for FPP Performance Analysis

```

RT-11. FORTRAN IV      V019-00      PAGE 001
C PROGRAM TO GENERATE J RANDOM NUMBERS IN THE RANGE 0-99.99
C FOUR EXIT CALLS ARE INSERTED FOR MONITORING *MAN* PERFORMANCE

0001      REAL MAX,MIN
0002      CALL IRRMT
0003      CALL EXIT(0)
0004      WRITE (7,20)
0005      20 FORMAT (1X,'# OF SAMPLES: ',S)
0006      READ (5,30) J
0007      30 FORMAT (1S)
0008      MINER
0009      MAX=99.99
0010      L=0
0011      N=0
0012      DO 100 K=1,J
0013      CALL ERMT(1)
0014      X=IRAN(L,M)*(MAX-MIN)+MIN
0015      CALL ERMT(2)
0016      WRITE (1,60) X
0017      CALL ERMT(3)
0018      60 FORMAT (F5.2)
0019      100 CONTINUE
0020      CALL ERMT(4)
0021      CALL EXIT
0022      END

```

APPENDIX C

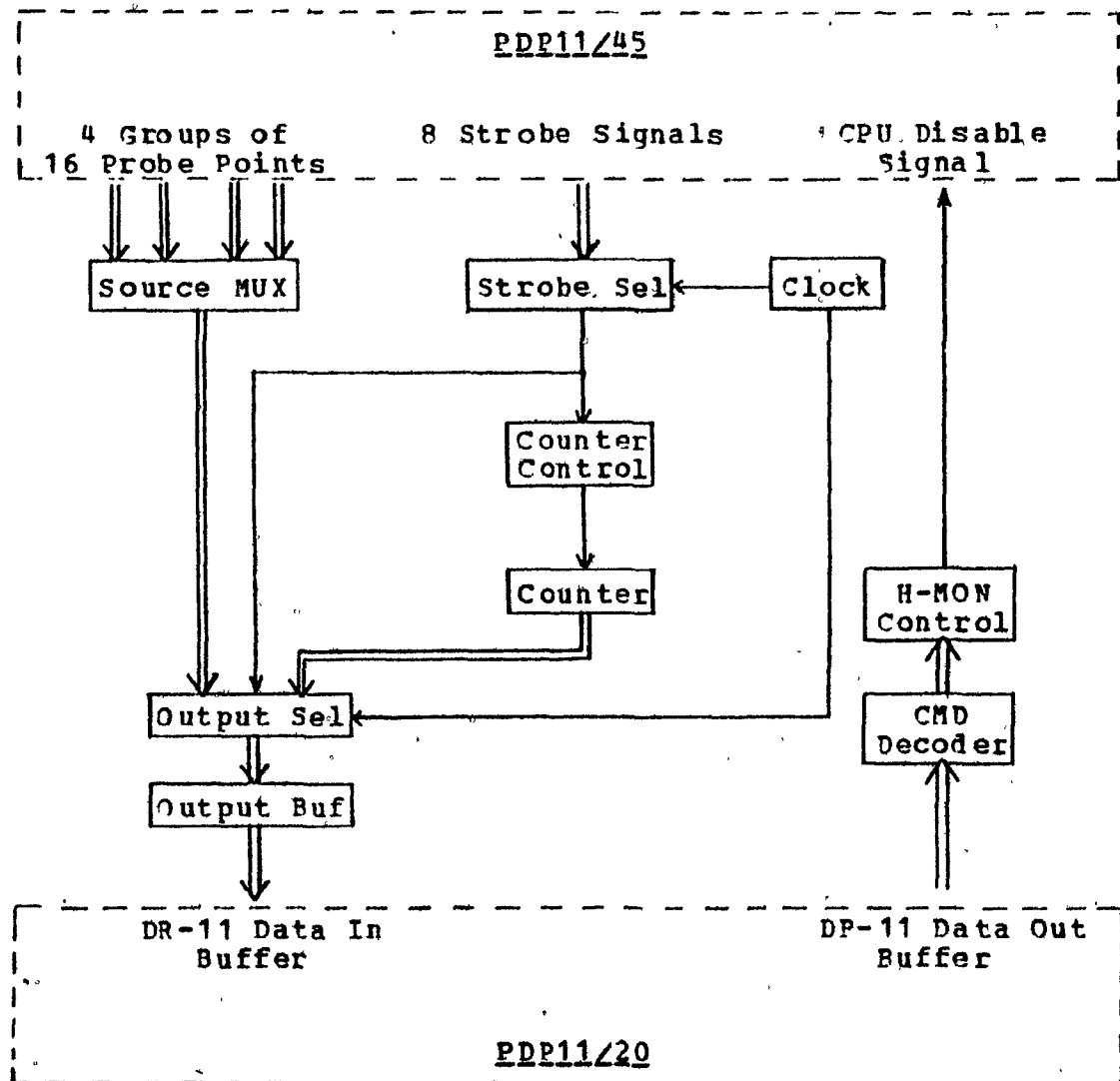
Hardware DocumentationC.1 Hardware Monitor Design

Figure C.1.1: H-MON, Data and Control

BIT											
0	2	3	4	5	7	8	-	9	11	12	15
Mode SEL <0..2>		Source SEL <0..1>		CLK SEL <0..2>		STOP/START		STROBE SEL <0..2>		RESERVED	
000 TRACE		00 SEL SOURCE 0		000 10 KHz		0 STOP		000 SEL STROBE 0			
000 COUNT		01 " " 1		001 5 KHz		1 START		001 " " 1			
010 EVENT DRIVEN		10 " " 2		010 2.5 KHz				010 " " 2			
011 TIME		01 " " 3		011 1.25 KHz				011 " " 3			
100 NOT VALID				100 500 Hz				100 " " 4			
101 NOT VALID				101 250 Hz				101 " " 5			
110 SAMPLE				110 125 Hz				110 " " 6			
111 FAST COUNT				111 62.5 Hz				111 " " 7			

Fig. C.1.2: H-MON Control Word, Bit Assignment.

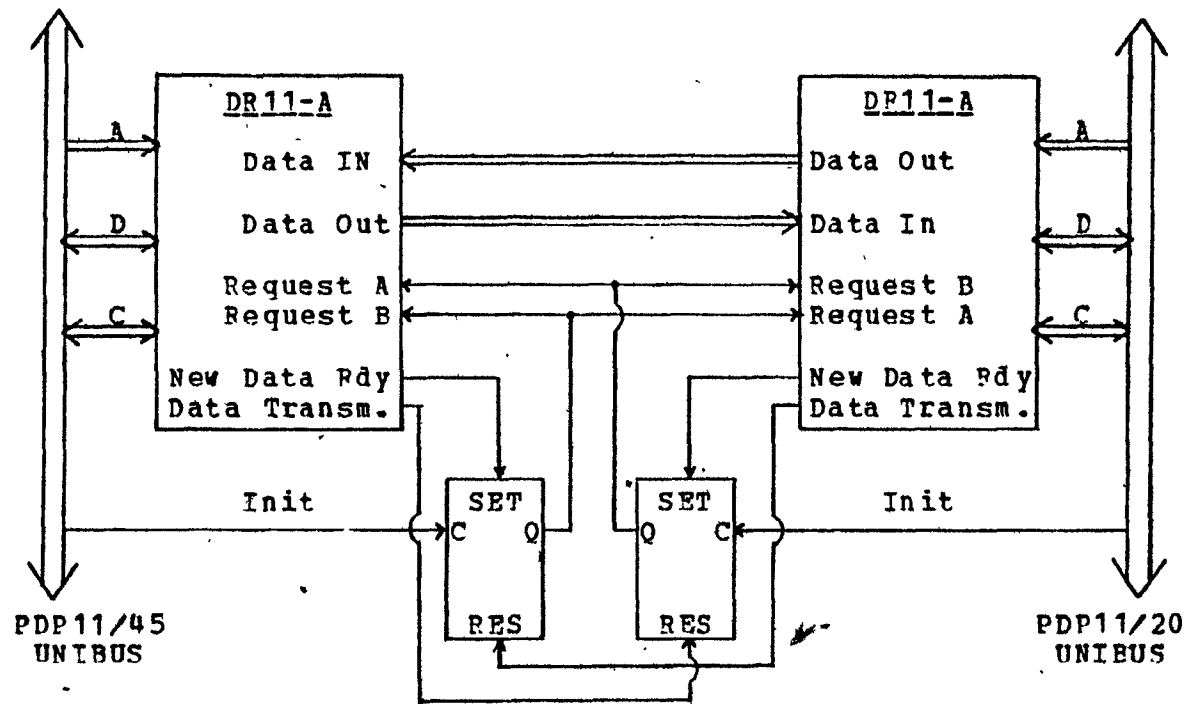
C.2 Parallel communication Link

Figure C.2.1: Communication Link

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